

INTERNATIONAL  
STANDARD

**ISO/IEC**  
**11172-1**

First edition  
1993-08-01

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**Information technology — Coding of  
moving pictures and associated audio for  
digital storage media at up to about  
1,5 Mbit/s —**

**Part 1:**  
Systems

*Technologies de l'information — Codage de l'image animée et du son  
associé pour les supports de stockage numérique jusqu'à environ  
1,5 Mbit/s —*

*Partie 1: Systèmes*



Reference number  
ISO/IEC 11172-1:1993(E)

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Printed in Switzerland.

## Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work.

In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

International Standard ISO/IEC 11172-1 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Sub-Committee SC 29, *Coded representation of audio, picture, multimedia and hypermedia information*.

ISO/IEC 11172 consists of the following parts, under the general title *Information technology — Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s*:

- Part 1: *Systems*
- Part 2: *Video*
- Part 3: *Audio*
- Part 4: *Compliance testing*

Annexes A and B of this part of ISO/IEC 11172 are for information only.

## Introduction

Note -- Readers interested in an overview of the MPEG Systems layer should read this Introduction and then proceed to annex A, before returning to the clauses 1 and 2. Since the system target decoder concept is referred to throughout both the normative and informative clauses of this part of ISO/IEC 11172, it may also be useful to refer to clause 2.4, and particularly 2.4.2, where the system target decoder is described.

The systems specification addresses the problem of combining one or more data streams from the video and audio parts of this International Standard with timing information to form a single stream. Once combined into a single stream, the data are in a form well suited to digital storage or transmission. The syntactical and semantic rules imposed by this systems specification enable synchronized playback without overflow or underflow of decoder buffers under a wide range of stream retrieval or receipt conditions. The scope of syntactical and semantic rules set forth in the systems specification differ: the syntactical rules apply to systems layer coding only, and do not extend to the compression layer coding of the video and audio specifications; by contrast, the semantic rules apply to the combined stream in its entirety.

The systems specification does not specify the architecture or implementation of encoder or decoders. However, bitstream properties do impose functional and performance requirements on encoders and decoders. For instance, encoders must meet minimum clock tolerance requirements. Notwithstanding this and other requirements, a considerable degree of freedom exists in the design and implementation of encoders and decoders.

A prototypical audio/video decoder system is depicted in figure 1 to illustrate the function of an ISO/IEC 11172 decoder. The architecture is not unique -- System Decoder functions including decoder timing control might equally well be distributed among elementary stream decoders and the Medium Specific Decoder -- but this figure is useful for discussion. The prototypical decoder design does not imply any normative requirement for the design of an ISO/IEC 11172 decoder. Indeed non-audio/video data is also allowed, but not shown.

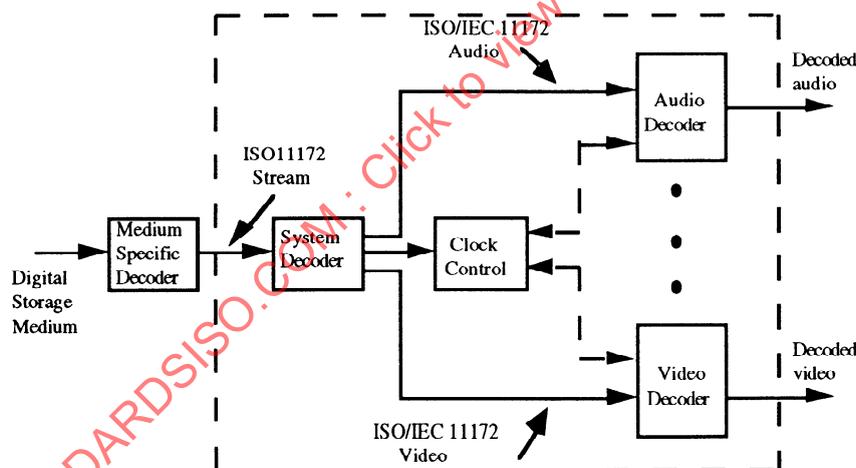


Figure 1 -- Prototypical ISO/IEC 11172 decoder

The prototypical ISO/IEC 11172 decoder shown in figure 1 is composed of System, Video, and Audio decoders conforming to Parts 1, 2, and 3, respectively, of ISO/IEC 11172. In this decoder the multiplexed coded representation of one or more audio and/or video streams is assumed to be stored on a digital storage medium (DSM), or network, in some medium-specific format. The medium specific format is not governed by this International Standard, nor is the medium-specific decoding part of the prototypical ISO/IEC 11172 decoder.

The prototypical decoder accepts as input an ISO/IEC 11172 multiplexed stream and relies on a System Decoder to extract timing information from the stream. The System Decoder demultiplexes the stream, and the elementary streams so produced serve as inputs to Video and Audio decoders, whose outputs are decoded video and audio signals. Included in the design, but not shown in the figure, is the flow of timing information among the System Decoder, the Video and Audio Decoders, and the Medium Specific Decoder.

The Video and Audio Decoders are synchronized with each other and with the DSM using this timing information.

ISO/IEC 11172 multiplexed streams are constructed in two layers: a system layer and a compression layer. The input stream to the System Decoder has a system layer wrapped about a compression layer. Input streams to the Video and Audio decoders have only the compression layer.

Operations performed by the System Decoder either apply to the entire ISO/IEC 11172 multiplexed stream ("multiplex-wide operations"), or to individual elementary streams ("stream-specific operations"). The ISO/IEC 11172 system layer is divided into two sub-layers, one for multiplex-wide operations (the pack layer), and one for stream-specific operations (the packet layer).

## 0.1 Multiplex-wide operations (pack layer)

Multiplex-wide operations include the coordination of data retrieval off the DSM, the adjustment of clocks, and the management of buffers. The tasks are intimately related. If the rate of data delivery off the DSM is controllable, then DSM delivery may be adjusted so that decoder buffers neither overflow nor underflow; but if the DSM rate is not controllable, then elementary stream decoders must slave their timing to the DSM to avoid overflow or underflow.

ISO/IEC 11172 multiplexed streams are composed of packs whose headers facilitate the above tasks. Pack headers specify intended times at which each byte is to enter the system decoder from the DSM, and this target arrival schedule serves as a reference for clock correction and buffer management. The schedule need not be followed exactly by decoders, but they must compensate for deviations about it.

An additional multiplex-wide operation is a decoder's ability to establish what resources are required to decode an ISO/IEC 11172 multiplexed stream. The first pack of each ISO/IEC 11172 multiplexed stream conveys parameters to assist decoders in this task. Included, for example, are the stream's maximum data rate and the highest number of simultaneous video channels.

## 0.2 Individual stream operations (packet layer)

The principal stream-specific operations are 1) demultiplexing, and 2) synchronizing playback of multiple elementary streams. These topics are discussed next.

### 0.2.1 Demultiplexing

On encoding, ISO/IEC 11172 multiplexed streams are formed by multiplexing elementary streams. Elementary streams may include private, reserved, and padding streams in addition to ISO/IEC 11172 audio and video streams. The streams are temporally subdivided into packets, and the packets are serialized. A packet contains coded bytes from one and only one elementary stream.

Both fixed and variable packet lengths are allowed subject to constraints in 2.4.3.3 and in 2.4.5 and 2.4.6.

On decoding, demultiplexing is required to reconstitute elementary streams from the ISO/IEC 11172 multiplexed stream. This is made possible by stream\_id codes in packet headers.

### 0.2.2 Synchronization

Synchronization among multiple streams is effected with presentation time stamps in the ISO/IEC 11172 multiplexed stream. The time stamps are in units of 90kHz. Playback of N streams is synchronized by adjusting the playback of all streams to a master time base rather than by adjusting the playback of one stream to match that of another. The master time base may be one of the N decoders' clocks, the DSM or channel clock, or it may be some external clock.

Because presentation time-stamps apply to the decoding of individual elementary streams, they reside in the packet layer. End-to-end synchronization occurs when encoders record time-stamps at capture time, when the time stamps propagate with associated coded data to decoders, and when decoders use those time-stamps to schedule presentations.

Synchronization is also possible with DSM timing time stamps in the multiplexed data stream.

### 0.2.3 Relation to compression layer

The packet layer is independent of the compression layer in some senses, but not in all. It is independent in the sense that packets need not start at compression layer start codes, as defined in parts 2 and 3. For example, a video packet may start at any byte in the video stream. However, time stamps encoded in packet headers apply to presentation times of compression layer constructs (namely, presentation units).

## 0.3 System reference decoder

Part 1 of ISO/IEC 11172 employs a "system target decoder," (STD) to provide a formalism for timing and buffering relationships. Because the STD is parameterized in terms of fields defined in ISO/IEC 11172 (for example, buffer sizes) each ISO/IEC 11172 multiplexed stream leads to its own parameterization of the STD. It is up to encoders to ensure that bitstreams they produce will play in normal speed, forward play on corresponding STDs. Physical decoders may assume that a stream plays properly on its STD; the physical decoder must compensate for ways in which its design differs from that of the STD.

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# Information technology — Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s —

## Part 1: Systems

### Section 1: General

#### 1.1 Scope

This part of ISO/IEC 11172 specifies the system layer of the coding. It was developed principally to support the combination of the video and audio coding methods defined in ISO/IEC 11172-2 and ISO/IEC 11172-3. The system layer supports five basic functions:

- a) the synchronization of multiple compressed streams on playback,
- b) the interleaving of multiple compressed streams into a single stream,
- c) the initialization of buffering for playback start up,
- d) continuous buffer management, and
- e) time identification.

An ISO/IEC 11172 multiplexed bit stream is constructed in two layers: the outermost layer is the system layer, and the innermost is the compression layer. The system layer provides the functions necessary for using one or more compressed data streams in a system. The video and audio parts of this specification define the compression coding layer for audio and video data. Coding of other types of data is not defined by the specification, but is supported by the system layer provided that the other types of data adhere to the constraints defined in clause 2.4.

#### 1.2 Normative references

The following International Standards contain provisions which, through reference in this text, constitute provisions of this part of ISO/IEC 11172. At the time of publication, the editions indicated were valid. All standards are subject to revision, and parties to agreements based on this part of ISO 11172 are encouraged to investigate the possibility of applying the most recent editions of the standards indicated below. Members of IEC and ISO maintain registers of currently valid International Standards.

ISO/IEC 11172-2:1993 *Information technology - Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s - Part 2: Video.*

ISO/IEC 11172-3:1993 *Information technology - Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s - Part 3 Audio.*

CCIR Recommendation 601-2 *Encoding parameters of digital television for studios.*

CCIR Report 624-4 *Characteristics of systems for monochrome and colour television.*

CCIR Recommendation 648 *Recording of audio signals.*

CCIR Report 955-2 *Sound broadcasting by satellite for portable and mobile receivers, including Annex IV Summary description of Advanced Digital System II.*

CCITT Recommendation J.17 *Pre-emphasis used on Sound-Programme Circuits.*

IEEE Draft Standard P1180/D2 1990 *Specification for the implementation of 8x 8 inverse discrete cosine transform".*

IEC publication 908:1987 *CD Digital Audio System.*

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## Section 2: Technical elements

### 2.1 Definitions

For the purposes of ISO/IEC 11172, the following definitions apply. If specific to a part, this is noted in square brackets.

**2.1.1 ac coefficient [video]:** Any DCT coefficient for which the frequency in one or both dimensions is non-zero.

**2.1.2 access unit [system]:** In the case of compressed audio an access unit is an audio access unit. In the case of compressed video an access unit is the coded representation of a picture.

**2.1.3 adaptive segmentation [audio]:** A subdivision of the digital representation of an audio signal in variable segments of time.

**2.1.4 adaptive bit allocation [audio]:** The assignment of bits to subbands in a time and frequency varying fashion according to a psychoacoustic model.

**2.1.5 adaptive noise allocation [audio]:** The assignment of coding noise to frequency bands in a time and frequency varying fashion according to a psychoacoustic model.

**2.1.6 alias [audio]:** Mirrored signal component resulting from sub-Nyquist sampling.

**2.1.7 analysis filterbank [audio]:** Filterbank in the encoder that transforms a broadband PCM audio signal into a set of subsampled subband samples.

**2.1.8 audio access unit [audio]:** For Layers I and II an audio access unit is defined as the smallest part of the encoded bitstream which can be decoded by itself, where decoded means "fully reconstructed sound". For Layer III an audio access unit is part of the bitstream that is decodable with the use of previously acquired main information.

**2.1.9 audio buffer [audio]:** A buffer in the system target decoder for storage of compressed audio data.

**2.1.10 audio sequence [audio]:** A non-interrupted series of audio frames in which the following parameters are not changed:

- ID
- Layer
- Sampling Frequency
- For Layer I and II: Bitrate index

**2.1.11 backward motion vector [video]:** A motion vector that is used for motion compensation from a reference picture at a later time in display order.

**2.1.12 Bark [audio]:** Unit of critical band rate. The Bark scale is a non-linear mapping of the frequency scale over the audio range closely corresponding with the frequency selectivity of the human ear across the band.

**2.1.13 bidirectionally predictive-coded picture; B-picture [video]:** A picture that is coded using motion compensated prediction from a past and/or future reference picture.

**2.1.14 bitrate:** The rate at which the compressed bitstream is delivered from the storage medium to the input of a decoder.

**2.1.15 block companding [audio]:** Normalizing of the digital representation of an audio signal within a certain time period.

**2.1.16 block [video]:** An 8-row by 8-column orthogonal block of pels.

**2.1.17 bound [audio]:** The lowest subband in which intensity stereo coding is used.

**2.1.18 byte aligned:** A bit in a coded bitstream is byte-aligned if its position is a multiple of 8-bits from the first bit in the stream.

**2.1.19 byte:** Sequence of 8-bits.

**2.1.20 channel:** A digital medium that stores or transports an ISO/IEC 11172 stream.

**2.1.21 channel [audio]:** The left and right channels of a stereo signal

**2.1.22 chrominance (component) [video]:** A matrix, block or single pel representing one of the two colour difference signals related to the primary colours in the manner defined in CCIR Rec 601. The symbols used for the colour difference signals are Cr and Cb.

**2.1.23 coded audio bitstream [audio]:** A coded representation of an audio signal as specified in ISO/IEC 11172-3.

**2.1.24 coded video bitstream [video]:** A coded representation of a series of one or more pictures as specified in ISO/IEC 11172-2.

**2.1.25 coded order [video]:** The order in which the pictures are stored and decoded. This order is not necessarily the same as the display order.

**2.1.26 coded representation:** A data element as represented in its encoded form.

**2.1.27 coding parameters [video]:** The set of user-definable parameters that characterize a coded video bitstream. Bitstreams are characterised by coding parameters. Decoders are characterised by the bitstreams that they are capable of decoding.

**2.1.28 component [video]:** A matrix, block or single pel from one of the three matrices (luminance and two chrominance) that make up a picture.

