

# INTERNATIONAL STANDARD

**IEC**  
**61162-410**

First edition  
2001-11

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**Maritime navigation and radiocommunication  
equipment and systems –  
Digital interfaces –**

**Part 410:  
Multiple talkers and multiple listeners –  
Ship systems interconnection –  
Transport profile requirements  
and basic transport profile**



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## Maritime navigation and radiocommunication equipment and systems – Digital interfaces –

### Part 410: Multiple talkers and multiple listeners – Ship systems interconnection – Transport profile requirements and basic transport profile

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## INTERNATIONAL ELECTROTECHNICAL COMMISSION

**MARITIME NAVIGATION AND RADIOCOMMUNICATION  
EQUIPMENT AND SYSTEMS –  
DIGITAL INTERFACES –**

**Part 410: Multiple talkers and multiple listeners –  
Ship systems interconnection –  
Transport profile requirements and basic transport profile**

## FOREWORD

- 1) The IEC (International Electrotechnical Commission) is a worldwide organization for standardization comprising all national electrotechnical committees (IEC National Committees). The object of the IEC is to promote international co-operation on all questions concerning standardization in the electrical and electronic fields. To this end and in addition to other activities, the IEC publishes International Standards. Their preparation is entrusted to technical committees; any IEC National Committee interested in the subject dealt with may participate in this preparatory work. International, governmental and non-governmental organizations liaising with the IEC also participate in this preparation. The IEC collaborates closely with the International Organization for Standardization (ISO) in accordance with conditions determined by agreement between the two organizations.
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International Standard IEC 61162-410 has been prepared by IEC technical committee 80: Maritime navigation and radiocommunication equipment and systems

The text of this standard is based on the following documents:

FDIS	Report on voting
80/311/FDIS	80/326/RVD

Full information on the voting for the approval of this standard can be found in the report on voting indicated in the above table.

This publication has been drafted in accordance with the ISO/IEC Directives, Part 3.

The special typographical conventions and nomenclature used in this standard are defined in IEC 61162-400 annex A.

Annexes A, B, C and D are for information only.

The committee has decided that the contents of this publication will remain unchanged until June 2005. At this date, the publication will be

- reconfirmed;
- withdrawn;
- replaced by a revised edition, or
- amended.

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## INTRODUCTION

International Standard IEC 61162 is a four-part standard which specifies four digital interfaces for applications in marine navigation, radiocommunication and system integration.

The four parts are:

- IEC 61162-1 Single talker and multiple listeners
- IEC 61162-2 Single talker and multiple listeners, high speed transmission
- IEC 61162-3 Multiple talkers and multiple listeners – Serial data instrument network
- IEC 61162-4 Multiple talkers and multiple listeners – Ship systems interconnection.

Part 4 of the standard is subdivided into a number of individual standards with part numbers in the 400 series.

This part of the standard contains the specification of the requirements to an IEC 61162-4 transport profile (T-profile) and also the specification of one implementation, based on redundant Ethernet and Internet protocol functionality. The T-profile is the protocol transport mechanisms that offer simple message or byte stream transport services to the higher protocol layers (defined in other parts of the standard). In addition, the T-profile also offers services for time distribution and physical network management.

The use of Internet and Ethernet protocols offer low cost and high efficiency data transport in any kind of system. However, for safety related applications, certain measures have to be taken to avoid that particulars of office-quality and off-the-shelf technology create safety risks. This part of the standard specifies mechanisms by which a certain degree of quality of service can be guaranteed from these networks, including the provision of redundancy.

Other T-profile documents will be prepared with specifications of the same T-profile requirements over other transport protocols. This will be issued in the same number series as this standard (IEC 61162-41x).

Relationship with the other parts of the IEC 61162 series of standards is defined in annex B to IEC 61162-400.

# MARITIME NAVIGATION AND RADIOCOMMUNICATION EQUIPMENT AND SYSTEMS – DIGITAL INTERFACES –

## Part 410: Multiple talkers and multiple listeners – Ship systems interconnection – Transport profile requirements and basic transport profile

### 1 Scope and object

This part of IEC 61162-4 defines the general requirements of the T-profile and three implementations of the T-profile over the Internet V4 (IPV4) protocol suite. Part 400 of this standard defines the relationship between the different protocol levels (T-profile, A-profile and companion standards) and part 401 defines the A-profile, the immediate user of the protocol level defined in this part.

The different components of the IEC 61162-4 standard are defined in IEC 61162-400. The T-profile is the specification of the communication services and the communication protocols used by the LNA to implement the A-profile functionality. Basically, the T-profile consists of the following components:

- a) a transport layer interface (TLI) definition that specifies the services and the semantics that will be available to the application level of the LNA (and in some cases the MAU). This includes data transport as well as time and network management services. The TLI will be general to all T-profiles and is defined in this part 410 of the standard,
- b) A T-profile protocol definition that specifies how the services provided by the TLI and additional time distribution and physical network management services are implemented on the protocol level. This part 410 contains a number of alternative T-profile protocol specifications using the Internet V4 series of standards. Additional parts of this standard will address other T-profiles based on other protocol families.

Note that the time distribution and network management functionality may or may not include specific TLI services. For some systems this functionality may be interfaced to directly by the underlying operating system. Note also that time distribution and network management are not strictly speaking transport related protocol functionality. However, the implementation of these services is normally dependent on the transport protocols in use and is, thus, placed in the T-profile part of the standard.

The purpose of this standard is to define and describe the services that will be provided at the transport level interface in a way which is completely independent of the underlying network environment as well as defining one possible implementation of these services over the Internet V4 protocols. The separation of service and protocol definitions allows the specification of several different transport profiles, each one dedicated to a specific network environment, and to use the same transport service interface in all cases.

Clause 4 defines the transport level services and clause 5 describes the transport layer interface through which the services are offered. These clauses define the general, network independent services.

Clause 6 defines the transport profile architecture for redundant Ethernet and Internet protocols version 4 (IPV4). Clause 7 defines the architecture for a local area non-redundant Internet network. These clauses define two specific implementations of the T-profile services.

Clause 8 defines a simple MAU-LNA protocol for use over wide area network (WAN) TCP/IP links. This can be used to implement a WAN architecture for the overall system. The WAN architecture is not intended for integrated ship control systems, but can be used for remote test integration and remote maintenance and diagnostics. The WAN protocol can also be used to support MAUs that are located in other host computers than the LNA, but on one local network (conformance class 4).

## 2 Normative references

The following normative documents contain provisions which, through reference in this text, constitute provisions of this part of IEC 61162. For dated references, subsequent amendments to, or revisions of, any of these publications do not apply. However, parties to agreements based on this part of IEC 61162 are encouraged to investigate the possibility of applying the most recent editions of the normative documents indicated below. For undated references, the latest edition of the normative document referred to applies. Members of IEC and ISO maintain registers of currently valid International Standards.

IEC 61162-400, *Maritime navigation and radiocommunication equipment and systems – Digital interfaces – Part 400: Multiple talkers and multiple listeners – Ship systems interconnection – Introduction and general principles*

ISO 8802-3, *Information technology – Telecommunications and information exchange between systems – Local and metropolitan area networks – Specific requirements – Part 3: Carrier sense multiple access with collision detection (CSMA/CD) access method and physical layer specifications*

ISO/IEC 9595: *Information technology – Open Systems Interconnection – Common management information service*

ISO/IEC 9596-1: *Information technology – Open Systems Interconnection – Common management information protocol – Part 1: Specification*

RFC 768:1980, *User Datagram Protocol (UDP), Internet Activities Board recommended standard*

RFC 793:1981, *Transmission Control Protocol (TCP), Internet Activities Board recommended standard*

RFC 826:1982, *Address Resolution Protocol (ARP), Internet Activities Board elective standard*

RFC 894:1984, *Internet Protocol on Ethernet Networks, Internet Activities Board elective standard*

RFC 1157:1990, *Simple Network Management Protocol (SNMP)*

RFC 1189:1990, *Common Management Information Services and Protocols for the Internet (CMOT and CMIP)*

RFC 1213:1991, *Management Information Base for Network Management of TCP/IP-based Internets: MIB-II*

RFC 1305:1992, *Network Time Protocol, Version 3 – Specification and Implementation*

RFC 2030:1996, *Simple Network Time Protocol (SNTP), Version 4 for IPv4, IPv6 and OSI*

RFC 2500:1999, *Internet Official Protocol Standards – Internet Activities Board standard*

NOTE RFC (request for comments) is a document, issued by the Internet engineering task force (IETF) the International standardization body for the Internet, that describes a part of the Internet protocol. Some RFCs are accepted as official Internet standards and listed in the "Internet Official Protocol Standards" itself an RFC.

### 3 Definitions

For the purpose of this part of IEC 61162, the definitions in IEC 61162-400 and the following definitions apply:

#### 3.1

##### **broadcast**

see multicast

#### 3.2

##### **CL – connectionless**

a connectionless communication means that sender and receiver do not have to know each other. There is no association established between sender(s) and receiver(s) before messages are sent (see also peer-to-peer and client-server)

#### 3.3

##### **client-server**

a client-server communication link is established by the client after a server has allowed the connection attempt by the establishment of a listening connection point. The client is the active part while the server allows the connection (see also connectionless and peer-to-peer)

#### 3.4

##### **CMIP – common management information protocol**

see ISO/IEC 9595 and ISO/IEC 9596-1

#### 3.5

##### **CMIS – common management information services**

see ISO/IEC 9595 and ISO/IEC 9596-1

#### 3.6

##### **CMOT – CMIS/CMIP over TCP/IP**

an Internet proposed standard protocol. Its status is elective (see RFC 1189)

#### 3.7

##### **CO – connection-oriented**

the opposite of CL (connectionless). A data exchange where an association between sender and receiver exists

#### 3.8

##### **connection point**

an entity that can represent a communication link end point (for established connections) or a connection attempt between two host computers in some state. It is also used for connectionless communication, but in this case it represents just the local host computer's port to the network

#### 3.9

##### **Ethernet**

refers to a carrier sense multiple access collision detect (CSMA/CD) local area network protocol standard as defined in ISO/IEC 8802-3. The medium access control (MAC) frame format shall use the Internet protocol type (0800) in the length/protocol field (see RFC 894).

Any type of Ethernet can be used in systems compliant with this standard as long as they fulfil relevant technical requirements and the system integrator ensures compatibility between the integrated components. The most relevant technologies are:

- 10Base-5 – thick coaxial, shared media bus – can be used in environments that require noise immunity;
- 10Base-2 – thin (RG58) coaxial, shared media bus – can be used in environments with low noise immunity requirements;
- 10Base-T – shielded or unshielded twisted pair – used in conjunction with a repeating or switching hub;
- 10Base-F – fibre optic media, as above, but for very high noise immunity applications;
- 100Base-TX – shielded or unshielded twisted pair, as above. Note that this medium requires special precautions to extend network lengths beyond 200 m;
- 100Base-FX – fibre optic, as previous, but for very high noise immunity.

The recommended solution is the use of non-duplex switching hubs with 10 and/or 100 megabit links to the host computers.

The system integrator must make sure that the relevant physical characteristics of the selected network solution satisfies the requirements defined in ISO/IEC 8802-3 and other applicable standards

### 3.10

#### **ICMP – Internet control message protocol**

an integral part of the Internet protocols (see 3.12)

### 3.11

#### **IGMP – Internet group management protocol**

a protocol used between hosts and multicast routers on a single physical network to establish hosts' membership in particular multicast groups. Multicast routers use this information, in conjunction with a multicast routing protocol, to support IP multicast forwarding across the Internet. The IGMP standard is part of the Internet RFC system, but is currently not used in this standard

### 3.12

#### **IP – Internet protocol.**

refers to RFC 2500. All required parts as defined by that document shall be implemented in an Internet compliant protocol stack. It is also required that the system in question implements the ARP (see RFC 826) address resolution protocol. Optionally, other address resolution schemes can be used at the discretion of the system integrator.

RFC 2500 will give additional information as to which additional standards apply to conforming implementations of the Internet protocol.

