
**Information technology — Security
techniques — Digital signatures with
appendix —**

Part 3:

Discrete logarithm based mechanisms

AMENDMENT 2: Optimizing hash inputs

*Technologies de l'information — Techniques de sécurité — Signatures
numériques avec appendice —*

Partie 3: Mécanismes basés sur un logarithme discret

*AMENDEMENT 2: Optimisation des entrées pour la fonction de
hachage*

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Case postale 56 • CH-1211 Geneva 20
Tel. + 41 22 749 01 11
Fax + 41 22 749 09 47
E-mail copyright@iso.org
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Foreword

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The main task of the joint technical committee is to prepare International Standards. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

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Amendment 2 to ISO/IEC 14888-3:2006 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 27, *IT Security techniques*.

It introduces an optimization for the Schnorr Elliptic Curve Digital Signature Algorithms specified in ISO/IEC 14888-3:2006/Amd.1:2010. Whereas this optimization is described in an informative (only) note of ISO/IEC 14888-3:2006/Amd.1:2010, Amendment 2 makes the optimization a normative option. It also corrects various errata in Annexes E and F and updates the date of a reference in the Bibliography.

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Information technology — Security techniques — Digital signatures with appendix —

Part 3: Discrete logarithm based mechanisms

AMENDMENT 2: Optimizing hash inputs

Subclause 6.9.1, Introduction to EC-SDSA

Replace the text in 6.9.1 with the following:

EC-SDSA (Elliptic Curve Schnorr Digital Signature Algorithm) is a signature mechanism with verification key $Y = [X]G$; that is, the parameter D is equal to 1. The message is prepared such that M_2 is empty and $M_1 = M$ the message to be signed. The witness R is computed as a hash-code of the message M and a random pre-signature $\Pi = [K]G$, by one of two methods, either

normal $R = h(\text{FE2BS}(\Pi_x) \parallel \text{FE2BS}(\Pi_y) \parallel M)$

or

optimized $R = h(\text{FE2BS}(\Pi_x) \parallel M)$.

The first method generates the witness by hashing the concatenation of the x-coordinate of Π , the y-coordinate of Π and the message M . The second method omits the y-coordinate from the hash calculation and thereby improves performance.

The second method is an optimized variant of EC-SDSA (see [40]).

Subclause 6.9.4.4, Computing the witness

Replace the text in 6.9.4.4 with the following:

The signing entity computes $R = h(\text{FE2BS}(\Pi_x) \parallel \text{FE2BS}(\Pi_y) \parallel M)$.

For the optimized variant of EC-SDSA, the signing entity instead computes $R = h(\text{FE2BS}(\Pi_x) \parallel M)$.

Annex A, ASN.1 Module

Add the following entries at the appropriate places in the OID assignments section.

```
id-dswa-dl-EC-SDSA-opt OID ::= { id-dswa-dl ec-sdsa-opt (13) }
```

```
    dswa-dl EC-SDSA-opt
```

```
dswa-dl-EC-SDSA-opt ALGORITHM ::= {
    OID id-dswa-dl-EC-SDSA-opt PARMS HashFunctions
}
```

F.11.2.3, Per message data

Append the following text to F.11.2.3:

For the optimized variant of EC-SDSA,

$$R = h(\text{FE2BS}(IT_x) \parallel M) =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| D7FB8135 | D8EA45E8 | FB3C9059 | F146E263 | 0EF4BD51 | C4006A92 |
| EDB4C8B0 | 849963FB | | | | |

F.11.2.4, Signature

Append the following text to F.11.2.4:

For the optimized variant of EC-SDSA,

$$R =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| D7FB8135 | D8EA45E8 | FB3C9059 | F146E263 | 0EF4BD51 | C4006A92 |
| EDB4C8B0 | 849963FB | | | | |

$$S =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| B46D1525 | 379E02E2 | 32D97928 | 265B7254 | EA2ED978 | 13454388 |
| C1A08F62 | DCCD70B3 | | | | |

F.11.2.5, Verification

Append the following text to F.11.2.5:

For the optimized variant of EC-SDSA,

$$R' = h(\text{FE2BS}(IT'_x) \parallel M) =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| D7FB8135 | D8EA45E8 | FB3C9059 | F146E263 | 0EF4BD51 | C4006A92 |
| EDB4C8B0 | 849963FB | | | | |

F.11.3.3, Per message data

Append the following text to clause F.11.3.3:

For the optimized variant of EC-SDSA,

$$R = h(\text{FE2BS}(IT_x) \parallel M) =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| 27D2F5B9 | 62A3ACF6 | 390A4718 | EA540DA7 | 9612A60E | AA15BEBB |
| 00B9E166 | 5783F7C7 | 91CCAC42 | 2CEE815A | 9C5DA367 | 8AC8D1F0 |

F.11.3.4, Signature

Append the following text to F.11.3.4:

For the optimized variant of EC-SDSA,

$$R =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| 27D2F5B9 | 62A3ACF6 | 390A4718 | EA540DA7 | 9612A60E | AA15BEBB |
| 00B9E166 | 5783F7C7 | 91CCAC42 | 2CEE815A | 9C5DA367 | 8AC8D1F0 |

$$S =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| 22CC89CE | B9E6BE84 | 15CC14B3 | 99BC66E6 | F3A21E5B | A38E09A6 |
| DE8DE670 | A145C0E4 | 74D5CC88 | BE8878F0 | 123CC662 | 25A1BA12 |

F.11.3.5, Verification

Append the following text to F.11.3.5:

For the optimized variant of EC-SDSA,

$$R' = h(\text{FE2BS}(\text{IT}'_x) \parallel M) =$$

| | | | | | |
|----------|----------|----------|----------|----------|----------|
| 27D2F5B9 | 62A3ACF6 | 390A4718 | EA540DA7 | 9612A60E | AA15BEBB |
| 00B9E166 | 5783F7C7 | 91CCAC42 | 2CEE815A | 9C5DA367 | 8AC8D1F0 |

Annex E

To provide a description of the reduced Tate pairing, which is used in examples of Annex F, insert the following clause at the end of Annex E:

E.4 The reduced Tate pairing

Let $l > 2$ be prime, and let P and Q be points on E with $[l]P = O$, the pairing $\langle P, Q \rangle$ can be computed in the following steps:

- choose some random point T on E , then
- compute $\langle P, Q \rangle = (d(P, Q - T) / d(P, -T))^{(p^k - 1) / l}$.

If during the computation of the pairing, a division by zero is attempted, then the computation should be restarted with a new point T .

NOTE 1 – More detailed information of pairing implementation can be found in [2, 14].

NOTE 2 – The reduced Tate pairing is used in numerical examples of clauses F.7 and F.8.

F.2.2, Signature key and verification key

To correct a one-digit error in the value of Y , in the last line of the formula for Y change “48CDF8DE” to “48CBF8DE”.

F.3, Pointcheval-Vaudenay mechanism

To identify which hash-function is used in the example of F.3, insert the following at the beginning of F.3:

This example uses the Secure Hash Algorithm (SHA-1) as the hash-function h . The hash-code is simply the value of SHA-1.

F.5.2.1, Parameters

To correct a one-digit error in the value of G_Y , in the formula of G_Y change “631011EC” to “631011ED”.

F.5.3, Example 2: Field F_p^m , 32-bit P and $m = 5$

To correct a minor typographical error in the title of F.5.3, change “Example 2” to “Example 3”.

F.7.2.1, Parameters

To correct the polynomial cited in F.7.2.1, change " $Y^2 = X^3 + 1$ " to " $Y^2 = X^3 + X$ ".

F.8.1.1, Parameters

To correct a multiple-digit error in the value of q in F.8.1.1, in the formula of q change

80000000 00000000 00FFFFFF FFFFFFFF FFFFFFFF

to

80000000 000FFFFFF FFFFFFFF FFFFFFFF FFFFFFFF.

Bibliography

Update reference [3] (ISO/IEC 11770-3) by changing "1999" to "2008".

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