**2.1.29 compression:** Reduction in the number of bits used to represent an item of data.

**2.1.30 constant bitrate coded video [video]:** A compressed video bitstream with a constant average bitrate.

**2.1.31 constant bitrate:** Operation where the bitrate is constant from start to finish of the compressed bitstream.

**2.1.32 constrained parameters [video]:** The values of the set of coding parameters defined in 2.4.3.2 of ISO/IEC 11172-2.

**2.1.33 constrained system parameter stream (CSPS) [system]:** An ISO/IEC 11172 multiplexed stream for which the constraints defined in 2.4.6 of this part of ISO/IEC 11172 apply.

**2.1.34 CRC:** Cyclic redundancy code.

**2.1.35 critical band rate [audio]:** Psychoacoustic function of frequency. At a given audible frequency it is proportional to the number of critical bands below that frequency. The units of the critical band rate scale are Barks.

**2.1.36 critical band [audio]:** Psychoacoustic measure in the spectral domain which corresponds to the frequency selectivity of the human ear. This selectivity is expressed in Bark.

**2.1.37 data element:** An item of data as represented before encoding and after decoding.

**2.1.38 dc-coefficient [video]:** The DCT coefficient for which the frequency is zero in both dimensions.

- 2.1.39 dc-coded picture; D-picture [video]:** A picture that is coded using only information from itself. Of the DCT coefficients in the coded representation, only the dc-coefficients are present.
- 2.1.40 DCT coefficient:** The amplitude of a specific cosine basis function.
- 2.1.41 decoded stream:** The decoded reconstruction of a compressed bitstream.
- 2.1.42 decoder input buffer [video]:** The first-in first-out (FIFO) buffer specified in the video buffering verifier.
- 2.1.43 decoder input rate [video]:** The data rate specified in the video buffering verifier and encoded in the coded video bitstream.
- 2.1.44 decoder:** An embodiment of a decoding process.
- 2.1.45 decoding (process):** The process defined in ISO/IEC 11172 that reads an input coded bitstream and produces decoded pictures or audio samples.
- 2.1.46 decoding time-stamp; DTS [system]:** A field that may be present in a packet header that indicates the time that an access unit is decoded in the system target decoder.
- 2.1.47 de-emphasis [audio]:** Filtering applied to an audio signal after storage or transmission to undo a linear distortion due to emphasis.
- 2.1.48 dequantization [video]:** The process of rescaling the quantized DCT coefficients after their representation in the bitstream has been decoded and before they are presented to the inverse DCT.
- 2.1.49 digital storage media; DSM:** A digital storage or transmission device or system.
- 2.1.50 discrete cosine transform; DCT [video]:** Either the forward discrete cosine transform or the inverse discrete cosine transform. The DCT is an invertible, discrete orthogonal transformation. The inverse DCT is defined in annex A of ISO/IEC 11172-2.
- 2.1.51 display order [video]:** The order in which the decoded pictures should be displayed. Normally this is the same order in which they were presented at the input of the encoder.
- 2.1.52 dual channel mode [audio]:** A mode, where two audio channels with independent programme contents (e.g. bilingual) are encoded within one bitstream. The coding process is the same as for the stereo mode.
- 2.1.53 editing:** The process by which one or more compressed bitstreams are manipulated to produce a new compressed bitstream. Conforming edited bitstreams must meet the requirements defined in ISO/IEC 11172.
- 2.1.54 elementary stream [system]:** A generic term for one of the coded video, coded audio or other coded bitstreams.
- 2.1.55 emphasis [audio]:** Filtering applied to an audio signal before storage or transmission to improve the signal-to-noise ratio at high frequencies.
- 2.1.56 encoder:** An embodiment of an encoding process.
- 2.1.57 encoding (process):** A process, not specified in ISO/IEC 11172, that reads a stream of input pictures or audio samples and produces a valid coded bitstream as defined in ISO/IEC 11172.
- 2.1.58 entropy coding:** Variable length lossless coding of the digital representation of a signal to reduce redundancy.
- 2.1.59 fast forward playback [video]:** The process of displaying a sequence, or parts of a sequence, of pictures in display-order faster than real-time.

**2.1.60 FFT:** Fast Fourier Transformation. A fast algorithm for performing a discrete Fourier transform (an orthogonal transform).

**2.1.61 filterbank [audio]:** A set of band-pass filters covering the entire audio frequency range.

**2.1.62 fixed segmentation [audio]:** A subdivision of the digital representation of an audio signal into fixed segments of time.

**2.1.63 forbidden:** The term "forbidden" when used in the clauses defining the coded bitstream indicates that the value shall never be used. This is usually to avoid emulation of start codes.

**2.1.64 forced updating [video]:** The process by which macroblocks are intra-coded from time-to-time to ensure that mismatch errors between the inverse DCT processes in encoders and decoders cannot build up excessively.

**2.1.65 forward motion vector [video]:** A motion vector that is used for motion compensation from a reference picture at an earlier time in display order.

**2.1.66 frame [audio]:** A part of the audio signal that corresponds to audio PCM samples from an Audio Access Unit.

**2.1.67 free format [audio]:** Any bitrate other than the defined bitrates that is less than the maximum valid bitrate for each layer.

**2.1.68 future reference picture [video]:** The future reference picture is the reference picture that occurs at a later time than the current picture in display order.

**2.1.69 granules [Layer II] [audio]:** The set of 3 consecutive subband samples from all 32 subbands that are considered together before quantization. They correspond to 96 PCM samples.

**2.1.70 granules [Layer III] [audio]:** 576 frequency lines that carry their own side information.

**2.1.71 group of pictures [video]:** A series of one or more coded pictures intended to assist random access. The group of pictures is one of the layers in the coding syntax defined in ISO/IEC 11172-2.

**2.1.72 Hann window [audio]:** A time function applied sample-by-sample to a block of audio samples before Fourier transformation.

**2.1.73 Huffman coding:** A specific method for entropy coding.

**2.1.74 hybrid filterbank [audio]:** A serial combination of subband filterbank and MDCT.

**2.1.75 IMDCT [audio]:** Inverse Modified Discrete Cosine Transform.

**2.1.76 intensity stereo [audio]:** A method of exploiting stereo irrelevance or redundancy in stereophonic audio programmes based on retaining at high frequencies only the energy envelope of the right and left channels.

**2.1.77 interlace [video]:** The property of conventional television pictures where alternating lines of the picture represent different instances in time.

**2.1.78 intra coding [video]:** Coding of a macroblock or picture that uses information only from that macroblock or picture.

**2.1.79 intra-coded picture; I-picture [video]:** A picture coded using information only from itself.

**2.1.80 ISO/IEC 11172 (multiplexed) stream [system]:** A bitstream composed of zero or more elementary streams combined in the manner defined in this part of ISO/IEC 11172.

- 2.1.81 joint stereo coding [audio]:** Any method that exploits stereophonic irrelevance or stereophonic redundancy.
- 2.1.82 joint stereo mode [audio]:** A mode of the audio coding algorithm using joint stereo coding.
- 2.1.83 layer [audio]:** One of the levels in the coding hierarchy of the audio system defined in ISO/IEC 11172-3.
- 2.1.84 layer [video and systems]:** One of the levels in the data hierarchy of the video and system specifications defined in this part of ISO/IEC 11172 and ISO/IEC 11172-2.
- 2.1.85 luminance (component) [video]:** A matrix, block or single pel representing a monochrome representation of the signal and related to the primary colours in the manner defined in CCIR Rec 601. The symbol used for luminance is Y.
- 2.1.86 macroblock [video]:** The four 8 by 8 blocks of luminance data and the two corresponding 8 by 8 blocks of chrominance data coming from a 16 by 16 section of the luminance component of the picture. Macroblock is sometimes used to refer to the pel data and sometimes to the coded representation of the pel values and other data elements defined in the macroblock layer of the syntax defined in ISO/IEC 11172-2. The usage is clear from the context.
- 2.1.87 mapping [audio]:** Conversion of an audio signal from time to frequency domain by subband filtering and/or by MDCT.
- 2.1.88 masking [audio]:** A property of the human auditory system by which an audio signal cannot be perceived in the presence of another audio signal.
- 2.1.89 masking threshold [audio]:** A function in frequency and time below which an audio signal cannot be perceived by the human auditory system.
- 2.1.90 MDCT [audio]:** Modified Discrete Cosine Transform.
- 2.1.91 motion compensation [video]:** The use of motion vectors to improve the efficiency of the prediction of pel values. The prediction uses motion vectors to provide offsets into the past and/or future reference pictures containing previously decoded pel values that are used to form the prediction error signal.
- 2.1.92 motion estimation [video]:** The process of estimating motion vectors during the encoding process.
- 2.1.93 motion vector [video]:** A two-dimensional vector used for motion compensation that provides an offset from the coordinate position in the current picture to the coordinates in a reference picture.
- 2.1.94 MS stereo [audio]:** A method of exploiting stereo irrelevance or redundancy in stereophonic audio programmes based on coding the sum and difference signal instead of the left and right channels.
- 2.1.95 non-intra coding [video]:** Coding of a macroblock or picture that uses information both from itself and from macroblocks and pictures occurring at other times.
- 2.1.96 non-tonal component [audio]:** A noise-like component of an audio signal.
- 2.1.97 Nyquist sampling:** Sampling at or above twice the maximum bandwidth of a signal.
- 2.1.98 pack [system]:** A pack consists of a pack header followed by one or more packets. It is a layer in the system coding syntax described in this part of ISO/IEC 11172.
- 2.1.99 packet data [system]:** Contiguous bytes of data from an elementary stream present in a packet.
- 2.1.100 packet header [system]:** The data structure used to convey information about the elementary stream data contained in the packet data.

**2.1.101 packet [system]:** A packet consists of a header followed by a number of contiguous bytes from an elementary data stream. It is a layer in the system coding syntax described in this part of ISO/IEC 11172.

**2.1.102 padding [audio]:** A method to adjust the average length in time of an audio frame to the duration of the corresponding PCM samples, by conditionally adding a slot to the audio frame.

**2.1.103 past reference picture [video]:** The past reference picture is the reference picture that occurs at an earlier time than the current picture in display order.

**2.1.104 pel aspect ratio [video]:** The ratio of the nominal vertical height of pel on the display to its nominal horizontal width.

**2.1.105 pel [video]:** Picture element.

**2.1.106 picture period [video]:** The reciprocal of the picture rate.

**2.1.107 picture rate [video]:** The nominal rate at which pictures should be output from the decoding process.

**2.1.108 picture [video]:** Source, coded or reconstructed image data. A source or reconstructed picture consists of three rectangular matrices of 8-bit numbers representing the luminance and two chrominance signals. The Picture layer is one of the layers in the coding syntax defined in ISO/IEC 11172-2. Note that the term "picture" is always used in ISO/IEC 11172 in preference to the terms field or frame.

**2.1.109 polyphase filterbank [audio]:** A set of equal bandwidth filters with special phase interrelationships, allowing for an efficient implementation of the filterbank.

**2.1.110 prediction [video]:** The use of a predictor to provide an estimate of the pel value or data element currently being decoded.

**2.1.111 predictive-coded picture; P-picture [video]:** A picture that is coded using motion compensated prediction from the past reference picture.

**2.1.112 prediction error [video]:** The difference between the actual value of a pel or data element and its predictor.

**2.1.113 predictor [video]:** A linear combination of previously decoded pel values or data elements.

**2.1.114 presentation time-stamp; PTS [system]:** A field that may be present in a packet header that indicates the time that a presentation unit is presented in the system target decoder.

**2.1.115 presentation unit; PU [system]:** A decoded audio access unit or a decoded picture.

**2.1.116 psychoacoustic model [audio]:** A mathematical model of the masking behaviour of the human auditory system.

**2.1.117 quantization matrix [video]:** A set of sixty-four 8-bit values used by the dequantizer.

**2.1.118 quantized DCT coefficients [video]:** DCT coefficients before dequantization. A variable length coded representation of quantized DCT coefficients is stored as part of the compressed video bitstream.

**2.1.119 quantizer scalefactor [video]:** A data element represented in the bitstream and used by the decoding process to scale the dequantization.

**2.1.120 random access:** The process of beginning to read and decode the coded bitstream at an arbitrary point.

**2.1.121 reference picture [video]:** Reference pictures are the nearest adjacent I- or P-pictures to the current picture in display order.

**2.1.122 reorder buffer [video]:** A buffer in the system target decoder for storage of a reconstructed I-picture or a reconstructed P-picture.

**2.1.123 requantization [audio]:** Decoding of coded subband samples in order to recover the original quantized values.

**2.1.124 reserved:** The term "reserved" when used in the clauses defining the coded bitstream indicates that the value may be used in the future for ISO/IEC defined extensions.

**2.1.125 reverse playback [video]:** The process of displaying the picture sequence in the reverse of display order.

**2.1.126 scalefactor band [audio]:** A set of frequency lines in Layer III which are scaled by one scalefactor.

**2.1.127 scalefactor index [audio]:** A numerical code for a scalefactor.

**2.1.128 scalefactor [audio]:** Factor by which a set of values is scaled before quantization.

**2.1.129 sequence header [video]:** A block of data in the coded bitstream containing the coded representation of a number of data elements.

**2.1.130 side information:** Information in the bitstream necessary for controlling the decoder.

**2.1.131 skipped macroblock [video]:** A macroblock for which no data are stored.

**2.1.132 slice [video]:** A series of macroblocks. It is one of the layers of the coding syntax defined in ISO/IEC 11172-2.

**2.1.133 slot [audio]:** A slot is an elementary part in the bitstream. In Layer I a slot equals four bytes, in Layers II and III one byte.

**2.1.134 source stream:** A single non-multiplexed stream of samples before compression coding.

**2.1.135 spreading function [audio]:** A function that describes the frequency spread of masking.

**2.1.136 start codes [system and video]:** 32-bit codes embedded in that coded bitstream that are unique. They are used for several purposes including identifying some of the layers in the coding syntax.

**2.1.137 STD input buffer [system]:** A first-in first-out buffer at the input of the system target decoder for storage of compressed data from elementary streams before decoding.

**2.1.138 stereo mode [audio]:** Mode, where two audio channels which form a stereo pair (left and right) are encoded within one bitstream. The coding process is the same as for the dual channel mode.

**2.1.139 stuffing (bits); stuffing (bytes) :** Code-words that may be inserted into the compressed bitstream that are discarded in the decoding process. Their purpose is to increase the bitrate of the stream.

**2.1.140 subband [audio]:** Subdivision of the audio frequency band.

**2.1.141 subband filterbank [audio]:** A set of band filters covering the entire audio frequency range. In ISO/IEC 11172-3 the subband filterbank is a polyphase filterbank.

**2.1.142 subband samples [audio]:** The subband filterbank within the audio encoder creates a filtered and subsampled representation of the input audio stream. The filtered samples are called subband samples. From 384 time-consecutive input audio samples, 12 time-consecutive subband samples are generated within each of the 32 subbands.