RFC 1918 lists a set of reserved private address spaces that can be used for a ship-board control network that must not be connected to off-ship or other ship-board Internets. For the purpose of this standard, it is recommended that the following network addresses are used:

172.16.0.1 to 172.16.255.254 (network mask 255.255.0.0 – class B)

### 3.13

#### **IPV4 – Internet protocol version 4**

the version used in this issue of the standard (see previous subclause)

### 3.14

#### **IPV6 – Internet protocol version 6**

the proposed next generation Internet protocol. It is currently not sufficiently accepted to be viable as the basis of this standard, but future generations of this standard are expected to migrate to IPV6

### 3.15

#### **LAN – local area network**

a network within a delimited physical area. Typical for control system networks (see also WAN, MAN)

### 3.16

#### **LNAC – LNA communication module interface**

a variant of TLI that is dedicated to MAU-LNA communication

### 3.17

#### **loop-back interface**

refers to the loop-back function in the Internet protocol whereby an Internet message can be sent between two entities on the same network host. The complete class A network number 127 is assigned the “loop-back” function (see RFC1060). The address 127.0.0.1 will be used for the loop-back

### 3.18

#### **MA – management agent**

an SNMP agent located in a network element (NE)

### 3.19

#### **MAN – metropolitan area network**

a network between LAN and WAN. The area the network covers is more than local, but not worldwide

### 3.20

#### **MIB – management information base**

basis for physical network management protocol

[RFC 1213]

### 3.21

#### **MTU – maximum transmission unit**

the largest message that can be transmitted over a given T-profile without fragmentation. For Ethernet, the MTU is 1 500 octets

### 3.22

#### **multicast**

a transport mechanism by which any number of computer hosts can be reached with one transmitted message. This transport mechanism is referred to as multicast or broadcast. The T-profile specification defines the actual physical mechanism that is used to implement this service

### 3.23

#### **NE – network element**

term used within SNMP to identify a managed entity in a network. This is typically a host computer with its protocol entities, a router, a gateway or a switch

### 3.24

#### **NMA – network management applications**

applications within a NMS

### 3.25

#### **NMS – network management station**

a host computer, with software implementing one or more NMA, that is responsible for physical network monitoring (see MIB and SNMP)

### 3.26

#### **NNN – network node number**

an identification of one node on a TP network. This is typically the Internet host number (excluding the network part of the address) for IPV4 networks. The NNN shall be unique and unambiguous for one node, even if this node is on a redundant network with two different network addresses. The NNN is a 32-bit unsigned integer

**3.27****NTP – network time protocol**

all references in this standard address the Internet time distribution and synchronization protocol described in RFC 1305

**3.28****OSI – open systems interconnect**

this term is defined in IEC 61162-400

**3.29****peer-to-peer**

a two-way communication link that can be established by both sides concurrently. The T-profile shall make sure that only one link is established and that the link has the same identification on both sides (see client-server and connectionless)

**3.30****primary network**

a redundant network has two independent transport paths between any two nodes (primary and secondary network). The primary network is the transport path defined as the one that is normally used when connections are established between the nodes, and that normally is to be used if non-redundant nodes are attached to the network. However, the T-profile will be able to use the secondary network in the same manner as the primary (i.e. accept incoming connects that arrive on the secondary network alone). The T-profile may process request on the secondary network after requests on the primary

**3.31****priority**

assigns importance to messages or message streams transported through the communication network. Three levels are defined: urgent, normal and low.

It must (as a minimum) act as a rudimentary bandwidth allocation scheme where high priority (importance) data is serviced before lower priority data

**3.32****QoS – quality of service**

degree to which a certain protocol (typically T-profile) supplies certain services to the higher levels

[IEC 61162-400]

**3.33****redundant IPV4**

reference to a redundant Internet means a T-profile as defined in clause 6 of this standard. Where the standard redundant IPV4 network is used (LNADR\_IPV4R), the following addresses are used:

172.16.0.1 to 172.16.255.254 (network mask 255.255.0.0 – “primary”).

172.17.0.1 to 172.17.255.254 (network mask 255.255.0.0 – “secondary”).

Other addresses can be used when the system integrator defines new T-profile networks for use in a specific implementation (LNADR\_IPV4RE).

Both networks will be of the same class (class B for the standard network) and each network node will have the same node address on the primary and secondary networks.

For all IP networks the following port numbers will be used:

**Table 1 – Internet IPV4 protocol port numbers**

Port type	Protocol	Port
LNA-LNA listening port (normal)	TCP	23305
LNA-LNA listening port (urgent)	TCP	23306
LNA-MAU listening port (normal)	TCP	23303
LNA-MAU listening port (urgent)	TCP	23304
MAU-MAU stream listening ports	TCP	22307-
ABC0 (standard broadcast)	UDP	23307 (23308)
ABC1	UDP	23309 (23310)
ABC2	UDP	23311 (23312)
ABC3	UDP	23313 (23314)
ABC4	UDP	23315 (23316)
ABC5	UDP	23317 (23318)

NOTE 1 The port numbers are different from the ones used in MiTS. This is done to support the co-existence of MiTS and IEC 61162-4 networks.

NOTE 2 The stream ports will be selected by trying different port numbers in sequence, starting with the one specified.

Two port numbers are listed for each UDP broadcast port. The first number (odd number) is the mandatory listening (destination) port. The even number in parenthesis is the recommended sending port (source) port. Other sending port numbers can be used if desired in an implementation.

NOTE Some implementations of UDP can cause problems when the same port number is used for sending and receiving. A typical problem is that sent broadcast messages are not received by the sending computer itself.

Likewise, two port numbers are used for the message ports to support differentiation between normal and urgent messages. Implementations that do not support differentiation will use the normal port (valid for MAU-LNA link only).

The same ports will be used for all IP type networks and on both the primary and secondary side of a redundant network.

**3.34**

**secondary network**

represents the second transport path

See primary network (3.30)

**3.35**

**SMI**

structure and identification of management information is a model for description of the structure of needed information

**3.36**

**SNMP – simple network management protocol**

Internet protocol used for physical network management

[RFC 1157]

**3.37**

**TCP/IP – transmission control protocol/Internet protocol**

refers to the reliable and stream oriented protocol defined by RFC 793. A system implementing TCP/IP will also implement the Internet protocol as described above

### 3.38

#### **TPN – TP network**

definition of a physical network architecture with all relevant address and network parameters. IEC 61162-410 defines a number of different TP networks based on the IPV4 family of protocols

### 3.39

#### **UDP – user datagram protocol**

the unreliable message based protocol defined by RFC 768. A system implementing UDP shall also implement the Internet protocol as described above

### 3.40

#### **WAN – wide area network**

a network that, in principle, can span the entire world. A typical example is the Internet

## **4 Transport layer service specification**

### **4.1 General**

The transport layer will provide the upper layers (the MAUs and the A-profile) with functions for the following services:

- communication services;
- time services;
- network management services;
- graceful degradation facility.

Communication and graceful degradation services are provided by the transport layer while time and management services typically involve the transport layer and/or other entities (upper layers, resources depending on actual implementation).

### **4.2 T-profile network and host computer addressing**

#### **4.2.1 Introduction**

This subclause defines requirements for the T-profile network and host computer addressing. The described addressing scheme shall ensure easy portability for higher level entities between different T-profiles.

#### **4.2.2 Addressing scheme**

Host computers that communicate via a specific T-profile use a two-level addressing scheme:

- each T-profile shall as a minimum, define one T-profile network (TP network) that specifies the physical architecture of the network. This network is logically a linear network without any further subdivisions or segmentation. Some T-profiles may define more than one TP network to subdivide the different functions supplied by the T-profile into separate modules.

NOTE A typical example is a redundant Internet profile that subdivides the TP network into one for MAUs (supporting only stream communication) and one for the LNA (supporting message-based communication).

- each host computer has a unique and singular network node number (NNN). This number when referenced to a specific TP network shall be sufficient to address the host computer by the T-profile implementation.

### 4.2.3 Unique and singular hosts on the network

The T-profile shall provide an addressing mechanism that makes each host computer on a T-profile network uniquely and singularly identifiable by a network node number (NNN). This applies for all types of T-profile networks, for those that are non-redundant as well as redundant. One NNN shall map to one and only one possible host computer.

NOTE For redundant networks the single NNN must normally be convertible to two physical network addresses, one for each of the active links.

### 4.2.4 Non-segmented logical network

The addressing scheme defined here creates a non-segmented logical network on which the host computers reside. However, a T-profile network may optionally provide physical segmentation between groups of host computers in the NNN address space. This standard does not specify how this can or should be done. This standard requires that all nodes on the T-profile network, independent of physical segmentation, shall be able to establish communication links between themselves and exchange data over these links.

NOTE A T-profile for the wide area Internet may need to handle broadcasts specially (e.g. by using a multicast protocol, for example IGMP) to enable general MAU name look-up over several networks.

### 4.2.5 Size of network

The NNN is an unsigned integer of at most 32 bits in length. Physical addressing schemes that require more bit-space must define attributes in the T-profile network description that supply the extra information so that a shorter NNN can be mapped to the correct address.

NOTE IPV6 uses a 128 bit long network address. To make host computers addressable by 32 bits, the remaining 96 bits must be determined from the T-profile network attributes. The same applies to redundant networks, for example for IPV4, where each host computer typically has two network addresses.

## 4.3 Communication services requirements

### 4.3.1 Introduction

This subclause defines the general attributes of the communication services that shall be made available through the T-profile.

### 4.3.2 Connectivity and data transfer internally and externally

The T-profile shall provide the following general services:

- the transfer of data between any nodes attached to the embedded system level network;
- allow connection to one or more other networks than the embedded system level network.

NOTE The concept of defining different logical T-profile networks allows several T-profiles to coexist on the same physical network. In this case, the *other* network may be another T-profile. More commonly, the other network is some form of administrative network accessed via a separate network interface card (NIC).

Different service classes are defined in this standard for data transfer on the embedded system network. Data transfer out of the embedded system network is not covered by this standard.

### 4.3.3 Service classes

#### 4.3.3.1 Overview

This subclause defines the service classes that shall be supported. Table 2 summarizes the service classes and their main characteristics.

**Table 2 – Service classes**

Communication	Connection	Data	Priority	Relation	Instances
LNA-LNA data	CO	Msg	U and N	Peer-to-peer	One
LNA-LNA MC in	CL	Short msg	U and / N	Multicast listen	Six MC ports
LNA-LNA MC out	CL	Short msg	U and / N	Multicast send	Six MC ports
MAU-MAU stream N	CO	Stream	N	Client/server	Any
MAU-MAU stream L	CO	Stream	L	Client/server	Any
MAU-LNA data	CO	Msg	U and N	Client/server	One

NOTE The first column specifies the service class and by that, its normal area of use.

The second column specifies whether the communication uses a connection-oriented (CO) or connectionless service.

The third column specifies whether data is sent as messages (with message boundaries) or as streams.

The fourth column specifies the priority classes that shall be supported for this service class.

The fifth column specifies the relationship between host computers using the service.

The final column specifies the various instances that are supported for each communication class. For LNA-LNA and MAU-LNA links, it is just one: The one standard path to the one listening LNA. For multicast services, the A-profile requires six predefined addresses, one for name look up and five for anonymous broadcast. For MAU-MAU link the TP network should support a reasonable number of CP to allow possible multiple MAUs on one host computer to have multiple connections to other MAUs.

The communication service classes are:

- LNA-LNA data: this is a communication link providing reliable bi-directional data transfer between any two LNAs (communication initiated by any party – peer-to-peer);
- LNA-LNA MC: this is a broad- or multicast service that provides name look-up, watchdog and anonymous broadcast services between LNAs. This is also essentially a peer-to-peer communication class, although connectionless. The service is divided into a transmitting and a receiving component;
- MAU-MAU stream: this is a service providing bulk data transfer between MAUs. The communication is initiated by a client connecting to a waiting server. There are two stream classes, one for normal and one for low priority (priority is associated with the link rather than with the data transported). The service is bi-directional;
- MAU-LNA data: this is a service (defined in the A-profile) providing communication between a MAU and its LNA. This is also a client-server type communication. The service is bi-directional.

#### 4.3.3.2 Implementation of multiple service classes

One implementation of a T-profile may provide one or more of the above service classes, dependent on the required use. The minimum requirements are listed in the table below.

**Table 3 – T-profile service implementation**

TP implementation for	Required service classes
LNA-LNA	LNA-LNA data, LNA-LNA MC
MAU-MAU	MAU-MAU stream N, MAU-MAU stream L
MAU-LNA	MAU-LNA data

NOTE Although all these services may be implemented by one T-profile, the implementation may offer the services to the users as a different TP network. Typically, the three groups of services will be offered as three distinct TP networks.

#### 4.3.3.3 Message or stream type communication

The following communication services shall be provided to the upper layers:

- Msg: reliable message transfer. The transfer mode will be a unicast (point to point) transfer mode. Message length may be limited by T-profile implementation. In particular, urgent messages may be limited to the MTU of the T-profile;
- Stream: reliable stream transfer of a unicast (point to point) transfer mode. Data is transmitted between sender and receiver without regard to message boundaries;
- Short msg: unreliable short message transfer. The transfer mode will be a multicast (point to multipoint) or broadcast transfer mode. The message length is typically limited by the MTU of the T-profile (and, hence, is “short”).

#### 4.3.3.4 Connection modes and flow control

A reliable transfer will be achieved through a connection-oriented (CO) mode protocol while an unreliable transfer will be achieved through a connectionless (CL) mode protocol.

A CO mode protocol provides for full duplex reliable (with guarantee of delivery to the destination), sequential and without duplication end to end data transfers by using sequence numbers, acknowledgements to indicate correct delivery of the data, and retransmissions in case of error (data loss, data damage).

A flow control mechanism shall be provided for CO services to avoid a receiver being congested by too many arriving messages.