**2.1.143 syncword [audio]:** A 12-bit code embedded in the audio bitstream that identifies the start of a frame.

**2.1.144 synthesis filterbank [audio]:** Filterbank in the decoder that reconstructs a PCM audio signal from subband samples.

**2.1.145 system header [system]:** The system header is a data structure defined in this part of ISO/IEC 11172 that carries information summarising the system characteristics of the ISO/IEC 11172 multiplexed stream.

**2.1.146 system target decoder; STD [system]:** A hypothetical reference model of a decoding process used to describe the semantics of an ISO/IEC 11172 multiplexed bitstream.

**2.1.147 time-stamp [system]:** A term that indicates the time of an event.

**2.1.148 triplet [audio]:** A set of 3 consecutive subband samples from one subband. A triplet from each of the 32 subbands forms a granule.

**2.1.149 tonal component [audio]:** A sinusoid-like component of an audio signal.

**2.1.150 variable bitrate:** Operation where the bitrate varies with time during the decoding of a compressed bitstream.

**2.1.151 variable length coding; VLC:** A reversible procedure for coding that assigns shorter code-words to frequent events and longer code-words to less frequent events.

**2.1.152 video buffering verifier; VBV [video]:** A hypothetical decoder that is conceptually connected to the output of the encoder. Its purpose is to provide a constraint on the variability of the data rate that an encoder or editing process may produce.

**2.1.153 video sequence [video]:** A series of one or more groups of pictures. It is one of the layers of the coding syntax defined in ISO/IEC 11172-2.

**2.1.154 zig-zag scanning order [video]:** A specific sequential ordering of the DCT coefficients from (approximately) the lowest spatial frequency to the highest.

## 2.2 Symbols and abbreviations

The mathematical operators used to describe this International Standard are similar to those used in the C programming language. However, integer division with truncation and rounding are specifically defined. The bitwise operators are defined assuming twos-complement representation of integers. Numbering and counting loops generally begin from zero.

### 2.2.1 Arithmetic operators

+	Addition.
-	Subtraction (as a binary operator) or negation (as a unary operator).
++	Increment.
--	Decrement.
*	Multiplication.
^	Power.
/	Integer division with truncation of the result toward zero. For example, 7/4 and -7/-4 are truncated to 1 and -7/4 and 7/-4 are truncated to -1.
//	Integer division with rounding to the nearest integer. Half-integer values are rounded away from zero unless otherwise specified. For example 3//2 is rounded to 2, and -3//2 is rounded to -2.
DIV	Integer division with truncation of the result towards-∞.
	Absolute value. $ x  = x$ when $x > 0$ $ x  = 0$ when $x == 0$ $ x  = -x$ when $x < 0$
%	Modulus operator. Defined only for positive numbers.
Sign( )	Sign(x) $= 1$ $x > 0$ $0$ $x == 0$ $-1$ $x < 0$
NINT ( )	Nearest integer operator. Returns the nearest integer value to the real-valued argument. Half-integer values are rounded away from zero.
sin	Sine.
cos	Cosine.
exp	Exponential.
√	Square root.
log <sub>10</sub>	Logarithm to base ten.
log <sub>e</sub>	Logarithm to base e.
log <sub>2</sub>	Logarithm to base 2.

### 2.2.2 Logical operators

	Logical OR.
&&	Logical AND.

! Logical NOT.

### 2.2.3 Relational operators

> Greater than.

>= Greater than or equal to.

< Less than.

<= Less than or equal to.

= Equal to.

!= Not equal to.

max [...,] the maximum value in the argument list.

min [...,] the minimum value in the argument list.

### 2.2.4 Bitwise operators

A twos complement number representation is assumed where the bitwise operators are used.

& AND.

| OR.

>> Shift right with sign extension.

<< Shift left with zero fill.

### 2.2.5 Assignment

= Assignment operator.

### 2.2.6 Mnemonics

The following mnemonics are defined to describe the different data types used in the coded bit-stream.

bslbf	Bit string, left bit first, where "left" is the order in which bit strings are written in ISO/IEC 11172. Bit strings are written as a string of 1s and 0s within single quote marks, e.g. '1000 0001'. Blanks within a bit string are for ease of reading and have no significance.
ch	Channel. If ch has the value 0, the left channel of a stereo signal or the first of two independent signals is indicated. (Audio)
nch	Number of channels; equal to 1 for single_channel mode, 2 in other modes. (Audio)
gr	Granule of 3 * 32 subband samples in audio Layer II, 18 * 32 sub-band samples in audio Layer III. (Audio)
main_data	The main_data portion of the bitstream contains the scalefactors, Huffman encoded data, and ancillary information. (Audio)
main_data_beg	The location in the bitstream of the beginning of the main_data for the frame. The location is equal to the ending location of the previous frame's main_data plus one bit. It is calculated from the main_data_end value of the previous frame. (Audio)
part2_length	The number of main_data bits used for scalefactors. (Audio)

rpchof	Remainder polynomial coefficients, highest order first. (Audio)
sb	Subband. (Audio)
sblimit	The number of the lowest sub-band for which no bits are allocated. (Audio)
scfsi	Scalefactor selection information. (Audio)
switch_point_l	Number of scalefactor band (long block scalefactor band) from which point on window switching is used. (Audio)
switch_point_s	Number of scalefactor band (short block scalefactor band) from which point on window switching is used. (Audio)
uimsbf	Unsigned integer, most significant bit first.
vlclbf	Variable length code, left bit first, where "left" refers to the order in which the VLC codes are written.
window	Number of the actual time slot in case of block_type=2, $0 \leq \text{window} \leq 2$ . (Audio)

The byte order of multi-byte words is most significant byte first.

### 2.2.7 Constants

$\pi$	3,14159265358...
e	2,71828182845...

## 2.3 Method of describing bit stream syntax

The bit stream retrieved by the decoder is described in 2.4.3. Each data item in the bit stream is in bold type. It is described by its name, its length in bits, and a mnemonic for its type and order of transmission.

The action caused by a decoded data element in a bit stream depends on the value of that data element and on data elements previously decoded. The decoding of the data elements and definition of the state variables used in their decoding are described in 2.4.4. The following constructs are used to express the conditions when data elements are present, and are in normal type:

Note this syntax uses the 'C'-code convention that a variable or expression evaluating to a non-zero value is equivalent to a condition that is true.

while ( condition ) { <b>data_element</b> . . . }	If the condition is true, then the group of data elements occurs next in the data stream. This repeats until the condition is not true.
do { <b>data_element</b> . . . } while ( condition )	The data element always occurs at least once. The data element is repeated until the condition is not true.
if ( condition ) { <b>data_element</b> . . . }	If the condition is true, then the first group of data elements occurs next in the data stream.
else { <b>data_element</b> . . . }	If the condition is not true, then the second group of data elements occurs next in the data stream.

for (expr1; expr2; expr3) { expr1 is an expression specifying the initialization of the loop. Normally it specifies the initial state of the counter. expr2 is a condition specifying a test made before each iteration of the loop. The loop terminates when the condition is not true. expr3 is an expression that is performed at the end of each iteration of the loop, normally it increments a counter.

```

data_element
  . . .
}

```

Note that the most common usage of this construct is as follows:

```

for ( i = 0; i < n; i++) { The group of data elements occurs n times. Conditional constructs
  data_element          within the group of data elements may depend on the value of the
  . . .                  loop control variable i, which is set to zero for the first occurrence,
}                          incremented to one for the second occurrence, and so forth.

```

As noted, the group of data elements may contain nested conditional constructs. For compactness, the {} may be omitted when only one data element follows.

**data\_element []** data\_element [] is an array of data. The number of data elements is indicated by the context.

**data\_element [n]** data\_element [n] is the n+1th element of an array of data.

**data\_element [m][n]** data\_element [m][n] is the m+1,n+1 th element of a two-dimensional array of data.

**data\_element [l][m][n]** data\_element [l][m][n] is the l+1,m+1,n+1 th element of a three-dimensional array of data.

**data\_element [m..n]** is the inclusive range of bits between bit m and bit n in the data\_element.

While the syntax is expressed in procedural terms, it should not be assumed that 2.4.3 implements a satisfactory decoding procedure. In particular, it defines a correct and error-free input bitstream. Actual decoders must include a means to look for start codes in order to begin decoding correctly, and to identify errors, erasures or insertions while decoding. The methods to identify these situations, and the actions to be taken, are not standardized.

### Definition of bytealigned function

The function bytealigned () returns 1 if the current position is on a byte boundary, that is the next bit in the bit stream is the first bit in a byte. Otherwise it returns 0.

### Definition of nextbits function

The function nextbits () permits comparison of a bit string with the next bits to be decoded in the bit stream.

### Definition of next\_start\_code function

The next\_start\_code function removes any zero bit and zero byte stuffing and locates the next start code.

Syntax	No. of bits	Mnemonic
next_start_code() {		
while ( !bytealigned() )		
<b>zero_bit</b>	1	"0"
while ( nextbits() != '0000 0000 0000 0000 0000 0001' )		
<b>zero_byte</b>	8	"00000000"
}		

This function checks whether the current position is byte aligned. If it is not, zero stuffing bits are present. After that any number of zero bytes may be present before the start-code. Therefore start-codes are always byte aligned and may be preceded by any number of zero stuffing bits.

## 2.4 Requirements

### 2.4.1 Coding structure and parameters

The system coding layer allows one or more elementary streams to be combined into a single stream. Data from each elementary stream are multiplexed and encoded together with information that allows elementary streams to be replayed in synchronism.

#### ISO/IEC 11172 multiplexed stream

An ISO/IEC 11172 stream consists of one or more elementary streams multiplexed together. Each elementary stream consists of access units, which are the coded representation of presentation units. The presentation unit for a video elementary stream is a picture. The corresponding access unit includes all the coded data for the picture. The access unit containing the first coded picture of a group of pictures also includes any preceding data from that group of pictures, as defined in 2.4.2.4 in ISO/IEC 11172-2, starting with the `group_start_code`. The access unit containing the first coded picture after a sequence header, as defined in 2.4.2.3 in part 2, also includes that sequence header. The `sequence_end_code` is included in the access unit containing the last coded picture of a sequence. (See 2.4.2.2 in ISO/IEC 11172-2 for the definition of the `sequence_end_code`). The presentation unit for an audio elementary stream is the set of samples that corresponds to samples from an audio frame (see 2.4.3.1, 2.4.2.1, and 2.4.2.2 in ISO/IEC 11172-3 for the definition of an audio frame).

Data from elementary streams is stored in packets. A packet consists of a packet header followed by packet data. The packet header begins with a 32-bit start-code that also identifies the stream to which the packet data belongs. The packet header may contain decoding and/or presentation time-stamps (DTS and PTS) that refer to the first access unit that commences in the packet. The packet data contains a variable number of contiguous bytes from one elementary stream.

Packets are organised in packs. A pack commences with a pack header and is followed by zero or more packets. The pack header begins with a 32-bit start-code. The pack header is used to store timing and bitrate information.

The stream begins with a system header that optionally may be repeated. The system header carries a summary of the system parameters defined in the stream.

### 2.4.2 System target decoder

The semantics of the multiplexed stream specified in 2.4.4 and the constraints on these semantics specified in 2.4.5 require exact definitions of decoding events and the times at which these events occur. The definitions needed are set out in this International Standard using a hypothetical decoder known as the system target decoder (STD).

The STD is a conceptual model used to define these terms precisely and to model the decoding process during the construction of ISO/IEC 11172 streams. The STD is defined only for this purpose. Neither the architecture of the STD nor the timing described precludes uninterrupted, synchronized play-back of ISO/IEC 11172 multiplexed streams from a variety of decoders with different architectures or timing schedules.

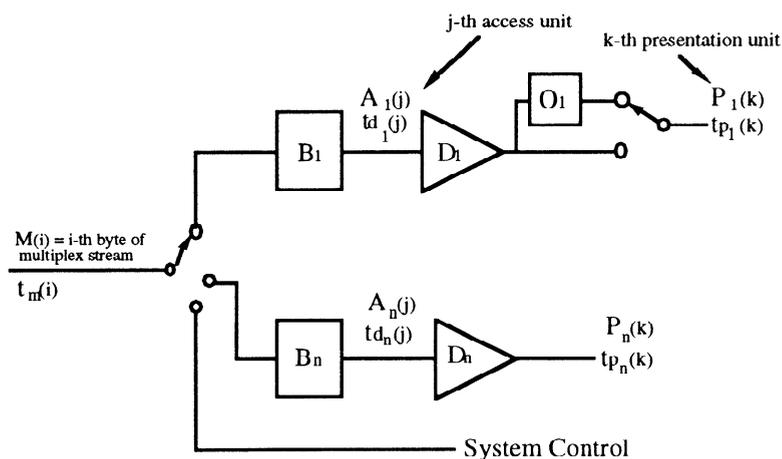


Figure 2 -- Diagram of system target decoder

### Notation

The following notation is used to describe the system target decoder and is partially illustrated in figure 2.

$i, i'$  are indices to bytes in the ISO/IEC 11172 multiplexed stream. The first byte has index 0.

$j$  is an index to access units in the elementary streams.

$k, k', k''$  are indices to presentation units in the elementary streams.

$n$  is an index to the elementary streams.

$M(i)$  is the  $i^{\text{th}}$  byte in the ISO/IEC 11172 multiplexed stream.

$t_m(i)$  indicates the time in seconds at which the  $i^{\text{th}}$  byte of the ISO/IEC 11172 multiplexed stream enters the system target decoder. The value  $t_m(0)$  is an arbitrary constant.

$SCR(i)$  is the time encoded in the SCR field measured in units of the 90 kHz system clock where  $i$  is the byte index of the final byte of the SCR field.

$A_n(j)$  is the  $j^{\text{th}}$  access unit in elementary stream  $n$ . Note that access units are indexed in decoding order.

$t_{d_n}(j)$  is the decoding time, measured in seconds, in the system target decoder of the  $j^{\text{th}}$  access unit in elementary stream  $n$ .

$P_n(k)$  is the  $k^{\text{th}}$  presentation unit in elementary stream  $n$ .

$t_{p_n}(k)$  is the presentation time, measured in seconds, in the system target decoder of the  $k^{\text{th}}$  presentation unit in elementary stream  $n$ .

$t$  is time measured in seconds.

$F_n(t)$  is the fullness, measured in bytes, of the system target decoder input buffer for elementary stream  $n$  at time  $t$ .

$B_n$  is the input buffer in the system target decoder for elementary stream  $n$ .

$BS_n$  is the size of the system target decoder input buffer, measured in bytes, for elementary stream  $n$ .

$D_n$  is the decoder for elementary stream  $n$ .

$O_n$  is the reorder buffer for elementary stream  $n$ .

## System Clock Frequency

Timing information is carried by several data fields defined in 2.4.3 and 2.4.4. This information is coded as the sampled value of a system clock.