A mechanism should be provided to detect that either side of an established connection has ceased to function (remote network node damage). This mechanism will in any case be augmented by a similar service in the A-profile so it will not be of critical importance.

A CL mode protocol provides for unreliable (without guarantee of delivery to the destination(s)) end to end message type data transfers. No notification is generated for either loss, duplication or out of sequence delivery of data. Damaged data need not be retransmitted, but shall in no case be delivered to the receiver (i.e. discard messages with bad checksums). The service is specified to be unidirectional (i.e. a listening service must be used together with a sending service).

Unreliable messages shall be short messages. It will be the responsibility of the upper layers to ensure that the length of the messages provided at the transport layer is within the permitted maximum value (medium MTU).

Reliable streams / long messages shall be segmented and reassembled at the transport layer when necessary. Short messages (typically of urgent priority) should, if possible, be transmitted in one physical layer message.

#### 4.3.3.5 Priority

In order to support the time dependency of the transmitted data (time critical or not) and to allow the intermixing of data flows of different significance on the network, importance levels shall be defined as follows:

- urgent;
- normal;
- low.

The implementation of these priorities will be T-profile dependent. In particular, there may not be a well-defined relationship between priority levels for CL and CO services as the CL and CO mode transfers may use different protocols. A T-profile may also in some cases handle different priority classes as the same (e.g. not distinguish between urgent and normal multicast priorities) due to limitations in underlying protocol. The T-profile shall in any case provide a best effort implementation of priorities.

The system integrator must verify that the implemented priority mechanism is suitable for the use to which the T-profile is being put.

#### **4.3.3.6 Logical communication paths and node relationships**

A logical communication path results from the association between two (point to point communication) or more (point to multipoint communication) transport level end points.

Dependent of communication class, different principles apply for the establishment of logical paths:

- peer-to-peer: communication between two LNAs can be established by any part at any time. The T-profile implementation must make sure that concurrent connection establishments result in only one logical path;

NOTE To facilitate a similar handling of all types of communication links, also the peer-to-peer service will be established in the same way as a client/server connection.

- client-server: communication between LNA and MAU and between MAUs is established by one of the parties acting as client (doing the connection request) and the other acting as server (waiting for connection attempt). This always results in one unambiguous communication path;
- multicast: this communication is connectionless and does not use logical paths.

### **4.4 Time services**

#### **4.4.1 General**

A common time reference shall be available in the system level network in order to timestamp all time dependent data transmitted by the nodes and to estimate the validity of the received data.

The time shall be generated and maintained at the network level at least at one dedicated time generator/server. The time distribution service shall allow the use of multiple (redundant) servers.

A time distribution service shall be provided to update and to synchronize the local clock available at each attached node to the global clock available at the time generator/server.

In addition, the time shall be maintained at each node in order to make available to their local user(s) a reference time even in case of failure of the dedicated common reference time generator/server(s).

The T-profile specification shall minimize synchronization jitters by, for example using urgent priorities for time-related messages. Other means can also be employed. The worst case jitters shall be specified.

#### **4.4.2 Services**

The time service shall maintain a synchronized time over the complete T-profile network. No particular services are provided to the applications. It is assumed that the time related services are made available through the operating system interface.

## 4.5 Management services

### 4.5.1 General

The management services shall enable monitoring and control of the system level network operation.

It requires monitoring and control of the different network elements: attached nodes and intelligent network equipment (e.g. switch, router or gateway).

Each one of these network elements shall maintain management information providing information (e.g. state, thresholds, counters, action results) on the local resources (software and hardware).

The management information involves standard management information (information dedicated to standard communication protocols), proprietary management information (information dedicated to manufacturer's equipment) and if necessary specific management information (information dedicated to specific system requirements).

Network configuration, fault and performance management functions are achieved from a dedicated network management station (NMS) providing the network maintainer with the current status of the network as an aggregate of the different management information distributed in the network.

The NMS accesses to the management information, local to each network element, through a standardized network management protocol, generally at the OSI application layer, while local management functions on a node rely on specific management services depending on the node implementation (operating system and protocol stack features).

### 4.5.2 Services

The following management services shall be provided:

- monitoring of a node attached to the network (load and fault conditions);
- network monitoring (load and fault conditions).

The following services are optional, but may be prohibited due to safety considerations:

- remedial actions;
- (re)configuration.

Each node shall detect internal errors and make available detailed status information to the local upper layers and to the remote NMS.

The NMS shall build up the overall network status (list of network elements in operation) by:

- receiving reports of changes from the nodes and intelligent network equipment (e.g. activation, deactivation and fault reports);
- requesting, with certain frequency, status, events, action results and statistical information from the same elements (if the information is not automatically broadcast).

The NMS shall monitor the network traffic and build traffic statistics.

Local node management shall be sufficient to enable a user to monitor the condition of its communication to the network and take remedial actions such as configuration change, network disconnection or test initialization. Any such actions shall only be available when the system is in a safe state (e.g. ship at berth) and in a way that cannot endanger the ship's safety.

The information returned from the network monitoring activity shall be sufficient to enable remedial actions to be taken such as network equipment configuration change, disconnection or reset.

The NMS must normally be duplicated for safety purposes.

#### **4.6 Graceful degradation facility**

The system level network shall be designed with a capacity for graceful degradation in order to retain maximum performance from the remaining available resources in case of a failure.

Redundancy should be used for safety critical parts of the system in order to provide services and performances even after a single failure.

In case of a number of faults exceeding its resilience, the network shall still provide its services to the maximum possible number of nodes considering that time budget performances may not be guaranteed and that a part of the network may not be available. The network is then in a degraded condition; the gravity of the degradation depends on the types of faults.

When operation of a resource is degraded or fails, an indication shall be provided to the NMS and to the A-profile where possible and relevant (see previous subclause).

Graceful degradation shall be automatic and transparent to the users at least on the occurrence of these events:

- failure of a single cable or single network interface hardware (no loss of functionality);
- fault, insertion or shutdown of one or more nodes (loss of function of that node). The node may have to be duplicated for safety purposes;
- loss of all the network equipment within a single compartment (due to fire or flood – loss of functionality located in the compartment).

The implementation of graceful degradation is partly dependent on physical architecture.

### **5 Transport layer interface (TLI)**

#### **5.1 Overview**

The A-profile shall get access to the transport layer services through a TLI. The TLI shall make the actual transport profile transparent to the upper layers.

The TLI shall provide the upper layers with a set of standard transport service primitives independent of the network environment while its implementation will depend on the actual environment (e.g. transport profile, operating system and local resources). The TLI will normally not supply time and network management services.

The TLI is illustrated in the following figure.

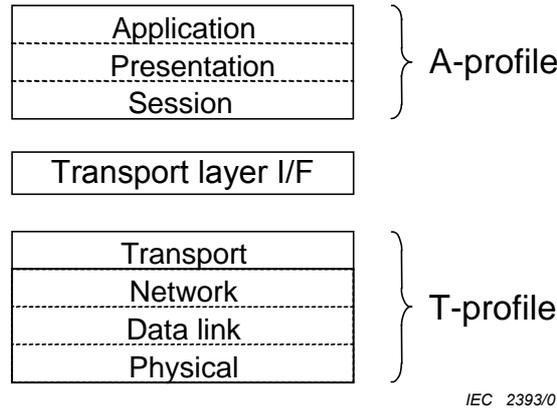


Figure 1 – Transport layer interface

The purpose of the TLI is to define a standard interface between the A-profile and the T-profile. The A-profile specification (IEC 61162-401) depends on the TLI service descriptions contained in this clause.

The service specification will not require any particular language binding or any particular modularization of the services into specific subroutine calls. This is left to the implementation. However, the service specification will be made in such a way so that it can be translated one to one into a normal procedural language.

Annex A describes a typical software library structure for the TLI and associated modules.

### 5.1.1 LNAC and TLI

A special variant of the TLI is the LNA communication module (LNAC). This module is used to set up communication between LNA and MAU. For MAU-LNA communication over a normal network communication path (e.g. TCP/IP) it is normally most convenient to implement the LNAC as a normal TLI with the same services and a specialized TP network. For IPC type communication between LNA and MAU, it may be more convenient to implement the LNAC in some other way.

This standard defines requirements to the TLI and associated T-profiles. These requirements can optionally be used also for the LNAC. An Internet LNAC profile is defined in clause 8.

## 5.2 General principles

### 5.2.1 TLI, TP network service classes and connection points

To have a general communication mechanism, the concept of a connection point is used. The connection point is a communication link end point and is directly associated with the TLI. However, the connection point derives address information and capabilities from the TP network and the service class it belongs to. The relationship between the TLI, TP networks, service classes, instances and connection points are illustrated in the ER diagram below.

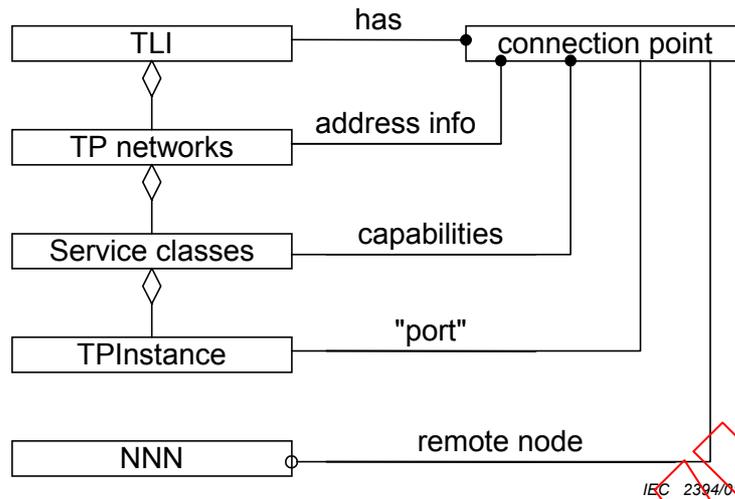


Figure 2 – Relationship between TLI, TP networks and ports

One TLI shall support the use of a number of different TP networks (minimum one, usually two or more, for example the LNA needs one network for the LNA-LNA communication and will typically use one or two more TP networks for LNA-MAU connections). Each TP network supports one or more service classes, dependent on the network.

NOTE Even though one TP network as specified in this standard can be able to support several service classes, it is often useful for the implementation to segment the one TP network into a set of different TP networks – normally based on the service classes supported. In many cases, it is easier to handle TP networks that can only derive one or a few types of CPs, for example with respect to how required CP parameters are specified (see table 4).

For each service class, it will normally be possible to create a number of communication service instances. The services may be just different instances of the service (e.g. concurrent MAU-MAU streams between two host computers) or some meaning may be pre-allocated to the instances (e.g. a pre-defined multicast group). One CP can be created for each instance. Finally, each instance of a connection point will normally also have a remote node associated with it. This is, however, not the case for connectionless services.

5.2.2 Connection point capabilities

The capabilities of a connection point (CP) are derived from one of the service classes supported by the network as defined in table 2. However, to support client-server type link establishment, it is necessary to distinguish between listening and connecting CPs. Thus, the complete list of CP types is as follows:

Table 4 – Connection point types

CP type	Connection	Priority	Type	Parameters
CPL_LNAMSG	CO-server	U and N	Listen for incoming LNA-LNA	None
CP_LNAMSG	CO-client	U and N	End point LNA-LNA	Remote NNN
CP_LNAMC_IN	CL	U and N	Short message input	ABC port
CP_LNAMC_OUT	CL	U and N	Short message output	ABC port
CPL_STREAM	CO-server	L or N	Listen for incoming stream	Port and priority
CP_STREAM	CO-client	L or N	End point reliable stream	Remote NNN, port, and priority
CPL_MAULNA	CO-server	U and N	Listen for incoming MAU-LNA	None
CP_MAULNA	CO-client	U and N	End point MAU-LNA	Remote NNN

The LNA-LNA connection point is defined as a peer-to-peer connection, but this standard requires that it is implemented as a client/server port by the TLI (i.e. have a listening variant). The difference from normal client/server links, for example stream links, is that any concurrent connection attempts from the two peers automatically are resolved into one CP at each side.

NOTE 1 The CO-server type CP must normally be able to create a new CP for each incoming connection. The way this is done is implementation dependent, but it is recommended that a fourth connection-type of CP is defined, e.g. CO-connect to represent these derived CPs.

NOTE 2 In the implementation of the T-profile, the CO-client and the CO-connect (see previous note) CP data structures must be co-ordinated so that only one of potentially two requested CP structures (listening and connecting) is returned to the application as valid. This is normally most conveniently done if it is the CO-connect CP that is returned.

The connectionless CP has one listening and one sending variant – this is to simplify the handling of multicasts where the sender is concurrently listening in on the same multicast group to which it is transmitting. Outgoing messages shall be copied to the incoming local CP (if open) by the T-profile.

Message connection points accept incoming and outgoing messages of any priority. Stream connection points have an inherent priority that is established when the port is created.

#### **5.2.2.1 Client-server connection establishment**

A client-server CP requires that one host computer acts as a server and another as client. The client (the connecting) host computer will not be allowed to establish a connection before the server has established a listening port. The server will only establish a connection as a result of the client doing a connection attempt.

The T-profile shall accept that two “identical” CP pairs are established when the two host computers concurrently establish a listening and a connection port each.

NOTE This is contrary to what is required of the peer-to-peer CP. The rationale is that client-server links are established as a result of a temporary need to exchange data (bulk data between MAUs). The situation outlined above can occur if both MAUs want to transmit data concurrently.

#### **5.2.2.2 Peer-to-peer connection establishment**

Two peer-to-peer CP pairs can be requested concurrently when two host computers want to communicate. The T-profile shall make sure that only one CP pair results from this process. In this standard, only the LNA-LNA connection is of this type.

NOTE This situation is typical of system start-up where many LNAs establish connections to each other. As there is no client-server relationship between LNAs, this requirement makes sure that each remote LNA is associated with exactly one CP.

To facilitate homogenous processing of connection points in the A-profile, the T-profile (and TLI) shall provide a listening as well as a connecting variant of the peer-to-peer CP. No incoming connection attempts shall be served before the listening CP has been established.

NOTE This requirement simplifies the use of protocol software that explicitly use a client-server model, for example sockets for TCP/IP. By having this requirement, the T-profile will get an explicit instruction to create a listening port when the LNA is ready to process incoming connection attempts.

#### **5.2.2.3 Connectionless connection points**

Connectionless communication (multicast) shall also use connection points. Although there is no connection between host computers in this case, the local T-profile shall establish a connection point to act as reference for incoming and outgoing messages.

Each connection point shall only be able to handle either incoming or outgoing messages. Two CPs are necessary to both send and receive multicasts.

NOTE If the T-profile uses multicast instead of broadcast, it must also automatically establish the multicast groups based on parameters in the TP network object and probably when the TP network is created. The user (LNA) will only request the establishment of a multicast CP and assume that it will work as specified.

### 5.2.3 Attributes of TP network

Each TLI, TP network will have a number of attributes that characterize its functionality. The minimum set of attributes is defined in the table below.

**Table 5 – TPN attributes**

Attribute code	Type	Description
TPA_NETWORK	int	Code (type) of TP-network
TPA_SERVICES	word32_m	Bitmap for available services
TPA_QOS	word32_m	Bit-map defining quality attributes
TPA_MTU	int	MTU of T-profile (zero if no limit)
TPA_STATE	int	State of TP (fully/reduced) operation
TPA_ADDR	address_m	Network address
TPA_LNNN	word32_m	Local network node number

#### 5.2.3.1 TP network address formats

The TP network address format will be dependent on the T-profile in use. The TLI shall adopt a T-profile independent address format as specified below.

- One `word16_m` entity specifying the type of network address (e.g. Internet IPV4, redundant Internet IPV4/Ethernet, redundant IPV6/Ethernet, redundant ARCNET or other). Table 6 and details in each T-profile definition will determine the identifier.
- One `word8_m` entity giving the length of the following address field.
- A sequence of maximum 48 `word8_m` entities containing the actual network address as required by the T-profile in use.

This is a record with the following format string (IEC 61162-401): “(w16[w8:48]w8)”.

#### 5.2.3.2 Defined TP network code ranges

The following table lists the TP networks that are defined by this part of IEC 61162 with required TP network and NNN address formats.

**Table 6 – TP network and NNN types**

Code	Values	Description	Notes
ANYADR_ANY	0	Undefined	For user specified addresses
MAUADR_IPC	1	IPC type MAU-LNA link	Not defined in this standard Reserved for LNAC
MAUADR_TCP	2	IPV4 TCP MAU-LNA link.	See clause 8.
MAUADR_TCPV6	3	IPV6 TCP MAU-LNA link.	Reserved for a future standard
LNADR_ARCNET	10 to 15	ARCNET range addresses	Reserved for a future standard.
LNADR_IPV4	120 to 139	IPV4 address formats.	See clauses 6 and 7.
LNADR_IPV6	140 to 159	IPV6 address formats.	Reserved for a future standard.

NOTE Even with the same TP network code, different address formats can be supported, for example for the user defined network and for the IPC network. In these cases, the address field must contain enough information to distinguish between the various address forms.

More address codes can be defined in future parts of this standard.

**5.2.3.3 Compatibility between redundant and non-redundant TP networks**

This standard allows the T-profile to define compatible redundant and non-redundant TP networks. In these cases, hosts on a redundant TP network shall be able to connect to hosts on the corresponding, non-redundant TP network and vice versa.

NOTE This is useful in a scenario where some non-critical and low-cost host computers shall be interfaced to the safety-related system network for, for example data acquisition or periodic maintenance.

It is allowed to let one (basically redundant) TP network accept both types of incoming connections. The T-profile is required to determine what kind of network the remote connection host computer resides on and report this to the local application.

NOTE An alternative is to let the local host computer define a TP network for both types of hosts that it will connect to or allow connections from. Typically, the local host computer will define one redundant and one non-redundant TP network.

**5.2.4 Attributes of connection point**

Each connection point will have a number of attributes that partly are set permanently during the creation of the CP and partly represent the state of the port. All attributes should be retrievable through the TLI calls. The minimum set of attributes is listed in the table below.

**Table 7 – CP attributes**

Attribute code	Type	Description
CPA_NETWORK	int	Identity of TP network
CPA_SERVICES	int	Code for selected service
CPA_PRIO	word32_m	Bit-map showing priorities available
CPA_STATE	int	State of CP
CPA_RNNN	word32_m	Remote network node number
CPA_ADDR	address_m	Local address (including, for example port number)

In addition it will maintain a dynamic state vector keeping information on queue status and traffic statistics data.

**5.2.5 Connection point states**

The following state diagrams use the same set of states:

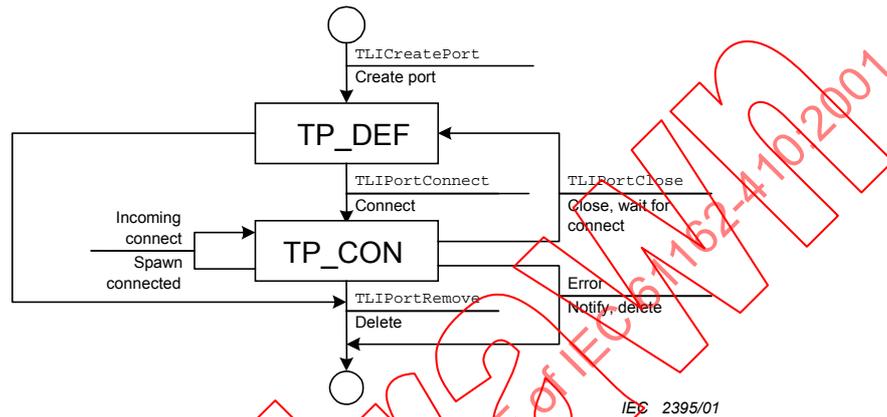
- a) TP\_DEF: the CP attributes are defined and made known to the TLI. The TLI is waiting for a connection request, either from local or remote host computer;
- b) TP\_REQ: the application has requested that the CP shall be made ready for communication. This state is used for a connecting CP that needs to establish a connection with another host computer's listening CP;
- c) TP\_CON: this state is used for CP that are in a communicable state, i.e. messages can be sent and received or connecting CP can be accepted;
- d) TP\_ERROR: errors, except time-out and forced close, make a CP go to this state. The CP will have to be reset (removed) before it can be used anew;
- e) TP\_CLOSE: a close request will normally make the CP go into this state (except for ports in the TP\_REQ state). This state allows queued messages to be sent before closing the connection.

NOTE 1 Common for most state transition diagrams is that they contain an explicit created and error state. These states are used so that all external events (close, error) cause a transition to a known state in a still defined object. The object can only be deleted through an explicit call from the higher level program. The exception is the CP that is derived from a listening CP (the "spawned" CP).

NOTE 2 The existence of the created state also makes it convenient to let the close call move the CP to created state instead of deleting it. This is not, strictly speaking, necessary but makes the diagrams more consistent.

**5.2.5.1 Listening connection points**

The following figure defines the states of a listening connection point.

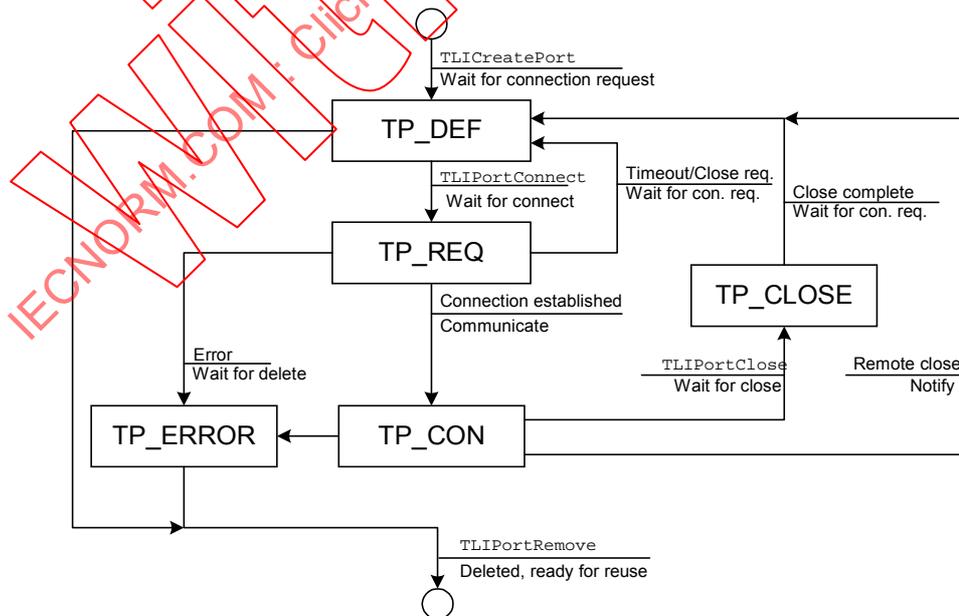


NOTE Note that this connection point is almost stateless and that it "spawns" new connection points during connection attempts.

**Figure 3 – Listening connection point states**

**5.2.5.2 Connecting connection point**

The following figure defines the states of the connecting connection points.



**Figure 4 – States of connecting connection points**

NOTE The connection point must remember its attributes even if the connection is broken or if the CP is closed by the application. This makes it easier to restore the connection afterwards. Note also that errors that inhibit further connection attempts put the connection point into an error state.

### 5.2.5.3 Accepting connection point

The following figure illustrates the states of a CP spawned from a listening connection point. This CP does not use a “created” state as it is derived from the listening CP. The notification given to the higher level application in conjunction with the close of the CP, shall be given while the CP is still in existence. The removal can take place after the notification has been given and processed.

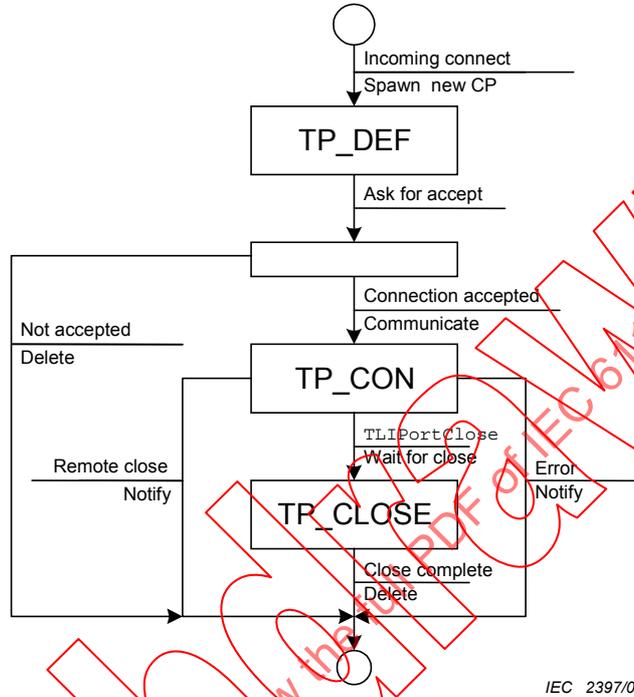


Figure 5 – States of accepting connection point

### 5.2.5.4 States of connectionless connection points

The states of connectionless CP are shown in the figure below. This is similar to the listening CP since this CP is essentially stateless.

Messages can be sent and received in the connected state. The connected state is reached immediately after requesting a port connection.