The value of the system clock frequency is measured in hertz and shall meet the following constraints:

$$90\,000 - 4,5 \leq \text{system\_clock\_frequency} \leq 90\,000 + 4,5$$

$$\text{rate of change of system\_clock\_frequency with time} \leq 250 * 10^{-6} \text{ Hz/s}$$

The notation "system\_clock\_frequency" is used in several places in this part of ISO/IEC 11172 to refer to the frequency of a clock meeting these requirements. For notational convenience, equations in which SCR, PTS, or DTS appear lead to values of time which are accurate to some integral multiple of  $(2^{33}/\text{system\_clock\_frequency})$ . This is due to the 33-bit encoding of timing information.

## Input to the System Target Decoder

Data from the ISO/IEC 11172 multiplexed stream enters the system target decoder. The  $i^{\text{th}}$  byte,  $M(i)$ , enters at time  $tm(i)$ . The time at which this byte enters the system target decoder can be recovered from the input stream by decoding the input system clock reference (SCR) fields encoded in the pack header. The value encoded in the  $SCR(i')$  field indicates time  $tm(i')$ , where  $i'$  refers to the last byte of the SCR field,  $M(i')$ .

Specifically:

$$SCR(i') = NINT ( \text{system\_clock\_frequency} * tm(i') ) \% 2^{33}$$

The input arrival time,  $tm(i)$ , for all other bytes shall be constructed from  $SCR(i')$  and the rate at which data arrive, where the arrival rate within each pack is the value represented in the  $mux\_rate$  field in that pack's header (see 2.4.3.2 and 2.4.4.2).

$$tm(i) = \frac{SCR(i')}{\text{system\_clock\_frequency}} + \frac{i - i'}{(mux\_rate * 50)}$$

Where:

- $i'$  is the index of the final byte of the system\_clock\_reference field in the pack header.
- $i$  is the index of any byte in the pack, including the pack header.
- $SCR(i')$  is the time encoded in the system\_clock\_reference field in units of the system clock.
- $mux\_rate$  is a field defined in 2.4.3.2 and 2.4.4.2.

After delivery of the last byte of a pack there may be a time interval during which no bytes are delivered to the input of the system target decoder.

Variable rate operation of the system target decoder is provided through the use of the  $mux\_rate$  field, the value of which may vary from pack to pack, and the fact that the data rate entering the system target decoder may drop to zero after the last byte of one pack arrives and before the following pack header arrives.

## Buffering

The packet data from elementary stream  $n$  is passed to the input buffer for stream  $n$ ,  $B_n$ . Transfer of byte  $M(i)$  from the system target decoder input to  $B_n$  is instantaneous, so that byte  $M(i)$  enters the buffer for stream  $n$ , of size  $BS_n$ , at time  $tm(i)$ .

Bytes present in the pack, system or packet headers of ISO/IEC 11172 multiplexed stream but not part of the packet data (for example the SCR, DTS, PTS, packet\_length fields, etc.- see 2.4.3) are not delivered to any of the buffers, but may be used to control the system.

The input buffer sizes  $BS_1$  through  $BS_n$  are given by parameters in the syntax (see 2.4.3 and 2.4.4).

At the decoding time,  $td_n(j)$ , all the data for the access unit that has been in the input buffer longest ( $A_n(j)$ ) is removed instantaneously. In the case of a video elementary stream, group of picture and sequence header data that precede the picture are removed at the same time. In the case of the first coded picture of a video sequence, any zero bit or byte stuffing immediately preceding the sequence header is removed at the same time. Note that this only applies to the first picture of a video sequence and not to additional occurrences of a sequence header within a video sequence. As the access unit is removed from the buffer it is instantaneously decoded into a presentation unit.

## Decoding

Elementary streams buffered in  $B_1$  through  $B_n$  are decoded instantaneously by decoders  $D_1$  through  $D_n$  and may be delayed in reorder buffers  $O_1$  through  $O_n$  before being presented to the viewer at the output of the system target decoder. Reorder buffers are used only in video decoding to store I-pictures and P-pictures while the sequence of presentation units is reordered before presentation.

In the case of a video elementary stream, some access units may not be stored in presentation order. These access units will need to be reordered before presentation. In particular, an I-picture or a P-picture stored before one or more B-pictures must be delayed in the reorder buffer,  $O_n$ , of the system target decoder before being presented. It should be delayed until the next I-picture or P-picture is decoded. While it is stored in the reorder buffer, the subsequent B-pictures are decoded and presented.

If  $P_n(k)$  is an I-picture or a P-picture that needs to be reordered before presentation, it is stored in  $O_n$  after being decoded and the picture previously stored in  $O_n$  is presented. Subsequent B-pictures are decoded and presented without reordering.

The time at which a presentation unit  $P_n(k)$  is presented to the viewer is  $tp_n(k)$ . For presentation units that are not reordered,  $tp_n(k)$  is equal to  $td_n(j)$  since the access units are decoded instantaneously. For presentation units that are reordered  $tp_n(k)$  and  $td_n(j)$  differ by the time that  $P_n(k)$  is delayed in the reorder buffer, which is a multiple of the nominal picture period.

Subclause 2.4.1 of ISO/IEC 11172-2 explains reordering of video pictures in greater detail.

## Presentation

The function of a decoding system is to reconstruct presentation units from compressed data and to present them in a synchronized sequence at the correct presentation times. Although real audio and visual presentation devices generally have finite and different delays and may have additional delays imposed by post-processing or output functions, the system target decoder models these delays as zero.

In the system target decoder the display of a video presentation unit (a picture) occurs instantaneously at its presentation time,  $tp_n(k)$ .

In the system target decoder the output of an audio presentation unit starts at its presentation time,  $tp_n(k)$ , when the decoder instantaneously presents the first sample. Subsequent samples in the presentation unit are presented in sequence at the audio sampling rate.

### 2.4.3 Specification of the system stream syntax

The following syntax describes a stream of bytes.

#### 2.4.3.1 ISO/IEC 11172 Layer

Syntax	No. of bits	Mnemonic
<pre>iso11172_stream() {   do {     pack()   } while (nextbits() == pack_start_code)   iso_11172_end_code }</pre>	32	bslbf

#### 2.4.3.2 Pack Layer

##### Pack

Syntax	No. of bits	Mnemonic
<pre>pack() {   pack_start_code   '0010'   system_clock_reference [32..30]   marker_bit   system_clock_reference [29..15]   marker_bit   system_clock_reference [14..0]   marker_bit   marker_bit   mux_rate   marker_bit    if (nextbits() == system_header_start_code)     system_header ()    while (nextbits() == packet_start_code_prefix)     packet() }</pre>	 32 4 3 1 15 1 15 1 1 1 22 1	 bslbf bslbf bslbf bslbf bslbf bslbf bslbf bslbf uimsbf bslbf

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**System header**

Syntax	No. of bits	Mnemonic
system_header () {		
system_header_start_code	32	bslbf
header_length	16	uimsbf
marker_bit	1	bslbf
rate_bound	22	uimsbf
marker_bit	1	bslbf
audio_bound	6	uimsbf
fixed_flag	1	bslbf
CSPS_flag	1	bslbf
system_audio_lock_flag	1	bslbf
system_video_lock_flag	1	bslbf
marker_bit	1	bslbf
video_bound	5	uimsbf
reserved_byte	8	bslbf
while (nextbits () == '1') {		
stream_id	8	uimsbf
'11'	2	bslbf
STD_buffer_bound_scale	1	bslbf
STD_buffer_size_bound	13	uimsbf
}		
}		

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## 2.4.3.3 Packet Layer

Syntax	No. of bits	Mnemonic
packet() {		
<b>packet_start_code_prefix</b>	<b>24</b>	<b>bslbf</b>
<b>stream_id</b>	<b>8</b>	<b>uimsbf</b>
<b>packet_length</b>	<b>16</b>	<b>uimsbf</b>
if (stream_id != private_stream_2) {		
while (nextbits() == '1111 1111') {		
<b>stuffing_byte</b>	<b>8</b>	<b>bslbf</b>
if (nextbits() == '01') {		
'01'	<b>2</b>	<b>bslbf</b>
<b>STD_buffer_scale</b>	<b>1</b>	<b>bslbf</b>
<b>STD_buffer_size</b>	<b>13</b>	<b>uimsbf</b>
}		
if (nextbits() == '0010') {		
'0010'	<b>4</b>	<b>bslbf</b>
<b>presentation_time_stamp[32..30]</b>	<b>3</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
<b>presentation_time_stamp[29..15]</b>	<b>15</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
<b>presentation_time_stamp[14..0]</b>	<b>15</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
}		
else if (nextbits() == '0011') {		
'0011'	<b>4</b>	<b>bslbf</b>
<b>presentation_time_stamp[32..30]</b>	<b>3</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
<b>presentation_time_stamp[29..15]</b>	<b>15</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
<b>presentation_time_stamp[14..0]</b>	<b>15</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
'0001'	<b>4</b>	<b>bslbf</b>
<b>decoding_time_stamp[32..30]</b>	<b>3</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
<b>decoding_time_stamp[29..15]</b>	<b>15</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
<b>decoding_time_stamp[14..0]</b>	<b>15</b>	<b>bslbf</b>
<b>marker_bit</b>	<b>1</b>	<b>bslbf</b>
}		
else		
'0000 1111'	<b>8</b>	<b>bslbf</b>
}		
for (i = 0; i < N; i++) {		
<b>packet_data_byte</b>	<b>8</b>	<b>bslbf</b>
}		

## 2.4.4 Semantic definition of fields in syntax

### 2.4.4.1 ISO/IEC 11172 Layer

**iso\_11172\_end\_code** -- The iso\_11172\_end\_code is the bit string "0000 0000 0000 0000 0000 0001 1011 1001" (000001B9 in hexadecimal). It terminates the ISO/IEC 11172 multiplexed stream.

### 2.4.4.2 Pack Layer

#### Pack

**pack\_start\_code** -- The pack\_start\_code is the bit string "0000 0000 0000 0000 0000 0001 1011 1010" (000001BA in hexadecimal). It identifies the beginning of a pack.

**system\_clock\_reference** -- The system\_clock\_reference (SCR) is a 33-bit number coded in three separate fields. It indicates the intended time of arrival of the last byte of the system\_clock\_reference field at the input of the system target decoder. The value of the SCR is measured in the number of periods of a 90kHz system clock with a tolerance specified in 2.4.2. Using the notation of 2.4.2, the value encoded in the system\_clock\_reference is:

$$\text{SCR}(i) = \text{NINT}(\text{system\_clock\_frequency} * (\text{tm}(i)) \% 2^{33})$$

for  $i$  such that  $M(i)$  is the last byte of the coded system\_clock\_reference field.

**marker\_bit** -- A marker\_bit is a one bit field that has the value "1".

**mux\_rate** -- This is a positive integer specifying the rate at which the system target decoder receives the ISO/IEC 11172 multiplexed stream during the pack in which it is included. The value of mux\_rate is measured in units of 50 bytes/s, rounded upwards. The value zero is forbidden. The value represented in mux\_rate is used to define the time of arrival of bytes at the input to the system target decoder in 2.4.2. The value encoded in the mux\_rate field may vary from pack to pack in an ISO/IEC 11172 multiplexed stream.

#### System Header

**system\_header\_start\_code** -- The system\_header\_start\_code is the bit string "0000 0000 0000 0000 0000 0001 1011 1011" (000001BB in hexadecimal). It identifies the beginning of a system header.

**header\_length** -- The header\_length shall be equal to the number of bytes in the system header following the header\_length field. Note that future extensions of this part of ISO/IEC 11172 may extend the system header.

**rate\_bound** -- The rate\_bound is an integer value greater than or equal to the maximum value of the mux\_rate field coded in any pack of the ISO/IEC 11172 multiplexed stream. It may be used by a decoder to assess whether it is capable of decoding the entire stream.

**audio\_bound** -- The audio\_bound is an integer, in the inclusive range from 0 to 32, greater than or equal to the maximum number of ISO/IEC 11172 audio streams in the ISO/IEC 11172 multiplexed stream for which the decoding processes are simultaneously active. For the purpose of this clause, the decoding process of an MPEG audio stream is active, if the STD buffer is not empty, or if the decoded access unit is being presented in the STD model.

**fixed\_flag** -- The fixed\_flag is a one-bit flag. If its value is set to "1" fixed bitrate operation is indicated. If its value is set to "0" variable bitrate operation is indicated. During fixed bitrate operation, the value encoded in all system\_clock\_reference fields in the multiplexed ISO/IEC 11172 stream shall adhere to the following linear equation:

$$\text{SCR}(i) = \text{NINT}(c1 * i + c2) \% 2^{33}$$

where

- c1 is a real-valued constant valid for all  $i$ ;
- c2 is a real-valued constant valid for all  $i$ ;
- $i$  is the index in the ISO/IEC 11172 multiplexed stream of the final byte of any system\_clock\_reference field in the stream.

**CSPS\_flag** -- The CSPS\_flag is a one-bit flag. If its value is set to "1" the ISO/IEC 11172 multiplexed stream meets the constraints defined in 2.4.6.

**system\_audio\_lock\_flag** -- The system\_audio\_lock\_flag is a one-bit flag indicating that there is a specified, constant rational relationship between the audio sampling rate and the system clock frequency in the system target decoder. Subclause 2.4.2 defines system\_clock\_frequency and the audio sampling rate is specified in ISO/IEC 11172-3. The system\_audio\_lock\_flag may only be set to "1" if, for all presentation units in all audio elementary streams in the ISO/IEC 11172 multiplexed stream, the ratio of system\_clock\_frequency to the actual audio sampling rate, SCASR, is constant and equal to the value indicated in the following table at the nominal sampling rate indicated in the audio stream.

$$\text{SCASR} = \frac{\text{system\_clock\_frequency}}{\text{audio sample rate in the STD}}$$

The notation  $\frac{X}{Y}$  denotes real division.

Nominal audio sampling frequency (kHz)	32	44,1	48
Ratio SCASR	$\frac{90\ 000}{32\ 000}$	$\frac{90\ 000}{44\ 100}$	$\frac{90\ 000}{48\ 000}$

**system\_video\_lock\_flag** -- The system\_video\_lock\_flag is a one-bit flag indicating that there is a specified, constant rational relationship between the video picture rate and the system clock frequency in the system target decoder. Subclause 2.4.2 defines system\_clock\_frequency and the video picture rate is specified in ISO/IEC 11172-2. The system\_video\_lock\_flag may only be set to "1" if, for all presentation units in all video elementary streams in the ISO/IEC 11172 multiplexed stream, the ratio of system\_clock\_frequency to the actual video picture rate, SCPR, is constant and equal to the value indicated in the following table at the nominal picture rate indicated in the video stream.

$$\text{SCPR} = \frac{\text{system\_clock\_frequency}}{\text{picture rate in the STD}}$$

Nominal picture rate (Hz)	23,976	24	25	29,97	30	50	59,94	60
Ratio SCPR	$\frac{15\ 015}{4}$	3 750	3 600	3 003	3 000	1 800	$\frac{3\ 003}{2}$	1 500

The values of the ratio SCPR are exact. The actual picture rate differs slightly from the nominal rate in cases where the nominal rate is 23,976, 29,97, or 59,94 pictures per second.