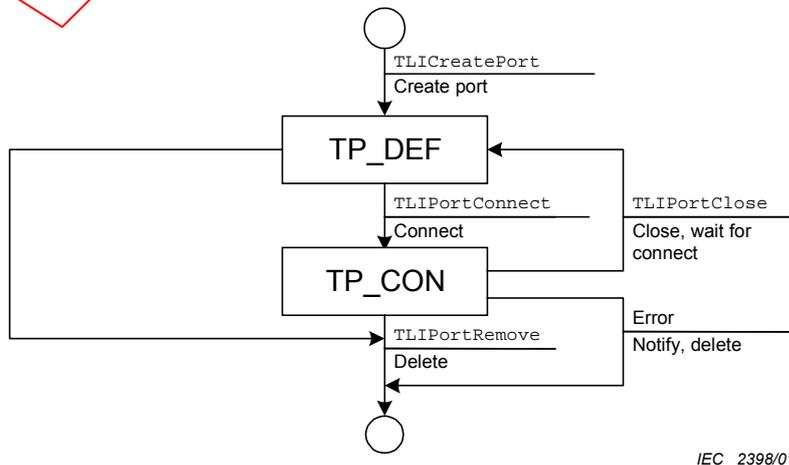


Figure 6 – Connectionless connection point states

### 5.3 TLI service overview

The following subclauses describe the required TLI service primitives. The services described in these clauses are required, but the actual TLI may be implemented by using other (e.g. other functions or names, fewer or more) service primitives.

NOTE The specification of the A-profile is partly modelled on the availability of the services below. It is recommended that an actual implementation of a TLI use these services as they are specified. Otherwise, a rewrite of the A-profile implementation may be required.

#### 5.3.1 General

##### 5.3.1.1 Programming paradigm

The TLI described here is based on an event-oriented principle where no service invocation blocks program execution (no waiting for external events) and where external events are delivered asynchronously to the callers event handler. The service description format follows that specified in clause A.7 of IEC 61162-400.

##### 5.3.1.2 Structural assumption

It is assumed that the TLI itself is a T-profile independent library that manages the interface between the T-profile and the A-profile. The TLI will thus have to be enhanced with TP network specific modules that deliver the actual functionality. This is reflected in the existence of the “add TP network” function.

#### 5.3.2 TLI management services

##### 5.3.2.1 TLI initialization (TLIInitialize)

Request TLI to establish necessary connections to operating system and enable the addition of TP networks. Initialize buffer management and other TP network independent parts of the TLI.

**Pre-condition:** none.

**Post-condition:** TLI ready for use. No TP networks defined.

**Input parameters:**

- upper limit size for buffer pool;
- other implementation dependent configuration limits, for example maximum number of CP.

NOTE To enable the creation of a TLI without using dynamic memory management, it is suggested that the upper limit size is transferred as a memory block of a suitable size. This memory block can then be used by the TLI for its buffer and object store requirements.

**Output parameters:** status of call.

**Call-back:** this call shall block caller until complete. No call back or exception handler is used later. Any later exceptions will only occur relative to TP networks or connection points.

**Errors:** various operating system errors:

- out of memory (application can be asked to give more memory or to close ports);
- various internal failures (bad pointers, inconsistencies).

### 5.3.2.2 Adding a TP network (TLIAddTPN)

This service adds a new TP network to the list of supported TP networks.

**Pre-condition:** TLI initialized.

**Post-condition:** new TP network can be used.

**Input parameters:** descriptor structure that allows TLI to dispatch calls to and from the new TP network modules.

**Output parameters:** TP network identifier.

**Call-back:** TP network related exceptions:

- network warnings (excessive errors, reduction of redundancy);
- complete TP network failure (all CP related to TPN back to defined state)

**Errors:** various operating system errors (e.g. out of memory).

NOTE This service can be implemented by the TLI TP-independent service layer as a way to let new TP networks specify their services that are available. This requires that the user layer calls some kind of initialization routine in the TP network specific module.

### 5.3.2.3 Change or inspect TPN attributes (TLITPNAttribute)

This service can be used to change or inspect certain TPN attributes, for example state (cannot be changed), address attributes, call-back routines, buffer sizes, etc. The attributes available will be dependent on implementation. As a minimum, it shall be possible to inspect state and to retrieve relevant address attributes.

The programmer must take care to check that attribute changes do not cause unexpected side effects, for example when a call-back routine is changed.

**Pre-condition:** TPN is existing.

**Post-condition:** possible change in attributes.

**Input parameters:** read/write flag, attribute code, optionally new value.

**Output parameters:** old value.

**Call-back:** none.

**Errors:** illegal operation, attribute or value.

## 5.3.3 Connection point management

### 5.3.3.1 Create connection point (TLICreateCP)

This service causes a CP description to be created within the TLI. This service creates a general CP object instance on which the rest of the A-profile can operate.

**Pre-condition:** TLI initiated and TP network defined.

**Post-condition:** CP descriptor created in TP\_DEF state.

**Input parameters:**

- TP network;
- service class requested;
- type of port (listening, connecting);
- remote NNN if applicable;
- priority if applicable;
- service instance code or port number if applicable;
- CP and communication (read and/or write) call-back routines.

Additional, possibly required parameters are:

- maximum message size if applicable (if not determined by T-profile);
- buffer parameters if applicable (maximum size, etc.).

Input parameters must be specified as appropriate to the service class that is requested and the type of TP network.

**Output parameters:** CP reference.

**Call-back:** communication and connection notifications when the connection state changes.

Only state the changes which are not directly a result of another TLI service call and which are defined in the state diagrams. Error messages giving causes for unexpected close, for example due to jabbering from the remote side, shall be reported to the TP network call-back and not to the port call-back routine. If the error results in a TP network close, it is not necessary to inform the CP call-back about the close.

**Errors:** bad input parameters, impossible to create descriptor.

### 5.3.3.2 Change or inspect CP attributes (TLICPAttribute)

This service can be used to change or inspect certain CP attributes, for example state (cannot be changed), address attributes, call-back routines, buffer sizes, etc. The attributes available will be dependent on implementation. As a minimum, it shall be possible to change call-back routines, check sizes of input and output buffers and to inspect state.

The programmer must take care to check that attribute changes do not cause unexpected side effects, for example when a call-back routine is changed.

**Pre-condition:** CP is existing.

**Post-condition:** possible change in attributes.

**Input parameters:** read/write flag, attribute code, optionally new value.

**Output parameters:** old value.

**Call-back:** none.

**Errors:** illegal operation, attribute or value.

### 5.3.3.3 Request port connection (TLICPConnect)

This service shall try to establish a connection to or establish as listening, a CP that is already defined.

**Pre-condition:** CP in state TP\_DEF.

**Post-condition:** CP in state TP\_CON, TP\_REQ or TP\_DEF (if error occurred).

**Input parameters:** CP descriptor, time-out for attempt.

**Output parameters:** new state.

**Call-back:** related to state changes out of the new state.

**Errors:** bad descriptor or bad state, attribute errors during connection request (e.g. bad address).

### 5.3.3.4 Request port close (TLICPClose)

Request the close of a CP in open or connecting state.

**Pre-condition:** CP in state TP\_REQ or TP\_CON.

**Post-condition:** CP in state TP\_DEF or TP\_CLOSE.

**Input parameters:** port descriptor, force flush of output data.

**Output parameters:** new state.

**Call-back:** if new state is TP\_CLOSE, a call-back will notify of transfer to TP\_DEF.

**Errors:** bad descriptor or state.

### 5.3.3.5 Request CP remove (TLICPRemove)

Request the removal of a CP in any state. This implies closing the CP and throwing away buffered data. The CP descriptor will be illegal immediately after a successful call.

**Pre-condition:** CP existing.

**Post-condition:** CP removed.

**Input parameters:** CP descriptor.

**Output parameters:** none.

**Call-back:** none.

**Errors:** bad descriptor.

### 5.3.3.6 CP state change call-back (TLICreateCP call back)

This call-back is called when a port changes state independently of an ongoing TLI service call (e.g. open, remote close, error). This will normally be a signal to the application to initiate new operations or new state changes.

**Pre-condition:** CP existing, the state has changed due to an external event.

**Post-condition:** usually the same as the state before call, i.e. the TLI shall set the new state before the call-back routine is activated.

**Input parameters:** CP descriptor, event code.

**Output parameters:** none.

**Call-back:** not applicable.

**Errors:** not applicable.

### 5.3.4 Data management

Data can be transmitted on a CP or received from a CP when it is in the TP\_CON state. The following services are used to send or receive data.

#### 5.3.4.1 Buffer handling

The TLI must provide mechanisms for input and output buffer management. This standard does not specify how this shall be done. In the following, it is assumed that a read from port returns a buffer, and that a write to a port consumes a buffer. It is also assumed that each read event delivers one data buffer and that the calling application (A-profile) can handle cases where write calls are not able to output data immediately.

The implementation must consider efficient buffer management (minimize copying), buffer pool management (dynamic or static allocation of buffers) and operating system related issues (e.g. allocation in the write sub-system and de-allocation from protocol level).

In the following service specifications, it is assumed that a read or write buffer follows the relevant events or service calls and that the ownership of this buffer follows the flow of control. It is also assumed that creation of write buffers and deletion of read buffers must be done explicitly.

#### 5.3.4.2 Read from CP (TLICreateCP call-back)

This service is implemented as a call to the port call-back routine when one buffer is ready for reading. This applies to both stream and message data. When this event is delivered, the corresponding input buffer should follow the event. If this is the case, the buffer must normally be freed by the application after processing.

**Pre-condition:** CP open with read data in a buffer.

**Post-condition:** no change, data buffer follows event.

**NOTE** It is assumed that the higher layer application is the new owner of the read data buffer and that it must explicitly free the buffer, either in the event handler or somewhere else.

**Input parameters:** CP identifier and read buffer.

**Output parameters:** not applicable.

**Call-back:** not applicable.

**Errors:** not applicable.

#### 5.3.4.3 Release read buffer (TLIDeleteBuffer)

This service deletes a read (or pre-allocated write) buffer that has been owned by the higher layer application.

**Pre-condition:** buffer owned by application.

**Post-condition:** buffer released to TLI pool.

**Input parameters:** buffer.

**Output parameters:** not applicable.

**Call-back:** not applicable.

**Errors:** not owner of buffer.

#### 5.3.4.4 Get write buffer (TLIAllocateBuffer)

This service creates a buffer that can be used for formatting of an output message.

**Pre-condition:** TLI initialized.

**Post-condition:** buffer owned by application.

**Input parameters:** maximum size of buffer.

**Output parameters:** buffer.

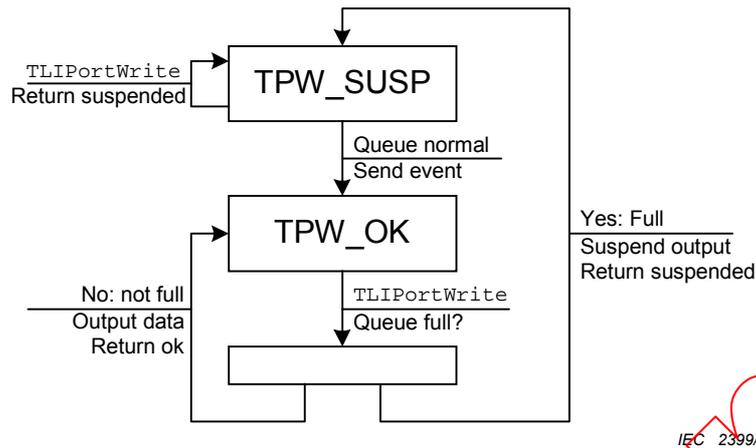
**Call-back:** not applicable.

**Errors:** buffer pool empty.

#### 5.3.4.5 Write data to a CP (TLICPWrite)

This service is used by the application to request the sending of a formatted stream of octets or a message out from a specified CP.

The service will normally queue data for output, but if queues are excessively long and the CP has a low priority, the calling application can be instructed to stop sending data until notified. In this case no part of the buffer will be queued for output. The notification that write can be attempted again is through a call-back from which this service can be invoked again. This is illustrated in the figure below, which represents sub-states of port state TP\_CON.



**Figure 7 – Write sub-states**

**Pre-condition:** CP is open and ready to accept write data. A buffer has been formatted with output data.

**Post-condition:** CP has data queued for output or port in write suspended sub-state. In the former case, the buffer has been consumed.

**Input parameters:** buffer (including length of data), port.

**Output parameters:** operation status (full – wait for call-back or ready for more).

**Call-back:** if state is full, call-back when CP state is ready.

**Errors:** bad CP or buffer.

#### 5.3.4.6 CP write call-back (TLICreateCP call-back)

This call-back is used to signal that the application can write more data to the port.

**Pre-condition:** CP output queue ready.

**Post-condition:** buffer consumed.

**Input parameters:** CP identifier, event code.

**Output parameters:** not applicable.

**Call-back:** not applicable.

**Errors:** not applicable.

## 6 T-profile for redundant Internet

### 6.1 Introduction

#### 6.1.1 Rationale for the use of Internet protocols

IP, currently in version 4, provides message and stream transmission services over any kind of LAN, MAN or WAN networks. The IP protocol is independent of the underlying network technology such as Ethernet, FDDI or ATM, thus enabling interoperability among these different network technologies and providing a safe migration path for the use of future technology.

The IPV4 transport profile specified in this clause relies on the IP network protocol and will also be independent of the underlying network technology.

This approach allows flexibility towards the choice of the network technology and ensures longevity of the system level network implementation due to, for example evolution of the network technology. The transport profile implementation will also take advantage of the many different commercial implementations that are already available on the market. This approach also facilitates the use of widespread Internet applications such as network management, time distribution and terminal emulation. In addition, the transport profile will benefit from the new emerging mechanisms and services that have been in development for several years by the Internet engineering task force (IETF), for example IP version 6 (IPV6).

#### 6.1.2 Relationship to basic Internet protocols

The T-profile is based on the Internet transport profile that provides:

- a connection-oriented stream transfer protocol: the transmission control protocol (TCP);
- a connectionless mode message transfer protocol: the user datagram protocol (UDP).

TCP supplies a robust service of transport including the following major properties:

- end to end transfer;
- full duplex data exchange;
- error detection and retransmission,
- guarantee of in-sequence delivery without duplication;
- flow control;
- congestion control.

TCP supports only a unicast transfer mode.

UDP provides a best-effort strategy for message delivery (without guarantee). UDP supports either a unicast, multicast or broadcast transfer mode.

Both TCP and UDP protocols support client-server type communication.

Internet transport protocols satisfy a subset of the required communication services. Unreliable short message transfer can be achieved through UDP protocol while reliable stream transfer can be achieved through TCP protocol. Reliable message transfer can be achieved through an additional protocol built over TCP. Redundancy can be achieved with the use of two independent networks (transmission channels). Additional protocols are also required for this. Neither TCP nor UDP provides priority. A simple priority mechanism will be provided through the use of the TCP urgent mechanism.

### 6.1.3 Rationale for use of Ethernet

Ethernet has been defined as the physical and data-link layer for this standard due to its wide use in the industry, its availability for various hardware platforms and its relatively high robustness. It has also been a major factor that Ethernet is a typical mainstream technology that is continuously being developed and enhanced.

Ethernet is not able to provide quality of service control and it uses a stochastic network arbitration scheme. These factors mean that it cannot be used for hard real-time control where accurate prediction of network latency is required. In general, the latency and performance of the Ethernet based network will depend on the load condition of the network.

For system integration purposes, where required latency is above 100 ms and no hard real-time deadlines are needed, Ethernet should perform well under most reasonable load conditions. However, care must be taken so that the load is not excessive in any operational state.

When switched Ethernet is used, many of the above assumptions become void. Switched Ethernet provides a full cross-bar switch between all host computers connected to a switch and will not normally use the stochastic medium access mechanism. This means a much higher degree of determinism in the network. However, due to the relatively high complexity of modern workstations, the system integrator should still not rely on deterministic end to end delay for this protocol. Unless special precautions are taken, the limit of 100 ms as the shortest expected response time should be observed.

## 6.2 Transport profile architecture

### 6.2.1 General

The transport communication, time and management services will be provided through the use of a set of protocols, based on Internet protocols, located either at the OSI communication layers (layers 3 to 4) or at upper layers (layers 5 to 7).

The general PISCES protocol architecture at the system level network, using the T-profile defined in this part of IEC 61162, is illustrated in the following figure. The shaded A-profile shows that it is not covered in this part of IEC 61162 but is still part of the complete protocol architecture.

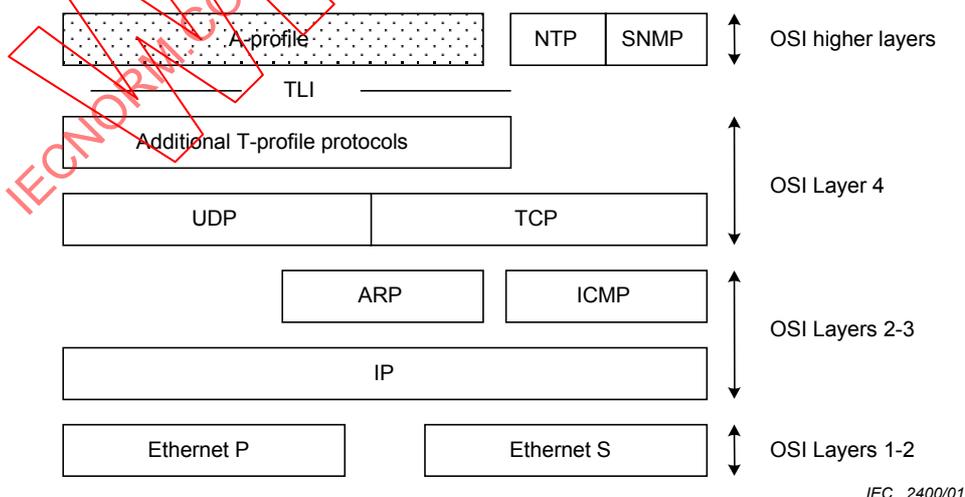


Figure 8 – General transport profile architecture

This architecture satisfies the needs currently identified at the system level network aboard a ship and also makes allowances for the future needs associated with enhanced multimedia applications.

### 6.2.2 Redundancy

Redundancy is implemented by using two physical network interfaces and two independent Ethernet (primary and secondary). Each of these networks has different Internet network addresses. However, only one Internet protocol stack is normally used in the host computer. This is made possible by using the multi-homing facilities of the Internet protocol stacks.

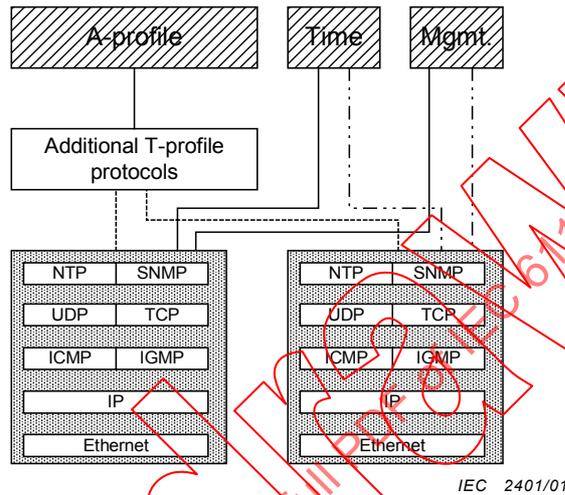


Figure 9 – Redundancy implementation

Higher level protocols (NTP and SNMP) will operate independently on the two networks with no modification of basic protocol or services. This means that the time and network management services will need to be aware of and interface to both networks, i.e. regard each host computer as having two distinct network addresses, one on each network.

The A-profile will make use of additional T-profile functionality as described in this subclause. This means that the A-profile (and, thus, the LNA and MAU) will relate to only one logical network (using one IP network address and one network node number for each host computer).

The additional T-profile services will cover both UDP and TCP. The basic principle is that all data is sent on both physical networks simultaneously and that duplicate data is filtered out at the receiving side. This provides zero switch-over time between the two networks and continuous availability.

Multicast/broadcast services will use UDP broadcasts. Each physical network must be set up so that broadcasts can reach all nodes on the network, for example any switches and routers should normally not do broadcast filtering. To handle duplicate messages, a sequence number will be added to all UDP messages so that the last arriving of the duplicate messages can be removed. This will also ensure that all messages arrive in sequence.

Reliable stream services will use an octet count to do duplicate removal. Reliable message based traffic uses the same principle and message assembly is done only after the duplicate octets have been removed. After errors on one network, where redundant communication is resumed at a later stage, the T-profile uses a special protocol to synchronize octet streams after reconnect.

## 6.2.3 Communication services

### 6.2.3.1 TP network types

The table below lists the TP networks that are defined in this clause. These TP networks are identical except that the first group has predefined IP network addresses. The second form allows user specified network addresses. Note that the second form shall be able to handle all classes of IP networks while the first is defined as a class B network.

The TP networks are divided into MAU-MAU (stream) and LNA-LNA (message) type networks. Note that connection points derived from the stream type networks require a 16-bit port number as additional address information which is not included in the network address format. For the message type, the TCP/UDP port is implicit in the selected service.

**Table 8 – Redundant TP networks**

Code	TP network address format	NNN address format	Notes
LNADR_IPV4R (129)	IPV4 standard redundant address for LNA. The address field is 0 octets long, being void.	The host Internet address with the network part, all zeros.	The two networks are the ones defined in 3.33.
MMADR_IPV4R (131)	IPV4 standard redundant address for MAU. The address field is 0 octets long, being void.		The two networks are the ones defined in 3.33.
LNADR_IPV4RE (130)	IPV4 free redundant address for LNA. The address field is 8 octets long, giving the Internet network addresses in binary format, primary network first, then secondary. The node number part of the addresses shall be all zero bits.		The two networks (of the same class) shall be as defined in 3.33
MMADR_IPV4RE (132)	IPV4 free redundant address for MAU. The address field is 8 octets long, giving the Internet network addresses in binary format, primary network first, then secondary. The node number part of the addresses shall be all zero bits.		The two networks (of the same class) shall be as defined in 3.33.

NOTE An implementation may subdivide one or more of these TP networks into several network modules. This may be done, for example to simplify or modularize implementations of a T-profile (see also 5.2.1).

### 6.2.3.2 Services supported

The services defined in 4.3.3 shall be implemented, with the exception of the MAU-LNA message service. The services may, however, be implemented in a number of different modules where suitable service sub-sets can be linked into the application, either a MAU or an LNA. This may mean that several different TP network instances use the same network code.

## 6.3 Message formats

### 6.3.1 General format

The A-profile (IEC 61162-401) defines a general message format that shall be used by all the A-profile protocols. This message format will also be used by this T-profile. The general message header is shown in table 9. Refer to IEC 61162-401 for a detailed discussion of the different components of the message.

**Table 9 – General message format**

Octet #	Data type	Field description
0 – 1	word16_m	MV4_MAGIC
2 – 3	word16_m	Message type
3 – 4	word16_m	Priority
4 – 5	word16_m	Total length of message in octets (x+1)
- x		Data

### 6.3.2 Special reliable link messages

To synchronize the two channels of a redundant link and to inhibit the creation of more than one communication channel in a peer-to-peer link, a special message is sent ahead of all other traffic. This message is described in table 10.

**Table 10 – Special reliable hello message format**

Octet #	Data type	Field description
0 – 1	word16_m	MV4_MAGIC
2 – 3	word16_m	LL_HELLO (0x0001)
4 – 5	word16_m	Priority
6 – 7	word16_m	Total length of message in octets (16)
8 – 9	word16_m	Type of network
10 – 13	word32_m	NNN of sending host computer
14 – 17	word32_m	Next octet number

The message identity code has the value one. The range from 1 and up to 255 is reserved for management messages.

The priority is not used and can be discarded. The total length is fixed at 16 octets.

The type of network is a code specifying the TP-network used by the sending host computer as defined in table 8. It can be used to detect non-redundant hosts on a redundant network.

The NNN is the network node number of the transmitting host computer. It is used to break ties if two host computers simultaneously act as client and server during a peer-to-peer connection establishment.

The next octet number is used to synchronize octet streams when a redundant link is re-established after a break in one of the communication channels. It is not used for non-redundant networks. It shall be zero for the first connection attempt. When a broken channel is reconnected it shall contain the sequence number of the data octet immediately following the LL\_HELLO message. The first data octet transmitted on a channel is counted as zero (consistent with the zero value of the octet number in the initial message).

### 6.3.3 Broadcast messages

To filter out duplicate messages, the NNN together with the sequence number is added to the broadcast messages by the T-profile. This is done as a prefix to the message header shown in table 9. To maintain compatibility with non-redundant TP-networks in the A-profile, IEC 61162-401, these extra octets are not counted in the total message length. This is indicated by counting octets from -6. Thus, the complete message format for broadcasts is as shown in table 11.

**Table 11 – Broadcast message format**

Octet #	Data type	Field description
-6 – -3	word32_m	NNN of sending host computer
-2 – -1	word16_m	Sequence number for this broadcast port
0 – 1	word16_m	MV4_MAGIC
2 – 3	word16_m	Message type
4 – 5	word16_m	Priority
6 – 7	word16_m	Total length of part 401 message in octets (x+1)
- x		Data

The NNN is the network node number of the sending host. It is used to differentiate between host computers. The sequence number is used to identify messages from one host computer being sent on two redundant channels.

The T-profile shall strip the message of the leading NNN and sequence number before passing it up to the higher layers. The T-profile shall also check message consistency (magic number and message length).

## 6.4 Reliable stream protocol

### 6.4.1 Introduction

The reliable stream protocol is also the basis for the reliable message protocol. It uses a TCP/IP connection between the two communicating host computers and provides a reliable, sequential and bi-directional octet stream.

### 6.4.2 Connection management

#### 6.4.2.1 Difference between peer-to-peer and client-server

The reliable stream protocol shall support two connection modes: Peer-to-peer and client to server. Peer-to-peer communication is established between a client (connecting) and a server (listening), but both parties may simultaneously act both as client and server. In the latter case, the two established connections shall be merged into one by the T-profile. In the client-server model the two links shall remain distinct.