**video\_bound** -- The video\_bound is an integer, in the inclusive range from 0 to 16, greater than or equal to the maximum number of ISO/IEC 11172 video streams in the ISO/IEC 11172 multiplexed stream for which the decoding processes are simultaneously active. For the purpose of this clause, the decoding process of an ISO/IEC 11172 video streams is active if the STD buffer is not empty, or if the decoded access unit is being presented in the STD model, or if the reorder buffer is not empty.

**reserved\_byte** -- This byte is reserved for future use by ISO/IEC. Until otherwise specified by ISO/IEC it shall have the value "1111 1111".

**stream\_id** -- The stream\_id indicates the type and number of the stream to which the following STD\_buffer\_bound\_scale and STD\_buffer\_size\_bound fields refer.

If stream\_id equals "1011 1000" the STD\_buffer\_bound\_scale and STD\_buffer\_size\_bound fields following the stream\_id refer to all audio streams in the ISO/IEC 11172 multiplexed stream.

If stream\_id equals "1011 1001" the STD\_buffer\_bound\_scale and STD\_buffer\_size\_bound fields following the stream\_id refer to all video streams in the ISO/IEC 11172 multiplexed stream.

If the stream\_id takes on any other value it shall be a byte value greater than or equal to "1011 1100" and shall be interpreted as referring to the stream type and number according to table 1. This table is used also to identify the stream type and number indicated by the stream\_id defined in 2.4.4.3.

Table 1 -- stream\_id table

stream_id	stream type
1011 1100	reserved stream
1011 1101	private_stream_1
1011 1110	padding stream
1011 1111	private_stream_2
110x xxxx	ISO/IEC 11172-3 audio stream - number xxxxx
1110 xxxx	ISO/IEC 11172-2 video stream - number xxxxx
1111 xxxx	reserved data stream - number xxxxx
The notation x means that the values 0 and 1 are both permitted and result in the same stream type. The stream number is given by the values taken by the x's.	

Each elementary stream present in the ISO/IEC 11172 multiplexed stream shall have its STD\_buffer\_bound\_scale and STD\_buffer\_size\_bound specified exactly once by this mechanism in each system header.

**STD\_buffer\_bound\_scale** -- The STD\_buffer\_bound\_scale is a one-bit field that indicates the scaling factor used to interpret the subsequent STD\_buffer\_size\_bound field. If the preceding stream\_id indicates an audio stream, STD\_buffer\_bound\_scale shall have the value "0". If the preceding stream\_id indicates a video stream, STD\_buffer\_bound\_scale shall have the value "1". For all other stream types, the value of the STD\_buffer\_bound\_scale may be either "1" or "0".

**STD\_buffer\_size\_bound** -- The STD\_buffer\_size\_bound is a 13-bit unsigned integer defining a value greater than or equal to the maximum system target decoder input buffer size,  $BS_n$ , over all packets for stream  $n$  in the ISO/IEC 11172 stream. If STD\_buffer\_bound\_scale has the value "0" then STD\_buffer\_size\_bound measures the buffer size bound in units of 128 bytes. If STD\_buffer\_bound\_scale has the value "1" then STD\_buffer\_size\_bound measures the buffer size bound in units of 1024 bytes. Thus:

```

if (STD_buffer_bound_scale == 0)
     $BS_n \leq STD\_buffer\_size\_bound * 128;$ 
else
     $BS_n \leq STD\_buffer\_size\_bound * 1024;$ 

```

#### 2.4.4.3 Packet Layer

**packet\_start\_code\_prefix** -- The packet\_start\_code\_prefix is a 24-bit code. Together with the stream\_id that follows, it constitutes a packet start code that identifies the beginning of a packet. The packet\_start\_code\_prefix is the bit string "0000 0000 0000 0000 0000 0001" (000001 in hexadecimal).

**stream\_id** -- The stream\_id specifies the type and number of the elementary stream as defined by the stream\_id table, table 1 in 2.4.4.2. Each elementary stream in an ISO/IEC 11172 multiplexed stream shall have a unique stream\_id.

**packet\_length** -- The packet\_length specifies the number of bytes remaining in the packet after the packet\_length field.

**stuffing\_byte** -- This is a fixed 8-bit value equal to "1111 1111" that can be inserted by the encoder for example to meet the requirements of the digital storage medium. It is discarded by the decoder. No more than sixteen stuffing bytes shall be present in one packet header.

**STD\_buffer\_scale** -- The STD\_buffer\_scale is a one-bit field that indicates the scaling factor used to interpret the subsequent STD\_buffer\_size field. If the preceding stream\_id indicates an audio stream,

STD\_buffer\_scale shall have the value "0". If the preceding stream\_id indicates a video stream, STD\_buffer\_scale shall have the value "1". For all other stream types, the value may be either "1" or "0".

**STD\_buffer\_size** -- The STD\_buffer\_size is a 13-bit unsigned integer defining the size of the input buffer, BS<sub>n</sub>, in the system target decoder. If STD\_buffer\_scale has the value "0" then the STD\_buffer\_size measures the buffer size in units of 128 bytes. If STD\_buffer\_scale has the value "1" then the STD\_buffer\_size measures the buffer size in units of 1024 bytes. Thus:

```
if (STD_buffer_scale == 0)
    BSn = STD_buffer_size * 128;
else
    BSn = STD_buffer_size * 1024;
```

The encoded value of the STD buffer size takes effect immediately when the STD\_buffer\_size field is received by the MPEG system target decoder.

**presentation\_time\_stamp** -- The presentation\_time\_stamp (PTS) is a 33-bit number coded in three separate fields. It indicates the intended time of presentation in the system target decoder of the presentation unit that corresponds to the first access unit that commences in the packet. The value of PTS is measured in the number of periods of a 90kHz system clock with a tolerance specified in 2.4.2. Using the notation of 2.4.2 the value encoded in the presentation\_time\_stamp is:

$$PTS = \text{NINT}(\text{system\_clock\_frequency} * (tp_n(k))) \% 2^{33}$$

where

$tp_n(k)$  is the presentation time of presentation unit  $P_n(k)$ .

$P_n(k)$  is the presentation unit corresponding to the first access unit that commences in the packet data. An access unit commences in the packet if the first byte of a video picture start code or the first byte of the synchronization word of an audio frame (see ISO/IEC 11172-2 and ISO/IEC 11172-3) is present in the packet data.

If there is filtering in audio, it is assumed by the system model that filtering introduces no delay, hence the sample referred to by PTS at encoding is the same sample referred to by PTS at decoding.

**decoding\_time\_stamp** -- The decoding\_time\_stamp (DTS) is a 33-bit number coded in three separate fields. It indicates the intended time of decoding in the system target decoder of the first access unit that commences in the packet. The value of DTS is measured in the number of periods of a 90 kHz system clock with a tolerance specified in 2.4.2. Using the notation of 2.4.2 the value encoded in the decoding\_time\_stamp is:

$$DTS = \text{NINT}(\text{system\_clock\_frequency} * (td_n(j))) \% 2^{33}$$

where

$td_n(j)$  is the decoding time of access unit  $A_n(j)$ .

$A_n(j)$  is the first access unit that commences in the packet data. An access unit commences in the packet if the first byte of a video picture start code or the first byte of the synchronization word of an audio frame (see ISO/IEC 11172-2 and ISO/IEC 11172-3) is present in the packet data.

**packet\_data\_byte** -- packet\_data\_bytes shall be contiguous bytes of data from the elementary stream indicated by the packet's stream\_id. The byte-order of the elementary stream shall be preserved. The number of packet\_data\_bytes, N, may be calculated from the packet\_length field. N is equal to the value indicated in the packet\_length minus the number of bytes between the last byte of the packet\_length field and the first packet\_data\_byte.

In the case of a video stream, packet\_data\_bytes are coded video data as defined in ISO/IEC 11172-2. In the case of an audio stream, packet\_data\_bytes are coded audio data as defined in ISO/IEC 11172-3. In the case of a padding stream, packet\_data\_bytes consist of padding bytes. Each padding byte is a fixed bit-string with the value "1111 1111". In the case of a private stream (type 1 or type 2), packet\_data\_bytes are user definable and will not be defined by ISO/IEC in the future. The contents of packet\_data\_bytes in reserved streams may be specified in the future by ISO/IEC.

## 2.4.5 Restrictions on the multiplexed stream semantics

### 2.4.5.1 Buffer management

The ISO/IEC 11172 multiplexed stream,  $M(i)$  in the notation described in 2.4.2, shall be constructed and  $t_m(i)$  shall be chosen so that the input buffers of size  $BS_1$  through  $BS_n$  neither overflow nor underflow in the system target decoder. That is:

$$0 \leq F_n(t) \leq BS_n \quad \text{for all } t \text{ and } n$$

and  $F_n(t) = 0$  instantaneously before  $t = t_m(0)$ .

$F_n(t)$  is the instantaneous fullness of STD buffer  $B_n$ .

For all ISO/IEC 11172 multiplexed streams, the delay caused by system target decoder input buffering shall be less than or equal to one second. The input buffering delay is the difference in time between a byte entering the input buffer and when it is decoded.

Specifically:

$$t_{d_n(j)} - t_m(i) \leq 1$$

For all bytes  $M(i)$  contained in access unit  $j$ .

### 2.4.5.2 Frequency of coding the system\_clock\_reference

The ISO/IEC 11172 multiplexed stream,  $M(i)$ , shall be constructed so that the time interval between the final bytes of system\_clock\_reference fields in successive packs shall be less than or equal to 0,7 s. Thus:

$$|t_m(i) - t_m(i')| \leq 0,7 \text{ s}$$

for all  $i$  and  $i'$  where  $M(i)$  and  $M(i')$  are the last bytes of consecutive system\_clock\_reference fields.

### 2.4.5.3 Frequency of presentation\_time\_stamp coding

The ISO/IEC 11172 multiplexed stream  $M(i)$  shall be constructed so that the maximum difference between coded presentation\_time\_stamps is 0,7 s. Thus

$$|t_{p_n(k)} - t_{p_n(k'')}| \leq 0,7 \text{ s}$$

for all  $n$  and all  $k, k''$  satisfying:

- 1)  $P_n(k)$  and  $P_n(k'')$  are presentation units for which presentation\_time\_stamps are coded;
- 2)  $k$  and  $k''$  are chosen so that there is no presentation unit,  $P_n(k')$  with a coded presentation\_time\_stamp and with  $k < k' < k''$ ;
- 3) no discontinuity (as defined in 2.4.5.4) exists in elementary stream  $n$  between  $P_n(k)$  and  $P_n(k'')$ .

### 2.4.5.4 Conditional coding of time stamps

For each elementary stream of an ISO/IEC 11172 stream, the presentation\_time\_stamp shall be encoded in the packet in which the first access unit of that elementary stream commences. For the purposes of this clause a video access unit commences in a packet if the first byte of the picture\_start\_code is present in the packet data (see ISO/IEC 11172-2). An audio access unit commences in a packet if the first byte of the synchronization word of the audio frame is present in the packet data (see ISO/IEC 11172-3).

A discontinuity exists at the start of presentation unit  $P_n(k)$  in an elementary stream  $n$  if the presentation time  $t_{p_n(k)}$  is greater than the largest value permissible given the specified tolerance on the system\_clock\_frequency. If a discontinuity exists in any elementary audio or video stream in the ISO/IEC

11172 multiplex then a presentation\_time\_stamp shall be encoded referring to the first access unit after each discontinuity.

Presentation\_time\_stamps may be present in any packet header with the following exception. If no access unit commences in the packet data, the presentation\_time\_stamp shall not be present in the packet header. If a presentation\_time\_stamp is present in a packet header it shall refer to the presentation unit corresponding to the first access unit that commences in the packet data.

A decoding\_time\_stamp shall appear in a packet header if and only if the following two conditions are met:

- 1) A presentation\_time\_stamp is present in the packet header.
- 2) The decoding time differs from the presentation time.

#### 2.4.5.5 Frequency of coding STD\_buffer\_size in packet headers

The STD\_buffer\_scale and STD\_buffer\_size fields shall occur in the first packet of each elementary stream and again whenever the value changes. They may also occur in any other packet.

#### 2.4.5.6 Coding of system header

The system header may be present in any pack, immediately following the pack header. The system header shall be present in the first pack of an ISO/IEC 11172 multiplexed stream. The values encoded in all the system headers in the ISO/IEC 11172 multiplexed stream shall be identical.

#### 2.4.6 Constrained system parameter stream

An ISO/IEC 11172 multiplexed stream is a "constrained system parameters stream" (CSPS) if it conforms to the bounds specified in this clause. ISO/IEC 11172 multiplexed streams are not limited to the bounds specified by the CSPS. A CSPS may be identified by means of the CSPS\_flag defined in the stream header (see 2.4.3.2). The CSPS is a subset of all possible ISO/IEC 11172 multiplexed streams.

#### Packet rate

In the CSPS, the maximum rate at which packets shall arrive at the input to the system target decoder is 300 packets per second if the value encoded in the mux\_rate field is less than or equal to 5 000 000 bits/s. For higher bit-rates the CSPS packet rate is bounded by a linear relation to the value encoded in the mux\_rate field.

Specifically, for all packs  $p$  in the ISO/IEC 11172 multiplexed stream,

$$NP \leq (tm(i') - tm(i)) * 300 * \max \left[ 1, \frac{R_{max}}{5 * 10^6} \right]$$

where

$$R_{max} = 8 * 50 * rate\_bound \quad \text{bits/s}$$

NP is the number of packet\_start\_code\_prefixes and system\_header\_start\_codes between adjacent pack\_start\_codes or between the last pack\_start\_code and the iso\_11172\_end\_code.

tm(i) is the time, measured in seconds, encoded in the system\_clock\_reference of pack  $p$ .

tm(i') is the time, measured in seconds, encoded in the system\_clock\_reference for pack  $p+1$ , immediately following pack  $p$ , or in the case of the final pack in the ISO/IEC 11172 multiplexed stream, the time of arrival of the last byte of the iso\_11172\_end\_code.

#### System target decoder buffer size

In the case of a CSPS the maximum size of each input buffer in the system target decoder is bounded. Different bounds apply for video elementary streams and audio elementary streams.

In the case of a video elementary stream in a CSPS the following applies:

In 2.4.3.2 of ISO/IEC 11172-2 the horizontal picture size, `horizontal_size`, and the vertical picture size, `vertical_size`, are defined. If the values encoded in `horizontal_size` and `vertical_size` meet the constraints on the picture size specified for the `constrained_parameters_flag` in 2.4.3.2 of ISO/IEC 11172-2, then

$$BS_n \leq 46 * 1\,024 \text{ bytes.}$$

For all other video elementary streams in a CSPS,

$$BS_n \leq \max \left[ 46 * 1\,024, \frac{R_{vmax} * 46 * 1024}{1\,856\,000} \right] \text{ bytes}$$

where  $R_{vmax}$  is the greatest value of video `bit_rate` specified or used in the elementary video stream; reference subclause 2.4.3.2 of ISO/IEC 11172-2.

In the case of an audio elementary stream in a CSPS the following applies:

$$BS_n \leq 4\,096 \text{ bytes.}$$

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## Annex A

(informative)

### Description of the system coding layer

#### A.1 Overview

This part of ISO/IEC 11172 specifies the syntax and semantics of the system layer coding of combined coded audio, video, and private data. The coding specification provides data fields and semantic constraints on the data stream to support the necessary system functions. These include the synchronized presentation of decoded information, the construction of a multiplexed stream, the management of buffers for coded data, start-up and random access, and absolute time identification.

While the specification applies to the coded bitstream, there are implications on the functions that encoders and decoders must perform and the degrees of freedom given to them.

The encoding system performs coding of audio and video data, coding of system layer information, and multiplexing. Coding the system layer information includes creating time-stamps: presentation time-stamps (PTS) and decoding time-stamps (DTS) fields are used for synchronization of audio and video. System clock reference (SCR) fields are used in conjunction with PTS and DTS fields for synchronization and buffer management. The use of a common time base, the system time-clock (STC), to unify the measurement of the timing of coded data (SCR) and the timing of the presentation of data (the PTS and DTS fields), ensures correct synchronization and buffer management.

The decoding system performs parsing and demultiplexing of the ISO/IEC 11172 multiplexed stream, the system functions listed above, and the decoding and presentation of elementary streams. The specification of the system is written in terms of an idealized reference decoder known as the system target decoder (STD). The purpose of the STD is to provide a clear, simple model of the decoding system so that the various terms can be specified unambiguously. In the case of video data the term "presentation unit" (PU) refers to decoded pictures, and in the case of audio data it refers to decoded audio frames. The term "access unit" (AU) refers to the coded representation of presentation units.

This annex explains how the system functions are provided by encoders and decoders and the degrees of freedom that are available. Clause A.2 outlines the operation of an encoding system, and clause A.3 outlines the functions of a decoder. The rationale behind constants in the normative clauses (eg 0,7 s or 90 kHz) is described in A.3.4. Clause A.4 describes a set of parameters suitable for use on CD-ROM devices, clause A.5 contains a worked example of a system bitstream, and clause A.6 illustrates the structure of the multiplex.

#### A.2 Encoder operations

##### A.2.1 Degrees of freedom

Flexibility in the system layer syntax allows a wide latitude in creating the multiplexed bitstream. Multiple elementary streams of audio, video, private and padding data can be combined in a practical way into a single stream. Two private data streams of different types are provided. One type is completely private and the other includes the syntax defined in this International Standard to support synchronization and buffer management. If more than two private data streams are needed, an unlimited number of sub-streams may be defined. Up to 32 ISO/IEC 11172-3 audio and 16 ISO/IEC 11172-2 video streams may be multiplexed simultaneously. Packs and packets (the elementary units of the multiplexed stream defined in 2.4.3.2 and 2.4.3.3) may vary in length in order to allocate different bitrates to different streams, or for other reasons. Elementary streams and multiplexed bitstreams may vary their rates from time to time, or operate at a fixed bitrate. The amount of buffering needed in the STD model decoder may be specified individually for each elementary stream. To facilitate random access, the encoding system may include frequent occurrences of the information needed to start decoding, such as SCR, PTS and system headers. Applications will be further assisted if this information is located in the data stream close to the information needed in elementary streams to start decoding (for example video sequence headers and I-Pictures - see ISO/IEC 11172-2).

The encoder has the option of following a set of specific constraints and setting the Constrained System Parameter Stream (CSPS) flag, the `system_audio_lock` flag, or the `system_video_lock` flag. These optional flags may be set independently of one another. The Constrained System Parameters are a defined sub-set of all possible system layer parameters and encoding options. Their purpose is to define a restricted set of parameters and options that can be decoded by economical decoders while being broad enough in application to gain widespread use. The `system_audio_lock_flag` indicates that all the audio streams have an exact relationship to the system clock frequency. The `system_video_lock_flag` indicates that all the video streams have an exact relationship to the system clock frequency.

## A.2.2 Synchronization

This part of ISO/IEC 11172 provides for end-to-end synchronization of the complete encoding and decoding process. This function is provided through the use of time-stamps, particularly Presentation Time-stamps (PTS). This end-to-end synchronization is illustrated in figure A.1 which includes a prototypical encoder and a prototypical decoder. While these prototypical encoding and decoding systems are not normative, they illustrate the functions expected of real systems.

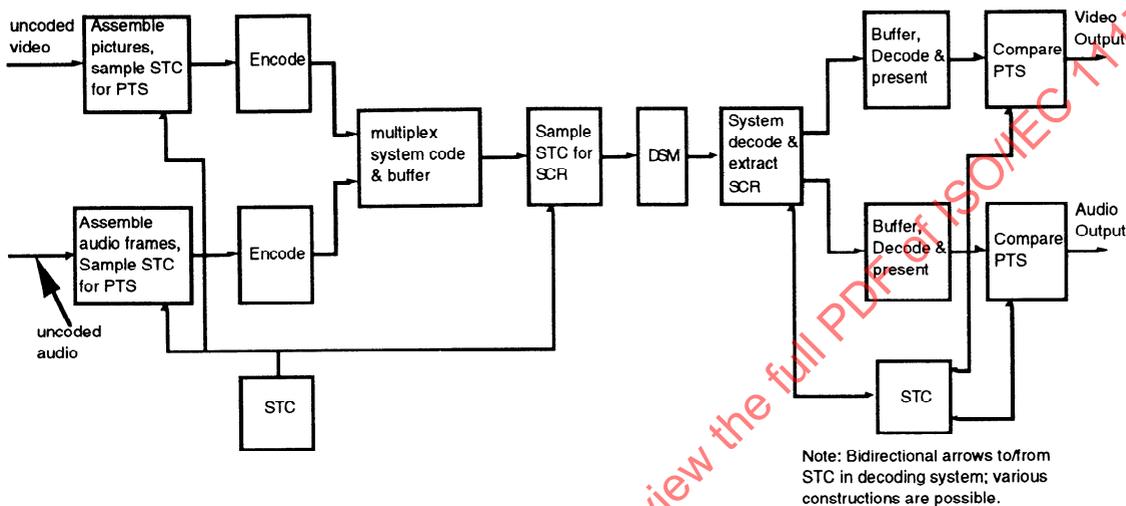


Figure A.1 -- Prototype encoder and decoder

In the prototypical encoding system, there is a single system time-clock (STC) which is available to the audio and video encoders. Audio samples entering the audio encoder are organized into audio presentation units (PU). Some, but not necessarily all, of the audio PUs have PTS values associated with them, which are samples of the STC at the time the first sample of the PU is input to the encoder. Likewise, video pictures enter the video encoder, and the STC values at the times that this occurs are used to create video PTS fields. SCR values represent the time when the last byte of the SCR field leaves the encoder.

This part of ISO/IEC 11172 specifies the encoder and decoder functions in terms of a reference decoder model known as the system target decoder (STD). In this model, video pictures and audio presentation units are presented to the user instantaneously. Actual decoders will generally introduce post-processing and presentation delays. These decoding delays should not be compensated by real encoders. Real encoders must generate bitstreams that play correctly on the idealised STD. Doing this may involve, for instance, choosing the value of the PTS at the time corresponding to the middle of a raster-scanned picture. Such an offset is acceptable providing that it is constant, does not introduce jitter into the sequence of PTS values, and the constraints on the bitstream buffering are respected. The delays that occur in any specific, real decoder must be compensated in that decoder, not the encoder.

SCR and PTS fields, and DTS where required, must be inserted by the encoder at intervals not exceeding 0,7 s as measured by the values contained in the fields. The time interval refers to coded data time for SCR fields, and presentation time for PTS and DTS fields. These fields need not be periodic, and they may be encoded more frequently than the minimum time specified.

Because clock frequencies generally deviate from their nominal values, the use of independent clocks for the generation of PTS, DTS and SCR fields would result in synchronization or buffer management problems. Therefore all the PTS, DTS and SCR fields in the multiplexed stream must be samples of the same STC or have values that are equivalent to those which would have been obtained from a single clock. It is not

permissible, for example, to use independent clocks to produce the PTS and SCR fields in the various streams. This requirement is specified via the definition of the "system\_clock\_frequency" in the definitions of PTS, DTS, and SCR.

While there is a specification for the frequency tolerance of the "system\_clock\_frequency" function used to create the time-stamps, there are no explicit specifications in ISO/IEC 11172-2 and ISO/IEC 11172-3 for the accuracy of the picture rate, audio sample rate, or bitrate, nor for the jitter of these parameters. This issue will be addressed in the compliance testing specification.

In practice the picture rate and audio sample rate will not exactly match the nominal rate unless they are specifically locked to the STC. There is no requirement that these rates should be locked to the system\_clock\_frequency. However, if either or both rates are locked and have an exact relationship to the system\_clock\_frequency the encoder may record this by setting the system\_audio\_lock\_flag and/or the system\_video\_lock\_flag. In the case where the audio sample rates and video picture rates are not locked to the STC, the values implied by the PTS fields should closely match the nominal rates indicated in the elementary data streams in order to avoid problems in decoders.

### A.2.3 Multiplexing

Data from elementary streams are kept distinct by the use of packets, each having a packet start code that identifies the stream. A data packet never contains data from more than one elementary stream and byte ordering is preserved. Thus, after removing the packet headers, packet data from all packets with a common stream identifier are concatenated to recover a single elementary stream.

There is wide latitude in how the multiplex is constructed (the size of packets and the relative placement of packets from different streams). The multiplex is constrained primarily by the STD model, including the specified buffer sizes. The multiplex must be constructed by the encoder in such a way as to ensure that the STD buffers do not overflow or underflow.

In general, short packets require less STD buffering but more system coding overhead than large packets. Other considerations, such as the need to align packets with the sectors on specific storage media, may influence the choice of packet length. A discussion of these factors is given for the particular case of CD-ROM in clause A.4.

Multiplexing may occur either in conjunction with the encoding of elementary streams or the operations may be independent. If multiplexing is combined with coding then the system is free to use the full range of the STD buffer. If multiplexing is independent from the coding, the elementary encoders must allow sufficient space in the STD buffers to allow for multiplexing. In the case of CD-ROM sector-based multiplexing, a headroom of  $6 * 1024$  bytes is generally sufficient. This is why the buffering limit is 40 kBytes in the video constrained parameters (see ISO/IEC 11172-2) and 46 kbytes in the constrained system parameters stream (see 2.4.6).

Coding for use with a bursty DSM or channel in general requires additional buffering in the STD model beyond that required with a constant-latency DSM or channel. The additional buffering required may be reduced through the careful use of multiplexing and the mux\_rate field. The STD uses a byte arrival schedule specified by the SCR and mux\_rate fields. In some cases the STD byte arrival schedule can be made to duplicate the actual delivery schedule of the bursty DSM or channel, permitting optimization of STD buffer usage.

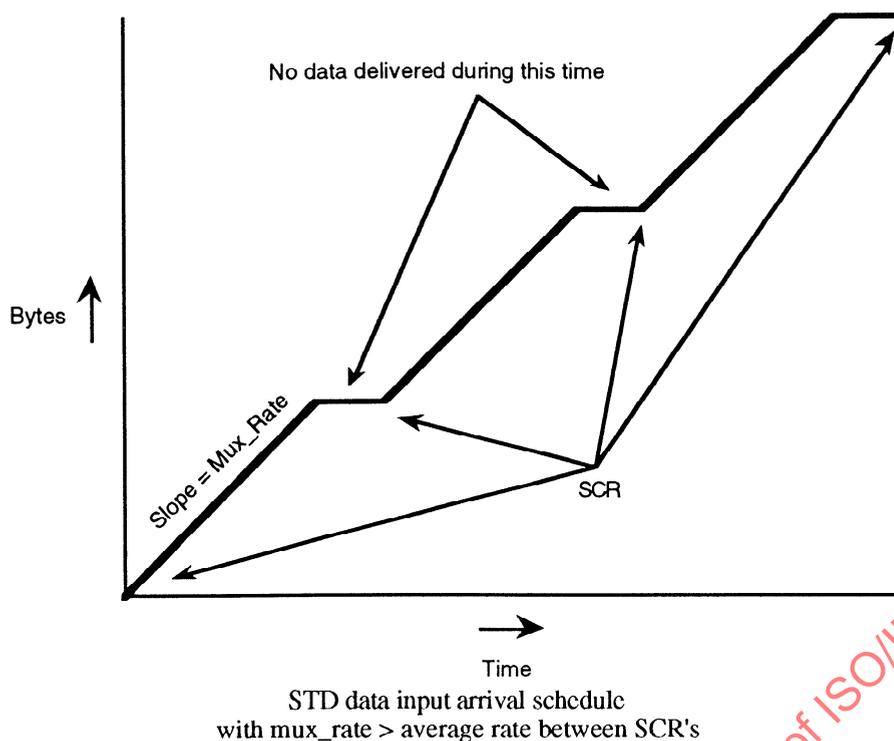


Figure A.2 -- STD data input arrival schedule

#### A.2.4 Encoder constraints caused by decoder buffering

It is the encoding system's responsibility to ensure that the STD model never overflows or underflows its buffers. Encoders must specify the sizes of the buffers used in the STD model for each stream, and limits on the sizes of the STD buffers valid for the entire ISO/IEC 11172 multiplexed stream must be placed in the system header packet. The STD model specifies the exact times when each byte of coded data enters and leaves each buffer in terms of the common system\_clock frequency. This timing is specified by a byte arrival schedule, which is specified in the data stream via the SCR and mux\_rate fields, and by a byte removal schedule which is specified by the PTS and DTS fields, the audio sample rate, audio AU and PU sizes, the video picture rate, and the STD model with instantaneous elementary stream decoders.

The STD is a model of a decoder. Encoders must apply the STD model in the creation of multiplexed bitstreams, but real decoders need not be implemented with the same architecture as the STD decoder. Variations in access unit size cause the elementary stream encoders to contribute to part of the STD buffer occupancy, and the action of multiplexing contributes to the rest. In the STD model a single buffer is used for both demultiplexing and elementary stream decoding.