As client-server connection, establishment can in this respect be looked at as a sub-set of the peer-to-peer establishment. This T-profile also implements the two services this way: A CP flagged as a peer-to-peer type CP will be subjected to additional processing to make sure that only one communication link is established between the parties. This processing is not used on a client-server type CP.

The next subclause gives details on the establishment of peer-to-peer connections. The client-server establishment uses normal TCP/IP mechanisms with a listening and a connecting CP. Redundancy is handled by using the same mechanisms as the peer-to-peer connection (see 6.4.2.3 and 6.4.2.4).

#### 6.4.2.2 Resolving peer-to-peer connection conflicts

Each of the two channels shall be established independently and concurrently. The below message sequence diagram shows the worst case exchange of messages when two hosts (1 and 2) simultaneously tries to connect to each other. For simplicity the connection attempt has been drawn so that host 1 connects from network number 1 while host number 2 connects from network number 2. The normal case will be that both attempts start on the primary network. Either case shall not influence the final result.

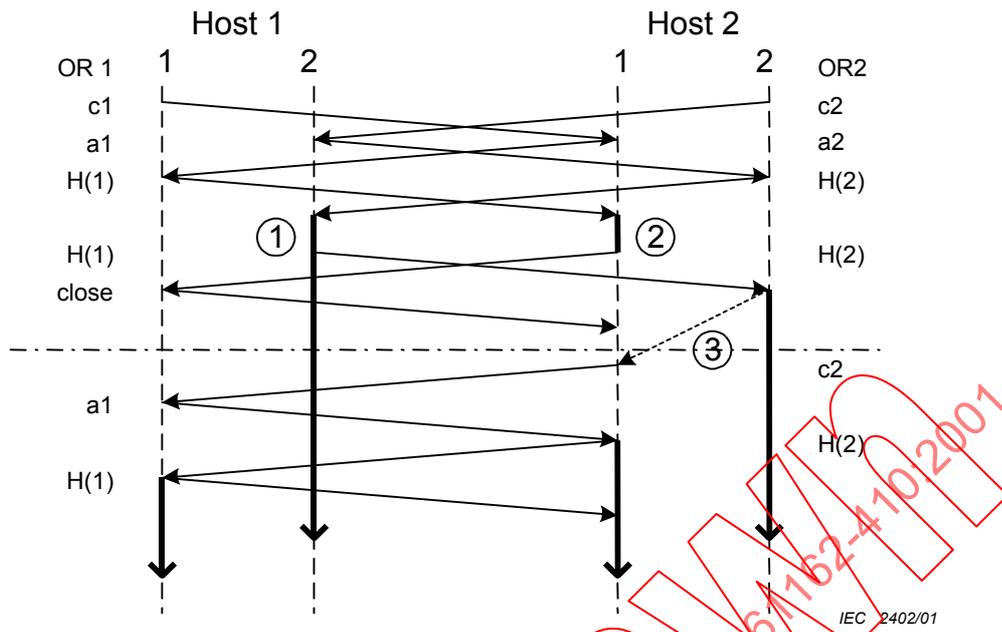


Figure 10 – Concurrent connection attempts

The initial events are that respectively Host 1 and Host 2 gets an open request from a higher layer (OR). This results in the attempt to create a TCP connection (c). On the remote side, both parties will respond with an accept (a). Each host then sends a hello message (H) with its respective host identity. The rest of the processing is explained in the following paragraphs, each referenced to the marks in the figure.

Item 1: Host 1 recognized that the incoming message is from a host with higher priority and accepts the connection. It sends its own host identity as acknowledgement and waits for the higher priority host to clean up links. It also recognizes that it itself has a connection attempt going towards the same host and leaves this to die by itself. The higher layer is, however, notified of a successful connection (although still not fully redundant) and return a connection point ID.

NOTE It is suggested that the transport network module returns the CP ID of the original outgoing connection attempt as this CP already is understood by the higher layer as going to host 2.

Item 2: Host 2 sees the incoming message as from a lower priority host, recognizes that it itself has started to connect to the same host and sends its own identity. It leaves it to the remote host to close the link. It is necessary to wait until the link closes so that this side is sure that the final message was received by the remote host.

Item 3: Host 2 finds that the first link is complete when the final acknowledgement is received from host 1 and notifies the higher layer about the successful (although not yet redundant) connection. It also starts the establishment of the secondary link, which completes without any tests.

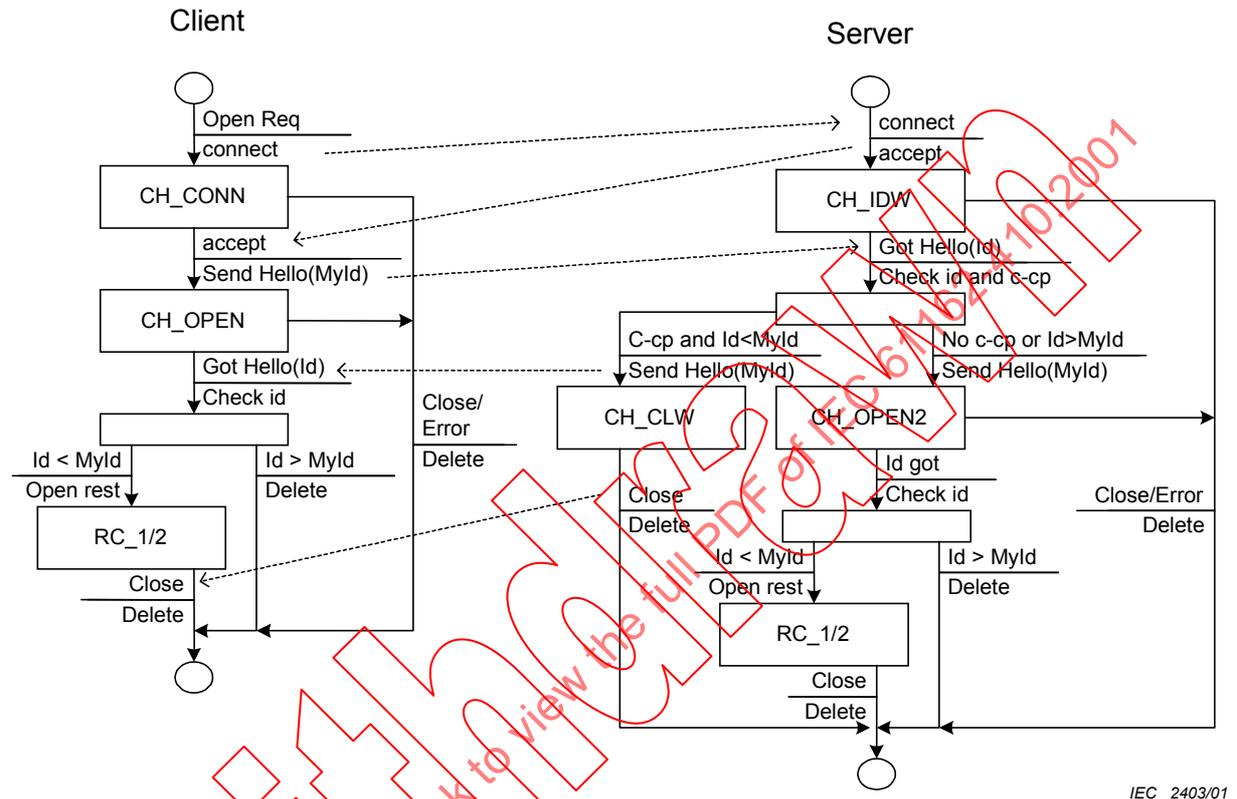
The area below the dashed horizontal line represents the handling of redundant link establishment, which is described in the next clause.

The state changes for each channel are described in the diagram below.

The left side shows the state changes for the client side, while the right side shows the state changes for the server side. Prior to starting the client process, the client shall check if the link is already open. The identity (MyId for local host computer or Id for remote host computer) is

the network node number. It is guaranteed to be unique and a simple check on relative numeric value can be used to break a tie when both sides try to initiate a connection. The result shall be that the side with the highest NNN shall succeed in being the client. The LL\_HELLO message is used to send the identity codes.

The dashed arrows show interaction between server and client side. These can also be found as arrows in the message sequence diagram above.



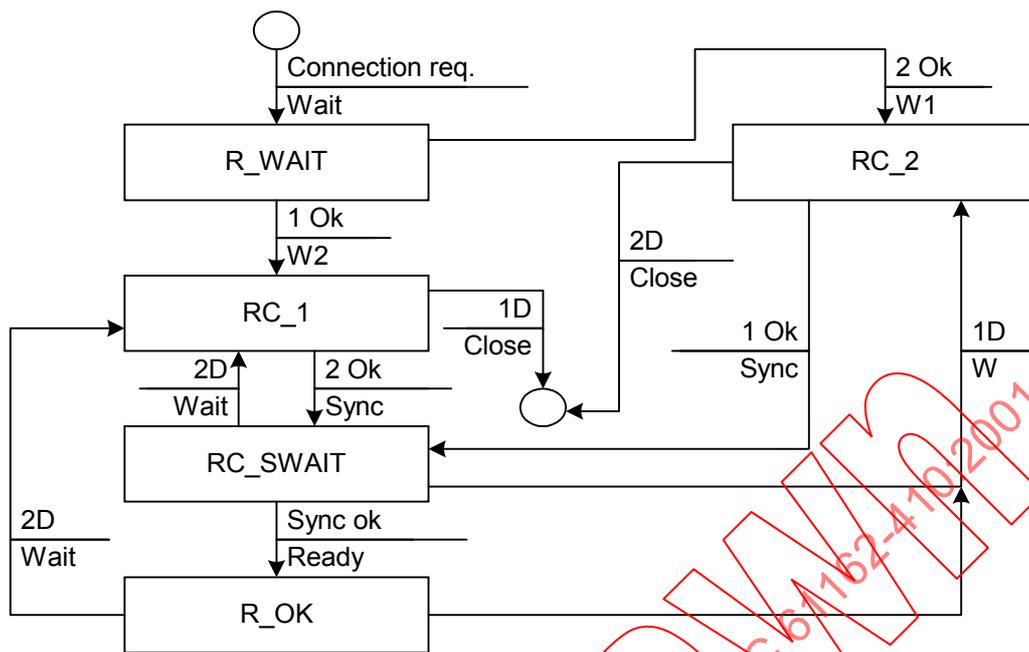
IEC 2403/01

Figure 11 – Connection establishment

### 6.4.2.3 Redundancy modes

The T-profile shall provide graceful degradation, i.e. it shall maintain communication even if one of the redundant channels goes down. The following state diagram shows the possible connection states.

When a dead link is re-established, this shall always be done by the higher numbered host computer (refer to message sequence diagram in figure 11). The exception is when both links go down and the connection is closed to the higher layers.



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Figure 12 – Redundancy communication modes

The events are channel down (e.g. 1D) or channel just connected to (e.g. 1 OK).

The states are:

- R\_WAIT: no communication established. Waiting for connection on both channels. This corresponds to states CH\_xx in figure 12;
- RC\_1 and RC\_2: connection established on primary or secondary channel respectively. Data transfer enabled on the connected channel;
- RC\_SWAIT: both channels connected to data reception enabled on the channel corresponding to the state just left (RC\_1 or RC\_2). Data is transmitted on both channels. Waiting for data synchronization on the other channel before full redundant reception is enabled;
- R\_OK: full redundant communication established.

The synchronization wait (RC\_SWAIT) is necessary to check that the remote host computer actually is transmitting correct information on both links before the second link is accepted as a full replacement for the first established link.

In all communication modes, data is transferred to the higher levels as soon as they have been received. The T-profile shall not delay data.

#### 6.4.2.4 Channel synchronization requirements

The T-profile shall provide a mechanism for consistency checks between the two channels. Different mechanisms can be employed, but the mechanism shall implement a check to find (with a suitable probability) the following errors in one of the channels.

- a) It shall be able to detect the omission of a data block.
- b) It shall be able to detect the repetition of a data block.
- c) It shall be able to detect the wrong ordering of two data blocks.
- d) It shall be able to detect the insertion of a number of random octets, for example some number of zero value octets.

These errors are the most probable that can occur in the sending T-profile. It is not necessary to detect errors due to transmission problems on the physical network. This is reasonably well handled in the TCP/IP and Ethernet layers.

NOTE 1 Requirement 3 (item c) above) is probably the most difficult to handle with a reasonably effective check. It requires relatively fine-grained checkpoint intervals or some form of positional memory (e.g. a CRC) in the algorithm.

Any errors detected by the synchronization algorithm shall immediately cause the communication link (both channels) to be closed and an exception to be generated for the higher layers. It is the responsibility of the higher layers to reconnect the link.

NOTE 2 The link shall also be closed when an error occurs during reopening of a redundant link in the RC\_SWAIT state. It is generally not possible to prescribe a safe action in this case, even when one channel is operating.

The LL\_HELLO message contains the necessary information for restoration of a broken redundancy link. It contains the octet sequence number for the immediately following data.