Encoders setting the CSPS\_flag must not specify STD buffer sizes larger than those permitted in the constrained system parameter limits. For audio, buffer  $B_n$  in the system target decoder must not be larger than 4 096 bytes. For video the maximum size of buffer  $B_n$  in the STD depends on both the picture size and the video bitrate,  $R_V$ . If either the picture size is less than or equal to the maximum allowed by the constrained parameter video stream (see ISO/IEC 11172-2) or  $R_V$  is less than or equal to 1 856 000 bits/s then the maximum size of  $B_n$  is 46 \* 1 024 bytes. Otherwise the maximum size is:

$$\frac{46 * 1\,024 * R_V}{1\,856\,000} \text{ bytes.}$$

Note that this maximum STD  $B_n$  size is larger than the 40 kbyte maximum size of the Video Buffer Verifier to accommodate demultiplexing. If the video bitrate is variable, the peak video bitrate throughout the video stream is used for  $R_V$  in the above formula.

### A.2.5 Stream characterization

The System Header is a special packet that contains no elementary stream data. Instead it indicates decoding requirements for each of the elementary streams. It indicates a number of limits that apply to the entire ISO/IEC 11172 multiplexed stream, such as data rate, the number of audio and video streams, and the STD buffer size limits for the individual elementary streams. A decoding system may use these limits to establish its ability to play the stream.

The system header contains a flag indicating whether or not the data stream is encoded for constant rate delivery to the STD. If the data rate averaged over the time intervals between SCRs is constant throughout the stream and, when rounded upwards in units of 50 bytes/s, is equal to the value in the mux\_rate field, the constant rate flag may be set. If the mux\_rate fields indicate a rate higher than this, data is delivered to the STD in bursts at the rates indicated by the mux\_rate fields. The mux\_rate field will never be lower than that implied by the SCR fields.

The system header must be in the first pack of the ISO/IEC 11172 multiplexed stream. It may be repeated within the stream as often as necessary. In broadcast applications this may be desirable.

Real-time encoding systems must calculate suitable limits for the values in the header before starting to encode. Non-real-time encoders may make two passes over the data to find suitable values.

### A.2.6 Padding stream

A padding stream is provided. It may be used to maintain a constant total data rate, to achieve sector alignment, or to prevent buffer underflow. As the padding stream is not associated with decoding and presentation, it has neither a buffer in the STD model nor PTS or DTS fields.

Stuffing of up to 16 bytes is allowed within each data packet. This can be used for purposes similar to that of the padding stream and is well suited to providing word (16-bit) or long word (32-bit) alignment in applications where 8-bit alignment is not sufficient. Use of stuffing bytes is the only available method of padding when the number of bytes required for stuffing is less than the minimum size of the padding packet, which is equal to the size of the stream header.

### A.2.7 Insertion of private data

Two private stream types, private\_stream\_1 and private\_stream\_2, are provided for applications not defined in ISO/IEC 11172. Private\_stream\_1 follows the same syntax as audio and video streams. It may contain stuffing bytes, a buffer size field, and PTS and DTS fields. The use of these fields is not specified in ISO/IEC 11172. Private\_stream\_2 is similar except that no syntax is specified for stuffing bytes, buffer sizes, PTS or DTS fields.

Although only two private stream identifiers are provided, private streams may be designed to include branching fields to support an unlimited number of private sub-streams. This mechanism is not defined in ISO/IEC 11172.

## A.3 Decoder operations

Figures A.3 and A.4 show two different models of an implementation of a decoding system that are used in the following clauses to illustrate the operation of the system. Both models represent possible implementations.

### A.3.1 Decoder synchronization

#### A.3.1.1 Time-Stamps

The PTS, DTS and SCR fields are the basis for synchronization in decoders. Decoders parse the data stream and extract the PTS or DTS fields contained in the packet coding layer together with the relevant coded data. The PTS and DTS fields are associated with the first access unit (AU) that commences in a packet containing a PTS and/or DTS field. Picture start codes and audio syncwords are not necessarily located at the start of packets, and there may be more than one AU commencing in a packet.

PTS and DTS fields are not necessarily encoded for each picture or audio PU. They are required to occur with intervals not exceeding 0,7 s. This bound allows the construction of a control loop using the PTS values which has guaranteed stability with a known bandwidth. For those PUs for which PTS is not encoded, the decoder can approximate the correct value as the sum of the most recent PTS and an increment. The increment is the nominal number of system\_clock\_frequency cycles per PU times the number of PUs since the last PTS.

DTSs specify the time at which all the bytes of an access unit are removed from the buffer of an elementary stream decoder in the STD model. The STD model assumes instantaneous decoding of access units. In audio streams, and for B-pictures in video streams, the decoding time is the same as the presentation time and so only the PTSs are encoded; DTS values are implied. In video streams, for I-pictures and P-pictures the DTS values are nominally equal to the PTS value minus the number of picture periods of video recording delay multiplied by the picture period, in units of the 90kHz STC. The DTS and PTS need not be encoded for every access unit. Intervening values may be calculated from known DTS and PTS values and the rate of PUs for each stream.

Similarly, SCR values, which measure the time of events in the coded data stream, are required to occur with intervals not exceeding 0,7 s. Again, this allows construction of a controller using SCR values with a guaranteed stability.

#### A.3.1.2 Clock relationships

A decoding system, including all of the synchronized decoders and the source of the coded data, must have exactly one independent time-master. This fact is a natural result of the requirement to avoid overflow and underflow in finite size buffers, while maintaining synchronization of the presentation of data. All other synchronized entities must slave the timing of their operation to the time-master. If a decoder attempts to have more than one simultaneous time-master it may experience problems with buffer management or synchronization.

A decoder system has complete freedom in choosing which entity is the time-master. Typically these entities include the video decoder, the audio decoder, a separate STC, or the data source. Whichever entity is the time-master must communicate to the others the correct value of the STC. A time slave will typically maintain a local STC which is incremented nominally at 90 kHz between updates or corrections. In this way each entity has a continuously updated value of the STC which is nominally correct and which it uses to compare with the time-stamps.

Two examples are presented to illustrate different approaches to designing a decoder. One uses the audio decoder's clock as the time-master, and the other relies on the DSM clock as the time-master.

#### A.3.1.3 Example: audio as time-master

In this first example, the audio decoder is the time-master in a decoding system. Its operation is described here and illustrated in figure A.3.

The system time clock (STC) is typically initialized to be equal to the value encoded in the first SCR field when that field enters the decoder's buffer. Thereafter the audio decoder controls the STC. As the audio decoder decodes audio AUs and presents audio PUs, it finds PTS fields associated with some of the audio PUs. As the beginning of each PU is output to the user, the associated PTS field contains the correct value of the decoder's STC in an idealized decoder following the STD model. The audio decoder may use this value to update the STC immediately, or to control the STC values via a control loop.

The other decoders then use this STC to determine the correct time to present their decoded data, at the times when their PTS fields are equal to the current value of the STC.

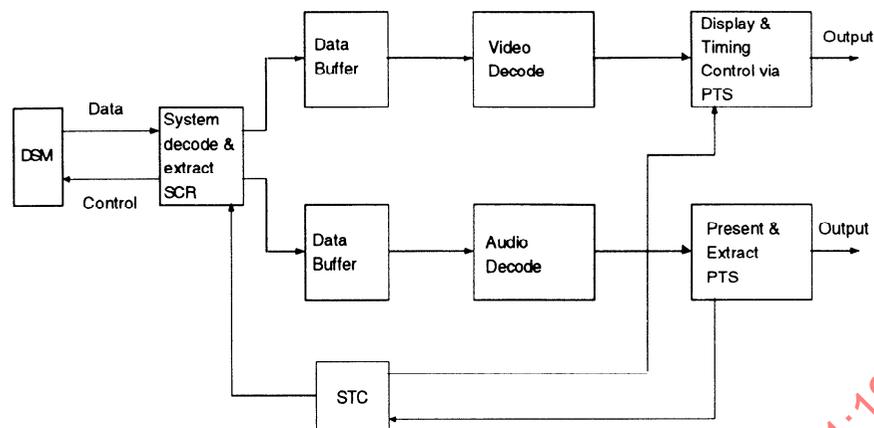


Figure A.3 -- Example of decoding system - audio time-master

Note that the data source (DSM or channel) must provide data to the decoder on a schedule determined by the SCR and mux\_rate field values and the decoder's STC. This time relationship is necessary in order to manage the decoder buffers. Buffer management is described further in A.3.3.

The DSM control mechanism obtains data from the DSM at a rate at least equal to that specified in the mux\_rate field for each pack until the next SCR field is received. It is not necessary to obtain more data from the DSM until the STC value equals the most recently received SCR value. If more data is read, more buffering will be required.

#### A.3.1.4 Example: DSM as time-master

In this second example, illustrated in figure A.4, the DSM is the time-master, and the audio and video decoders are implemented as separate decoder subsystems, each receiving the complete multiplexed data stream and extracting and using only that portion of the stream needed.

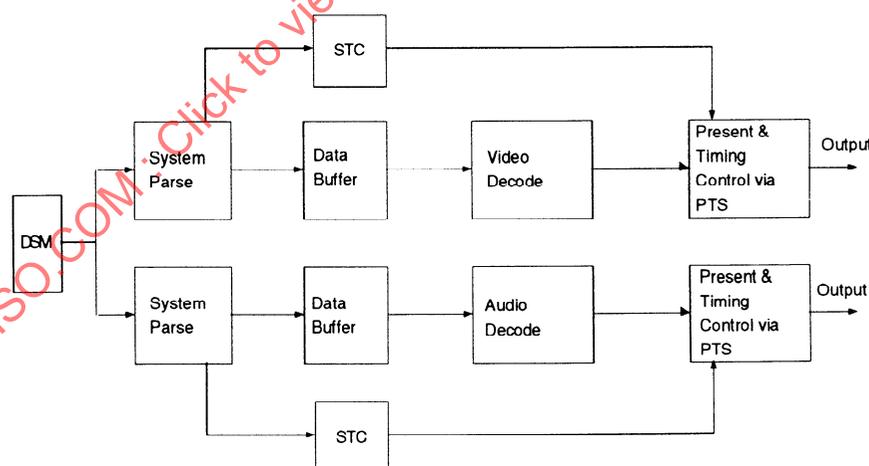


Figure A.4 -- Example of decoding system - DSM time-master

Each decoder receives and parses the complete multiplexed data stream and extracts the system layer information and the coded data needed by that decoder. Synchronization is implemented by the individual decoders slaving their timing to the DSM. The DSM timing is indicated by the SCR fields, which contain the expected value of the decoder's STC at the time that the last byte of the SCR is received by the decoder. Each decoder has a separate STC that is initialized to the first value of SCR received and increments at a nominal 90 kHz rate. The correct timing of the STC is maintained by ensuring that the STC is equal to the SCR values at the time that the SCRs are received. The STC may be maintained either by updating the STC with the value of the SCRs or via a control loop, using the SCR values as reference inputs.

### A.3.2 Decoder start-up synchronization

#### A.3.2.1 Finding start codes at random access

The syntax specified in this part of ISO/IEC 11172 fully specifies the location of the system start codes. If decoding of the stream begins with the first byte, all start codes can be found without uncertainty. After random access to the stream, or in broadcast applications, where decoding may begin at an arbitrary byte, the problem of finding start codes must be solved.

The thirty-two bit pack and packet start codes are constructed so that they cannot occur in video data, which is expected to be the largest portion of the data in most applications. Therefore parsing can start after a random access with a low probability of incorrectly identifying packet data as a start code. Nonetheless, the probability is not zero because start code emulation may occur in audio or private streams.

If additional protection is desired, for example where a large number of audio streams are multiplexed or there is a large amount of private data, a parser could detect a 32-bit pack\_start\_code followed 8 bytes later by a 24-bit packet\_start\_code\_prefix. Once a probable system start code is found by a parser, the packet\_length field may be used to predict the position of the next start code. In this way the probability of incorrectly identifying coded data as a start code decreases geometrically as each successive start code is found. A decoder performing this function has the options of either discarding or saving data until parsing is operating with a sufficient level of confidence. Decoding can then begin with the earliest available data for which system start codes are known.

If the application includes a means of directly addressing known system start codes, then the probability of incorrect parsing of start codes can be made zero.

It is possible for a decoder to "switch channels"; that is, for it to stop decoding one ISO/IEC 11172 multiplexed stream and to start decoding another. This function generally requires that decoding of the second stream start at an unknown byte location. Switching channels is possible, but involves the flushing of decoder buffers and introduces delay. The amount of delay depends on the frequency of start codes in the second stream, as well as on the exact location where decoding starts.

#### A.3.2.2 System layer startup considerations

Once the decoding system has locked on to the data stream it can begin decoding data.

Decoding systems determine the correct time to start decoding by comparing the DTS (or PTS) fields extracted from the stream or computed as described above, with the current value of the STC. The delay from the time the decoder begins to process data until it can begin to present decoded data is bounded from below by the start-up delay implied by the SCR and PTS fields. A decoder following the STD model may produce decoded output as soon as the following conditions are met:

- a) At least one SCR field has been extracted and the STC is synchronized with the DSM via the SCRs and mux\_rate fields.
- b) At least one PTS has been extracted.
- c) A complete coded AU is available and the correct associated PTS value (coded or computed) is known.
- d) The PTS for an AU which is available is equal to the current STC value.

DTS fields may be used by a decoder to control input and reorder buffering.

In addition there may be other constraints imposed by the elementary decoders (for example the need for sequence layer information and an I-picture in video coded according to ISO/IEC 11172-2).

This procedure guarantees that the streams will be synchronized, but it does not necessarily ensure that the start up will be simultaneous. It may be necessary to discard some decoded audio or video and to wait until all elementary decoders are ready. In general there will not be audio and video PUs that start at the same time or with the same PTS values. Thus exactly simultaneous start up of audio and video may require muting some audio samples.

### A.3.2.3 Coding layer startup considerations

Apart from the system layer decoding considerations, start-up may not be possible immediately, particularly in the case of video. On startup a video decoder requires video sequence layer information and must begin decoding at an I-picture at the start of a group of pictures. In some applications it may be possible to retain the video sequence layer information from a previous access. For more information see ISO/IEC 11172-2

### A.3.2.4 Compensation of actual decoding delays

The system target decoder (STD) is a model of a decoder. Encoders must apply the STD model in the creation of a multiplexed stream, but real decoders need not be implemented with the STD architecture. In practice real decoders will not be instantaneous. If a real decoder cannot remove and decode an entire access unit instantaneously, it will need to delay completion of the processing and presentation compared with the values specified in the STD, and a buffer larger than that specified in the STD will be needed. The effective values of PTS and DTS used for timing may be modified in the decoder to accommodate the decoding delay.

For example, if a video decoder requires one picture period to decode a picture, it may delay completion of decoding all the pictures by one picture period with respect to the values indicated by DTS. This in turn requires additional video decoder buffering (the size of the average coded picture). Proper synchronization can be maintained at the output of the entire system by adding one picture period to the effective values of PTS and DTS for all the elementary streams to compensate for this delay.

### A.3.2.5 Channel smoothing

ISO/IEC 11172 multiplexed streams are, in their most general form, channel independent. They do not assume a specific set of channel characteristics. Decoding from a bursty DSM or channel in general requires additional smoothing buffers not present in the STD model. It is up to decoders to compensate for deviations of a real channel from the STD byte arrival schedule derived from the SCR and mux\_rate fields.

In some cases the STD byte arrival schedule can be made to duplicate the actual performance of a bursty DSM or channel. In these cases no extra channel smoothing is required and the performance of the system will be optimized.

### A.3.3 Buffer management in the decoder

Buffer management uses the system target decoder (STD) model. Elementary stream decoder buffers are guaranteed not to overflow or underflow during decoding as long as the data stream conforms to the specification, and the complete decoding system is synchronized in terms of SCR and DTS. The STD model precisely specifies the times at which each data byte enters and leaves the buffer of each elementary stream decoder in terms of a common system time-clock.

The upper bound on STD delay specified in 2.4.5.1 is motivated by decoder performance considerations; in conjunction with elementary stream bitrate, the STD delay bounds the size of STD buffers.

ISO/IEC 11172-2 specifies buffer management and start-up delay for video, using the vbv\_delay parameter. When video is combined with system layer coding these specifications may conflict. In this case the information in the system layer prevails.

### A.3.4 Time identification

The absolute time of presentation of the material contained in the coded data stream is indicated in the PTS fields. These fields are defined as modulo  $2^{33}$  values of the 90 kHz STC. If required, PTS fields can be transcoded into other formats such as SMPTE time-codes. There is no requirement that the PTS values be initialized to any particular value at the start of the stream.

An application can find the coded data associated with a particular value of presentation time by searching the system layer coding for values of PTS equal to or within an appropriate range of the desired presentation time. Note that SMPTE-like time-codes are also defined in the video coding layer defined in ISO/IEC 11172-2.

Selection of 90 kHz as the STC frequency is based on the divisibility of 90 KHz by the nominal video picture rates of 24 Hz, 25 Hz, 29,97 Hz, and 30 Hz. Use of 33 bit encoding for PTS fields allows elapsed times of up to 24 h to have distinct coded values. Finally, the maximum 0,7 s SCR, PTS and DTS

intervals specified in 2.4.5.2 is small enough for phase-lock loop stability, but large enough to permit one PTS value per I-picture even when I-pictures occur slightly less often than twice per second.

#### A.4 Parameters for CD-ROM multiplexing

In this clause an example of a multiplex method for CD-ROM is presented. The example is developed for one video and one audio elementary stream. The ISO/IEC 11172 multiplexed stream is stored on a CD-ROM without additional error correction - a mode with 2 324 bytes in each sector. Packs are constructed to be this length so that they may be stored one in each sector. The duration of one sector equals 1/75 s, resulting in a total bitrate of  $8 * 2\,324 * 75 = 1\,394\,400$  bits / s.

The audio stream is coded in stereo with the ISO/IEC 11172-3 audio layer II coding method at a bitrate of 192 000 bits / s = 24 000 bytes / s. The sample rate used is 44 100 samples / s. Audio presentation units are 1 152 samples each, and so the size of an audio access unit equals:

$$\frac{1\,152 * 24\,000}{44\,100} \text{ bytes}$$

As this result is not an integer, most audio access units are 627 bytes but some are only 626 bytes.

The video stream is coded with a bitrate of 1 158 000 bits / s = 144 750 bytes / s. The value of  $B_{vby}$  used is 36 kBytes, leaving sufficient headroom in the 46 kbyte STD buffer of the Constrained System Parameters for the multiplexing.

The packs are to coincide with the sectors. Each pack contains a pack header, one packet of coded audio or coded video and one packet of a padding stream. Each packet of coded audio or coded video data contains exactly 2 250 data bytes. The padding stream ensures that each pack, including the pack header, consists of the number of bytes available in the data field of the sector in which the pack is stored. In sectors where all 2 324 bytes are available for the ISO/IEC 11172 multiplexed stream, packs are 2 324 bytes long. In sectors where less than 2 324 bytes are available, the size of the pack is reduced accordingly by decreasing the size of the padding packet.

The coded audio rate of 24 000 bytes / s, with each sector containing 2 250 bytes, requires an average  $24\,000 / 2\,250 = 10^{2/3}$  audio sectors/s. Similarly, the coded video bitrate of 144 750 bytes / s, with 2 250 bytes per sector, requires an average of  $144\,750 / 2\,250 = 64^{1/3}$  video sectors per second. In total, therefore, exactly 75 sectors of audio and video data are required each second for the combined bitstream, exactly filling the total bandwidth of the CD-ROM.

Interleaving the audio and video sectors must not cause the STD buffers to overflow or underflow. Many interleaving schemes are possible that will lead to a multiplexed stream following the Constrained System Parameters. In this example a simple interleaving scheme is used that repeats every 3 s (225 sectors). The scheme starts with 6 video sectors followed by one audio sector. This pattern is repeated 31 times, resulting in an interleave of 217 sectors. The last pattern in the interleave scheme consists of 7 video sectors followed by 1 audio sector. The three second period of 225 sectors contains 32 audio sectors and 193 video sectors. On average there are  $193/3 = 64\,1/3$  video sectors/second and  $32/3 = 10\,2/3$  audio sectors/second, as required.

## A.5 Example of an ISO/IEC 11172 multiplexed stream

A sample ISO/IEC 11172 multiplexed stream is presented here to illustrate the syntax and semantic rules governing generation of such streams. This example does not use the same parameters defined in the previous clause. The sample stream is a constrained system parameter stream combining two elementary streams: one video and one audio. The elementary streams are assumed to have been generated with the following specifications:

### A.5.1 Audio

Layer II encoding  
 48 kHz sample rate  
 24 000 bytes/s rate for a pair of stereo channels  
 1 152 samples per presentation unit  
 576 bytes per access unit

The stream so generated, with place holders for coded audio and video data, is listed in A.5.9 "Sample data stream".

### A.5.2 Video

Constrained parameter video encoding at 150 000 bytes/s.  
 25 Hz picture rate source.  
 40 \* 1 024 Byte video buffer verifier

The order of pictures at the decoder input is

1I 4P 2B 3B 7P 5B 6B 10P 8B 9B 13I 11B 12B 16P 14B 15B 19P 17B 18B 22P 20B 21B 25I

I pictures coded at 19 000 bytes each  
 P pictures coded at 10 000 bytes each  
 B pictures coded at 2 800 or 2 900 (2 875 byte average) each

### A.5.3 Multiplexing strategy

The example employs packets of length 2 048 bytes for both audio and video. The multiplex starts with thirteen video packets to limit audio buffering requirements. Thereafter, one audio packet is interleaved with every 6 to 7 video packets to match the 6,25 ratio of video bitrate to audio bitrate.

For simplicity, packets are constructed with a common number of packet\_data\_byte entries. Stuffing bytes are used to ensure that all packets have 20 header bytes and 2 028 data bytes.

A pack is generated every third packet. This structure is somewhat arbitrary, but leads to a pack rate of roughly 29 Hz, comfortably over the 1 to 2 Hz requirement of 2.4.5.2 (Coding of the system\_clock\_reference). The cost of such frequent pack formation is not great: all pack headers except the first are 12 bytes long, so pack headers account for some 0,2% of the total bitrate.

The sample bitstream is long word aligned. That is, all packets and all packet data (except the initial padding stream packet) start at 32-bit boundaries. Because the first pack header is 30 bytes long (it contains 18 bytes of system header information), a special padding stream packet appears in the first pack. This 10-byte packet guarantees long word alignment for subsequent packets.

To summarize, the stream is composed of packs and packets as follows:

#### Pack 1

header (includes system_header)	30 bytes
Padding stream packet	10 bytes
Video packet #1	2 048 bytes
Video packet #2	2 048 bytes
Video packet #3	2 048 bytes

## Pack 2

header	12 bytes
Video packet #4	2 048 bytes
Video packet #5	2 048 bytes
Video packet #6	2 048 bytes

## Pack 3

header	12 bytes
Video packet #7	2 048 bytes
Video packet #8	2 048 bytes
Video packet #9	2 048 bytes

## Pack 4

header	12 bytes
Video packet #10	2 048 bytes
Video packet #11	2 048 bytes
Video packet #12	2 048 bytes

## Pack 5

header	12 bytes
Video packet #13	2 048 bytes
Audio packet #1	2 048 bytes
Video packet #14	2 048 bytes

⋮

#### A.5.4 System clock reference (SCR)

Bytes 5 to 9 of every pack header contain encoded system\_clock\_reference fields. The multiplexed stream's data rate is computed from the data in A.5.1 and A.5.3, and the following formula:

$$R_{\text{mux}} = (\text{video data rate} + \text{audio data rate}) * \left(1 + \frac{(\text{packet header\_size} + \text{pack header\_size/packs/packet})}{\text{packet\_data\_size}}\right)$$

$$R_{\text{mux}} = (150\,000 + 24\,000) \left(1 + \frac{20 + 12/3}{2\,028}\right)$$

$$= 176\,059,1717 \text{ bytes/s}$$

This value can be rounded to 176 059 bytes/s without affecting the values in the data stream in this particular example.

$R_{\text{mux}}$  and the 90 kHz clock frequency are used by the encoder to convert SCR field byte indices to system\_clock\_reference values. The first SCR field, equal to 3 904, simply reflects a non-zero starting value for the encoder's clock. Subsequent SCR fields evaluate to

Pack	system_clock_reference
1	3 904
2	7 065
3	10 212
4	13 359
5	16 506
6	19 653
7	22 800
8	25 947

To understand the source of these numbers, consider the second pack's SCR value, SCR2. The SCR2 field occurs 6 180 bytes after the first pack's. SCR2 is related to SCR1 in terms of the elapsed time. For this example's constant rate byte delivery, SCR2 is

$$\begin{aligned} \text{SCR2} &= \text{SCR1} + 6\,180 * 90\,000/176\,059 \\ &= 7\,065 \end{aligned}$$

### A.5.5 Presentation time-stamps (PTS)

The video coding model used for this example leads to coded pictures of the type:

1I 4P 2B 3B 7P 5B 6B 10P 8B 9B 13I 11B 12B 16P 14B 15B 19P 17B 18B 22P 20B 21B 25I

Recalling that coded I, P, and B pictures are assumed to be 19 000, 10 000 and 2 800 to 2 900 bytes, respectively, and that packets contain 2 028 bytes of data each, it follows that picture start codes occur in video packet#1 (I picture), video packet#10 (P picture), video packet#15 (B-picture), etc. This is reflected by the presence of PTS fields in video packets 1, 10, 15, etc., in the sample stream listing.

In this example, N, the number of coded pictures between I pictures equals 12. The number of consecutive B pictures (M-1) between I or P pictures equals two, and thus M=3.

The audio coding model used for this example employs 576 byte access units, hence every 2 048-byte audio packet contains an access unit start code. All audio packets contain PTS fields.

The value of an elementary stream's first Decoding Time-stamp (DTS) field (or PTS if the two are equal) when compared with the initial SCR field, determines the decoder start-up delay for that stream. In the example, the first video DTS field has the value 22 804. The difference between the first pack's SCR value and the first video packet's DTS value is:

$$\begin{aligned} \text{start-up delay} &= (22\,804 - 3\,904 \text{ cycles}) * (1\,000 \text{ ms/s}) / (90\,000 \text{ cycles/s}) \\ &= 210 \text{ ms} \end{aligned}$$

This delay is required to prevent overflow or underflow in the system target decoder. It tells the decoder that the first I picture should be decoded 210 ms and presented 250 ms after reading the last byte of the first SCR field in the multiplexed stream.

Note that the first PTS field in the audio stream equals 26 395, a number slightly lower than the video's. This inequality arises if the video and audio encoders are not turned on at exactly the same instant, and does not imply synchronization error.

The system\_audio\_lock\_flag is set in the system header packet of the the sample bitstream, but the system\_video\_lock\_flag is reset. Therefore, decoders may assume a rational relationship between the audio clock and the system time clock, but may not assume such a relationship between the video clock and the system time clock. PTS and DTS value present in the stream are consistent with exact clocks for both video and audio; in practice, however, because the video clock is not locked some drift would appear in video time-stamps. Over one second, or 90 000 clock cycles, errors of 50 parts per million would lead to PTS values differing from the nominal values by 4 or 5. The discrepancy accumulates over time.

### A.5.6 Decoding time-stamp (DTS)

For I and P pictures, it is generally true that system target decoder operations for decoding and presentation occur at different times. Steady state operation with this example's GOP structure (M=3, N=12) leads to I-

and P-pictures being decoded three picture periods before their presentation. Thus, video packet #10 has DTS equal to 26 404 but PTS is equal to 37 204. This 10 800 clock cycle, 120 ms difference requires the P-picture to be stored in the system target decoder's reorder buffer for 3 picture periods.

Analysis of DTS and PTS values for the first I-picture (video packet #1) reveals a relationship needed to initialize the reorder buffer. The I-picture is decoded when the decoder's clock reaches 22 804, but nothing is displayed. The initialization is complete 40 ms later when the P-picture discussed in the previous paragraph is decoded, and the I-picture is displayed.

The second audio PTS field (value of 35 035) lags the first by 8 640 clock cycles, or 96 ms. Audio presentation units are 1 152 samples long, which at a 48 kHz sampling rate, corresponds to 24 ms. The second audio PTS field, therefore, appears in the stream after the start code for the fifth audio access unit.

### A.5.7 Buffer sizes

The example documents a constrained system parameter stream with pictures conforming to the video constrained parameters defined in Part 2 of this International Standard. The maximum allowable buffer sizes in the STD for such streams are used. These are:

Video streams: 46 \* 1 024 bytes  
Audio streams: 4 \* 1 024 bytes

### A.5.8 Adherence to System Target Decoder (STD)

For a stream to be a valid ISO/IEC 11172 multiplexed stream, it must play on the system target decoder without overflow or underflow of any STD buffer. Tables A.1 and A.2 track buffer occupancy for the STD video and audio buffers, respectively. The tables demonstrate that the one-second long sample bitstream complies with the STD buffering requirements.

Table A.1 -- System target decoder video buffer occupancy

Input Picture Index and Type (in coded order)	End-of-picture delivery time (ms)	Decoding / Presentation time (ms)	Buffer Occupancy (bytes)
--	0	----	----
1I	109	210/250	34 568
4P	178	250/370	21 560
2B	194	290	17 468
3B	211	330	20 928
7P	280	370/490	23 584
5B	297	410	20 196
6B	313	450	22 676
10P	382	490/610	26 664
8B	399	530	21 692
9B	427	570	25 756
13I	548	610/730	27 884
11B	564	650	15 848
12B	580	690	17 900
16P	650	730/850	21 964
14B	678	770	16 816
15B	694	810	20 880
19P	763	850/970	23 016
17B	780	890	20 072
18B	796	930	22 460
22P	866	970/1 090	26 088
20B	882	1 010	23 052
21B	898	1 050	27 308
25I	1 019	1 090/1 210	31 372