Annex B provides an example of a possible implementation of such an algorithm.

#### **6.4.2.5 Priority**

TCP/IP can to a certain degree support two priority levels: Normal and urgent. Urgent priority is achieved by setting the out of band (urgent) flag after having transmitted a block of urgent data. Normal priority data shall not use the out of band flag.

NOTE The actual effect of using out of band flags to implement priority will depend on the implementation. The TCP/IP stack is required to make a best effort to send this data before other data, but implementations vary in how they achieve this.

To make this mechanism work, only urgent messages shall be sent on an urgent link and only normal priority messages on a normal priority link. This means that a link that shall support both normal and urgent priority data (e.g. the LNA-LNA message link) must consist of four TCP/IP connections: Two for redundancy in each of the two priority links. Each pair of connections (normal or urgent priority) can be handled as one link on which the T-profile sends a certain sub-set of data from the higher level CP.

Data links that shall support only one priority level (urgent or normal, e.g. for stream links) should normally not create the other connection. The parameters determining the priorities to be supported are input to the TLI.

Low priority cannot be supported by TCP/IP as only two priority levels exist.

#### **6.4.3 Data transmission**

##### **6.4.3.1 General**

Data is transmitted on the active links as soon as they are open. This means that data transmission shall start in RC\_1 and RC\_2 modes and that a synchronization is necessary when fully redundant transmission starts at some later stage.

The T-profile must keep track of the octet count used for transmission so that the correct LL\_HELLO message can be sent when a link opens. It must also make sure that the transmitted octet count is consistent with the actual octets transmitted.

The T-profile shall generally send data as soon as data is requested sent by the higher layers. The T-profile may, however, buffer data internally or refuse transmission if the data channel is congested.

### 6.4.3.2 Transmission of urgent data

The urgent flag shall be set after each complete block of data handed over to the TCP layer, for example after each write call.

### 6.4.3.3 Graceful degradation between priority levels

If one of the links (urgent or normal) fail, transmission can continue on the other link provided that the following requirements are fulfilled:

- normal data shall not be transmitted on the urgent link (signal normal link failure to higher layers);
- urgent data can be transmitted on the normal link, but no normal messages can be transmitted on this link when urgent messages are transmitted (signal normal link failure to higher layers);
- connection re-establishment shall be attempted immediately;
- as soon as both links are re-established, the urgent messages shall be redirected to the urgent link and normal link functionality resumed.

Graceful degradation is only provided between urgent and normal links established at the same time. The T-profile shall not reallocate link descriptors, for example to give a low priority stream to a normal priority message link. This shall be handled, if necessary, by the higher layers.

### 6.4.3.4 Local link close

The link close command from the higher layer shall cause all data still being in the communication pipeline to be transmitted before the link is physically closed, unless a throw flag has been included in the close command. In the latter case, the link shall be closed immediately. Both redundant channels shall be closed at the same time. This may mean that the link cannot be reopened immediately after a close (if buffers need to empty). The throw flag can be used to enable immediate reopening.

NOTE The normal TCP/IP close will normally allow buffers to empty before the close is completed. The throw flag requires special handling. This has to be checked during implementation. Note also that TCP/IP sockets may linger and that this may have to be handled in the implementation by setting, for example no-linger flags.

### 6.4.3.5 Remote link close

The same provisions as for the local close also applies to remote link close. The CP may or may not have to flush its output buffers. Note in particular that data may be arriving although the connection does not allow the transmission of data (to empty the sender buffers).

NOTE Errors in the communication layer may be detected only during write operations for some error modes. Note also that the keep-alive timer used in TCP/IP is not very useful for control systems. The time-out is typically several minutes or, sometimes, hours.

## 6.4.4 Data reception

### 6.4.4.1 General

Data received shall be transferred to the higher layers as soon as they have been received on one of the open channels. The synchronization mechanism described in 6.4.2.4 shall be used to check consistency between channels, but the mechanism shall not delay the transfer of data to higher layers.

#### **6.4.4.2 Flow control**

The receiver may employ normal TCP/IP flow control mechanisms if a link becomes congested, i.e. stop doing read on the relevant connection point. This is, however, not legal on urgent links. These shall be processed immediately and any congestion control must be performed on other link types.

#### **6.4.4.3 Priority**

Urgent data received on an urgent link must process the out of band flag. This means that, for example urgent signals from the operating system must be trapped. It also means that all data on an urgent link must be read, up to the last pending urgent flag.

#### **6.4.4.4 Graceful degradation between priority levels**

The receiving end should normally only get urgent messages on the urgent link. However, observe the possible change from normal to urgent in the case of an urgent link failure.

#### **6.4.5 Link close**

The provisions described in 6.4.3.4 apply also on the listening side.

### **6.5 Reliable message service**

#### **6.5.1 Introduction**

The reliable message service shall use the reliable stream protocol and use the general message format (see 6.3.1) to find message boundaries. This means that most services described in 6.4 are implemented in the same manner.

The exceptions are noted in the following subclause(s).

#### **6.5.2 Data reception**

The T-profile is required to deliver only complete messages, one and one at a time, to the higher layers. This may require buffering of data to make sure that a complete message has been assembled.

#### **6.5.3 Data transmission on the urgent link**

The T-profile shall set the urgent flag only on message boundaries. If, for some reason, two or more messages are transmitted immediately after each other, the transmitter may set the urgent flag only after the last message.

### **6.6 Unreliable message service**

#### **6.6.1 Introduction**

The unreliable message service uses the UDP broadcast mechanism. This is a connectionless and unreliable service. However, the TLI shall be defined in such a way that the users of this service can use the normal concepts of connection points and data input and output on connection points.

#### **6.6.2 Connection management**

##### **6.6.2.1 Introduction**

There is no protocol level connection management associated with the unreliable message service other than maintaining and checking sequence number for all messages. However, the

TLI specifies that a CP shall be associated with each listening or transmitting broadcast address.

This standard specifies five different listening UDP ports for broadcast messages. Each of these ports can be associated with exactly one transmitting and one receiving CP.

### 6.6.2.2 Listening CP

The T-profile shall maintain state associated with each remote transmitting host computer that transmits to a specific UDP port. The state is limited to the identity (NNN) of the remote host computer and its current sequence number. When opening a new listening CP, the T-profile shall reset this state information to no remote host computers. As soon as a message from a previously unlisted host computer arrives, the state shall be updated with the identity of that computer and the sequence number of the message.

A remote transmitting host computer may go off and come on line due to a reset or other internal exception. This will cause a reset of the sequence number to one. The receiving end must be able to detect this and continue reporting incoming messages from the new sequence number. In systems that employ routers and alternative transmission paths, one may see duplicates of UDP messages, and hence, out of sequence messages, without an exception on the transmitting end. This situation is extremely unlikely for a LNA type ship control system and this standard prescribes the following handling of out of sequence packets:

- a) out of sequence messages on the same physical network shall normally be regarded as the result of a transmitter exception and lead to a reset of the listener's sequence number counter;
- b) exceptionally, if the newly arriving message has a high sequence number (significantly higher than zero) which is close to the current sequence number, the message may be thrown away.

NOTE "High" in this context should be at least 100. "Close" should be less than 10. There are no firm limits here and the decision is left to the implementation. The duplication situation is much more unlikely than a host computer reset and the latter situation must be given priority.

The standard (IEC 61162-4) defines the multicast service as unreliable and potentially with messages out of sequence, so the A-profile is constructed with these limitations in mind. This means that the alternative two above should be used with great care.

This standard specifies five different listening UDP ports for broadcast messages. Separate state shall be maintained for each port, i.e. distinct sequence number series will be used for each port.

Only one CP is allowed for each UDP port. The T-profile need not support duplication of incoming messages to possibly several ports.

The sequence number for the transmitting host shall be the same for both links, when redundant communication paths are used.

### 6.6.2.3 Transmitting CP

The T-profile shall maintain a sequence number for all outgoing messages on a given UDP port. The sequence numbers shall be separately maintained for each port, but it shall be common to redundant transmission paths for that port. It shall be set to one for all ports when the TP network is established. Closing and opening a CP on a specific port shall not reset the sequence number for that port.

The sequence number shall be incremented by one for each new transmitted message.

#### 6.6.2.4 Redundancy management

This service shall transmit and listen on both redundant networks unless instructed by higher level management services to use only one of the networks. There are no particular connection or redundancy management provisions required for this service.

Although sequence number management shall be distinguished between on UDP ports, but it shall be common to redundant links. More details are given in the following subclauses.

#### 6.6.3 Message length limitation

Outgoing messages are limited in size to that supported by the low-level network. For IP over Ethernet, the maximum transmission unit (MTU) is 1 500 octets (RFC 894). The limit is reduced by various headers as listed in the table below.

**Table 12 – Maximum message size**

Protocol header	Size (Octets)
IP header	20
UDP header	8
T-profile broadcast header	6
Sum	34

The maximum payload for higher layer protocols is 1 466 octets.

#### 6.6.4 Data transmission

Outgoing messages shall be checked for maximum length before they are transmitted. They shall be formatted with an additional header containing the NNN of the transmitting host and a sequence code as described in 6.3.3. The first message on a newly established transmitting CP shall use the sequence code one.

NOTE It is suggested that the TLI uses some form of data buffer management that reduces the need to copy data back and forth to insert the necessary header information.

The formatted message shall be transmitted on both the primary and the secondary network in two identical copies.

The transmitting service is not required to retransmit messages when errors occur locally (e.g. output buffers full). However, the T-profile shall make a best effort to transmit all outgoing data (e.g. by allocating sufficient memory for buffers).

#### 6.6.5 Data reception

Messages received on a listening and redundant CP shall be filtered so that

- the minimum of duplicate messages is delivered to higher levels;
- the minimum of out of sequence messages is delivered to higher levels;
- messages are delivered with minimum delay in the T-profile;
- the receiving part of the T-profile shall cater to the situation where a remote host is reset so that sequence numbers are reset (see 6.6.2.2).

The filtering shall be done based on the NNN of the transmitting host computer and the sequence number. The T-profile shall immediately deliver a message to the higher layer when a message with a higher sequence number is received on one of the two networks, even when this number shows that one or more messages can have been lost. Any lost messages shall be handled by the higher layers.

NOTE 1 The reset of transmitting host causes problems which are “contrary” to the filtering requirements. Priority must be given to the reset issue as that is more likely in most contexts. However, one can examine the sequence numbers on the same network together with those on the redundant network to be able to make a better decision on this issue.

NOTE 2 The sequence number will start at one (see previous subclause). This means that the sequence number of a newly established listening CP can be set to zero to enable the first message to pass through. Note also that sequence numbers are 16-bit and that a wrap around will cause the number zero to be used. The wrap around will also have to be handled in conjunction with message filtering.

NOTE 3 As two independent networks are used, there is theoretically a possibility that out of sequence messages can be received, even on a local area network. This can occur if one or more messages are lost on the fastest transmitting network. It is recommended that this possibility is ignored by the T-profile and that it passes sequential messages (with possible gaps) to the higher layers as soon as possible.

## 6.7 Time services

### 6.7.1 Introduction

This standard uses the Network Time Protocol – NTP (RFC 1305) and the Simple Network Time Protocol – SNTP (RFC 2030) to implement a distributed time reference in the IEC 61162-4 network. A short overview of these protocols can be found in annex C.

### 6.7.2 Time management architecture

This standard requires that each of the two redundant networks shall have at least one time server, each supporting NTP (version 3 or higher). The time servers (program modules) shall reside in host computers that have a connection to a reliable UTC source, for example a GPS receiver.

The two time servers should use independent time sources.

The time servers shall send all multicast messages on both networks and shall also be able to serve unicast or anycast queries arriving at any of the two networks. The time servers shall support all SNTP services used by clients (unicast, multicast and anycast). The two networks shall, otherwise, be regarded as two separate networks with different IP address ranges.

The two time servers should work together in symmetric peer mode (NTP) to back each other up in case the time source for either fails.

NOTE 1 Bringing up a NTP primary server requires a radio or satellite receiver or modem. Support for the commonly available GPS receivers is included in the default configuration of many of the commonly available NTP daemons, for example `ntpd` for UNIX. Refer to the manuals for the various NTP modules for more details on configuration.

NOTE 2 A GPS receiver requires support in the form of a serial port. On a typical work-station the timing jitters contributed by the serial port hardware and software driver can accumulate to several milliseconds. A set of special line disciplines and stream modules can be configured in the operating system kernels in order to reduce these errors. These routines intercept special characters or signals provided by the radio or modem clock and save a local timestamp for later processing.

### 6.7.3 Client interfaces

Clients shall have an interface to both time servers. The client shall select the service level that gives the time accuracy as required by the server. SNTP is normally sufficient for all but the most strict requirements. It is recommended that the client uses the SNTP anycast or multicast mode to avoid unnecessary configuration in the system.

The expected accuracy for the various service levels shall be specified by the system integrator. This will be dependent on accuracy of time sources, of configuration files and of general network configuration.